

# Computer Architecture

## Lecture 6a: RowHammer II

Prof. Onur Mutlu

ETH Zürich

Fall 2019

4 October 2019

# Recall: The Story of RowHammer

- One can **predictably induce bit flips** in commodity DRAM chips
  - >80% of the tested DRAM chips are vulnerable
- First example of how a **simple hardware failure mechanism** can create a **widespread system security vulnerability**

**WIRED**

Forget Software—Now Hackers Are Exploiting Physics

BUSINESS	CULTURE	DESIGN	GEAR	SCIENCE
----------	---------	--------	------	---------

ANDY GREENBERG SECURITY 08.31.16 7:00 AM

SHARE



SHARE  
18276



TWEET

# FORGET SOFTWARE—NOW HACKERS ARE EXPLOITING PHYSICS

# RowHammer Solutions

# Two Types of RowHammer Solutions

---

## ■ Immediate

- ❑ To protect the vulnerable DRAM chips in the field
- ❑ Limited possibilities

## ■ Longer-term

- ❑ To protect future DRAM chips
- ❑ Wider range of protection mechanisms

## ■ Our ISCA 2014 paper proposes both types of solutions

- ❑ Seven solutions in total
- ❑ PARA proposed as best solution → already employed in the field



# Apple's Patch for RowHammer

---

- <https://support.apple.com/en-gb/HT204934>

Available for: OS X Mountain Lion v10.8.5, OS X Mavericks v10.9.5

Impact: A malicious application may induce memory corruption to escalate privileges

Description: A disturbance error, also known as Rowhammer, exists with some DDR3 RAM that could have led to memory corruption. This issue was mitigated by increasing memory refresh rates.

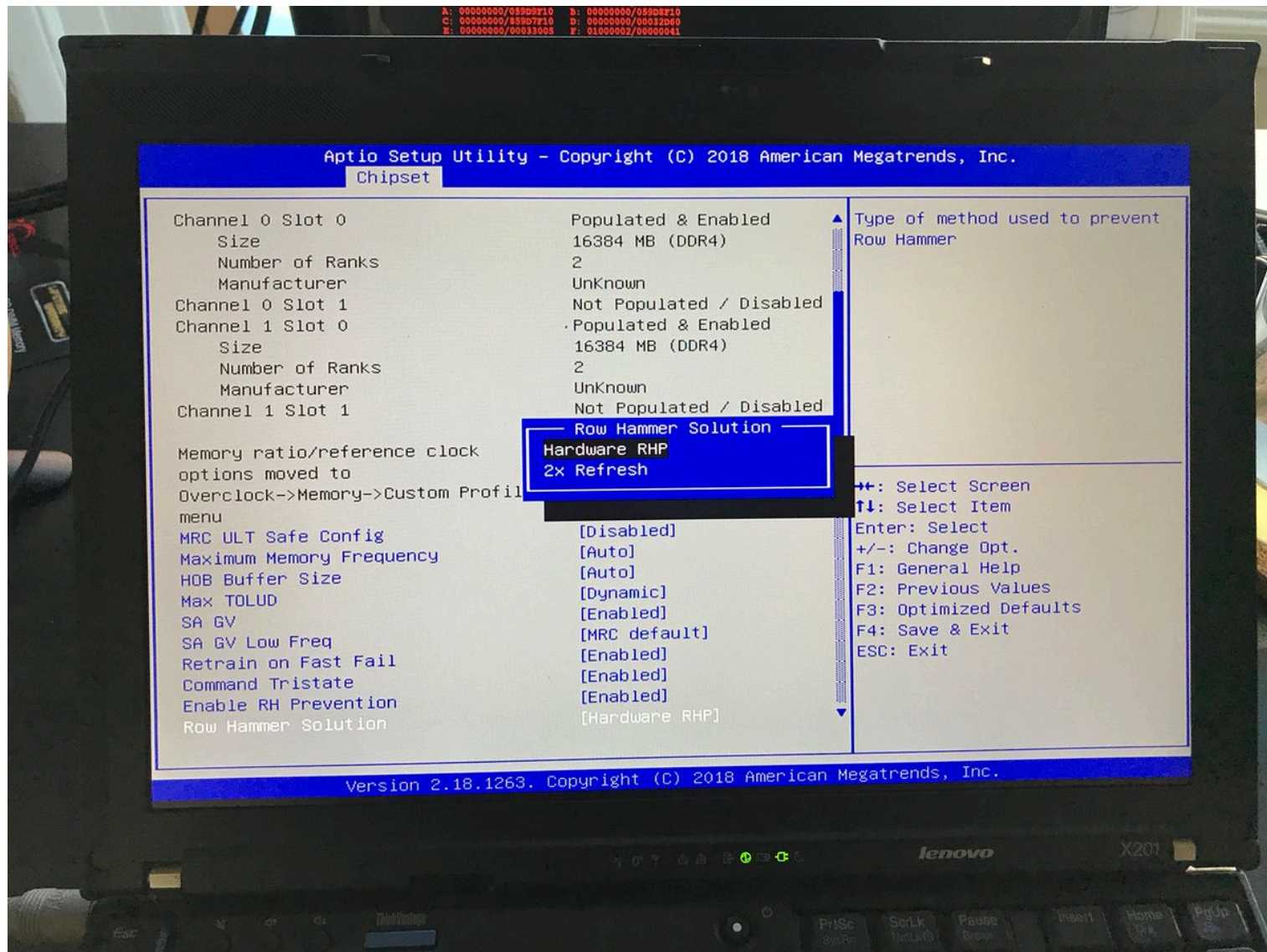
CVE-ID

CVE-2015-3693 : Mark Seaborn and Thomas Dullien of Google, working from original research by Yoongu Kim et al (2014)

HP, Lenovo, and other vendors released similar patches

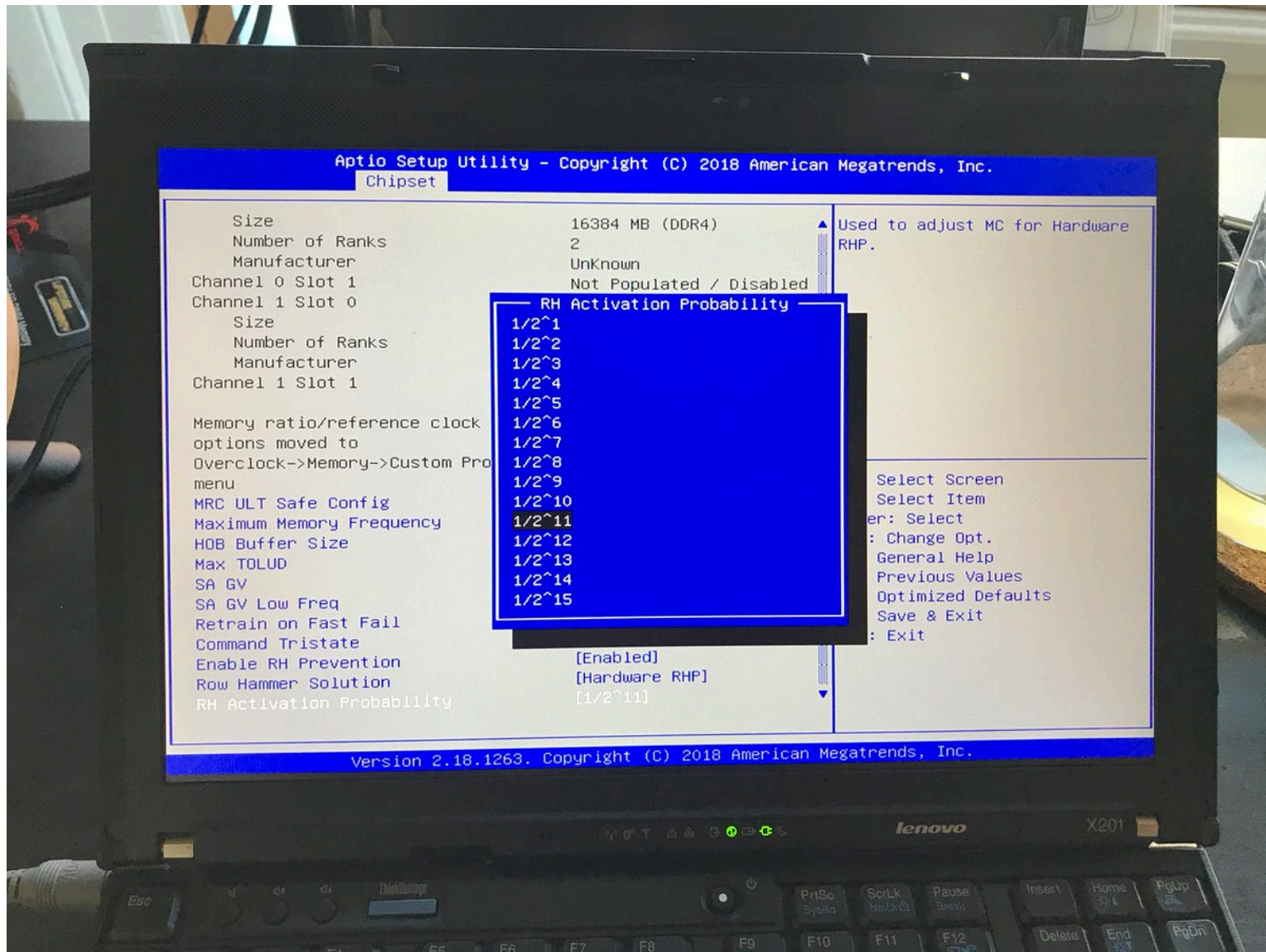
---

# Probabilistic Activation in Real Life (I)





# Probabilistic Activation in Real Life (II)



# Seven RowHammer Solutions Proposed

---

- Yoongu Kim, Ross Daly, Jeremie Kim, Chris Fallin, Ji Hye Lee, Donghyuk Lee, Chris Wilkerson, Konrad Lai, and Onur Mutlu,  
**"Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors"**  
*Proceedings of the 41st International Symposium on Computer Architecture (ISCA), Minneapolis, MN, June 2014.*  
[\[Slides \(pptx\) \(pdf\)\]](#) [\[Lightning Session Slides \(pptx\) \(pdf\)\]](#) [\[Source Code and Data\]](#)

## Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors

Yoongu Kim<sup>1</sup>   Ross Daly\*   Jeremie Kim<sup>1</sup>   Chris Fallin\*   Ji Hye Lee<sup>1</sup>  
Donghyuk Lee<sup>1</sup>   Chris Wilkerson<sup>2</sup>   Konrad Lai   Onur Mutlu<sup>1</sup>

<sup>1</sup>Carnegie Mellon University   <sup>2</sup>Intel Labs

# Main Memory Needs Intelligent Controllers for Security

# Industry Is Writing Papers About It, Too

## DRAM Process Scaling Challenges

### ❖ Refresh

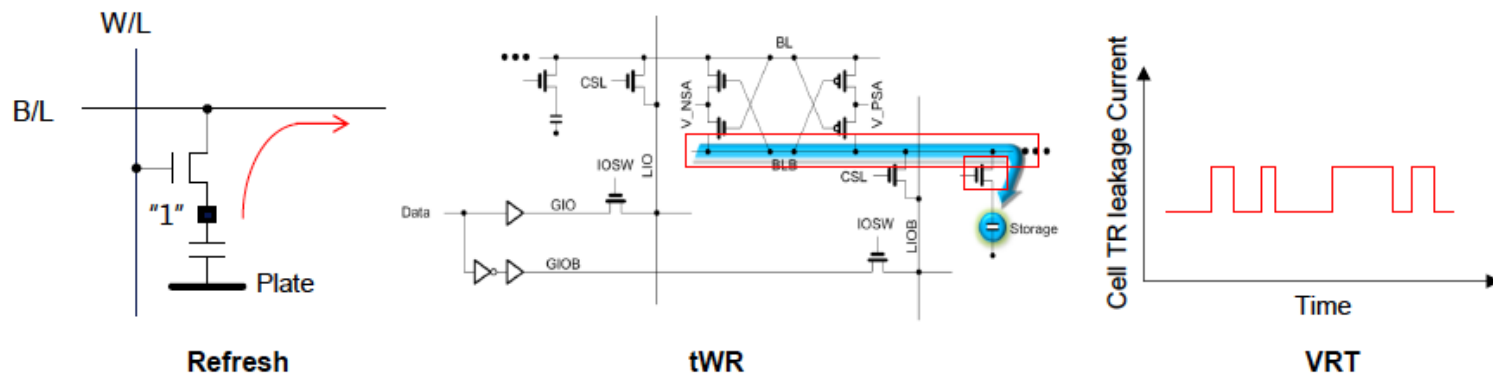
- Difficult to build high-aspect ratio cell capacitors decreasing cell capacitance
- Leakage current of cell access transistors increasing

### ❖ tWR

- Contact resistance between the cell capacitor and access transistor increasing
- On-current of the cell access transistor decreasing
- Bit-line resistance increasing

### ❖ VRT

- Occurring more frequently with cell capacitance decreasing



# Call for Intelligent Memory Controllers

## DRAM Process Scaling Challenges

### ❖ Refresh

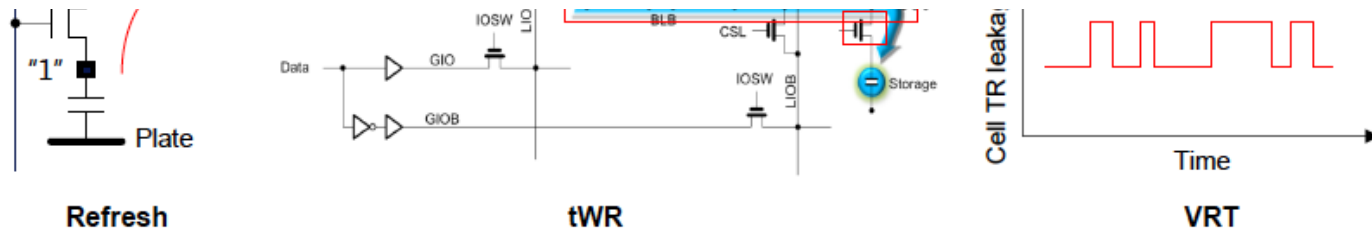
- Difficult to build high-aspect ratio cell capacitors decreasing cell capacitance

THE MEMORY FORUM 2014

## Co-Architecting Controllers and DRAM to Enhance DRAM Process Scaling

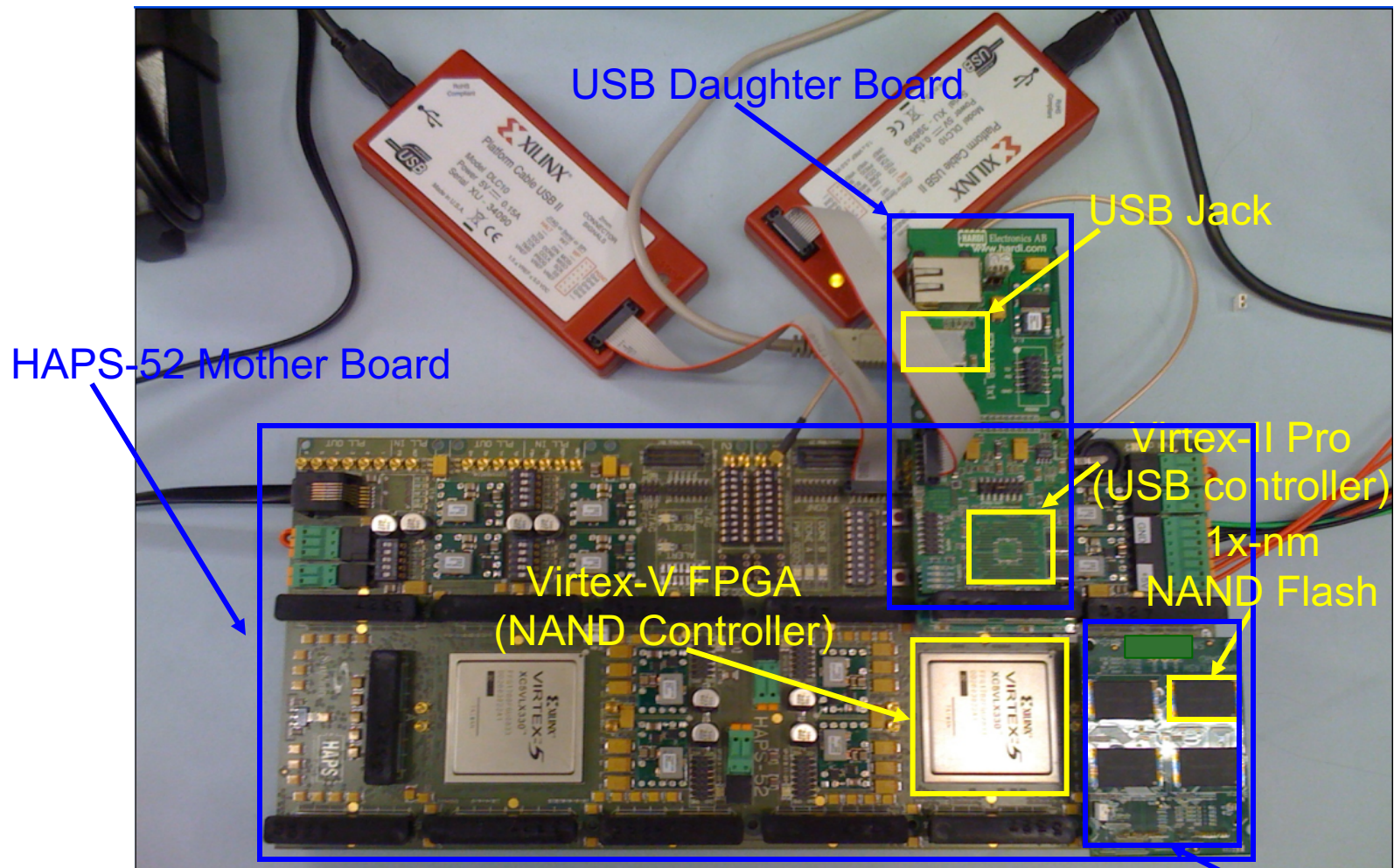
Uksong Kang, Hak-soo Yu, Churoo Park, \*Hongzhong Zheng,  
\*\*John Halbert, \*\*Kuljit Bains, SeongJin Jang, and Joo Sun Choi

*Samsung Electronics, Hwasung, Korea / \*Samsung Electronics, San Jose / \*\*Intel*





# Aside: Intelligent Controller for NAND Flash



[DATE 2012, ICCD 2012, DATE 2013, ITJ 2013, ICCD 2013, SIGMETRICS 2014, HPCA 2015, DSN 2015, MSST 2015, JSAC 2016, HPCA 2017, DFRWS 2017, PIEEE 2017, HPCA 2018, SIGMETRICS 2018]

Cai+, "Error Characterization, Mitigation, and Recovery in Flash Memory Based Solid State Drives," Proc. IEEE 2017.





*Proceedings of the IEEE, Sept. 2017*



## Error Characterization, Mitigation, and Recovery in Flash-Memory-Based Solid-State Drives

*This paper reviews the most recent advances in solid-state drive (SSD) error characterization, mitigation, and data recovery techniques to improve both SSD's reliability and lifetime.*

By YU CAI, SAUGATA GHOSE, ERICH F. HARATSCH, YIXIN LUO, AND ONUR MUTLU

<https://arxiv.org/pdf/1706.08642>

**Main Memory Needs  
Intelligent Controllers**

# Future Memory Reliability/Security Challenges

# Future of Main Memory

---

- DRAM is becoming less reliable → more vulnerable

# Large-Scale Failure Analysis of DRAM Chips

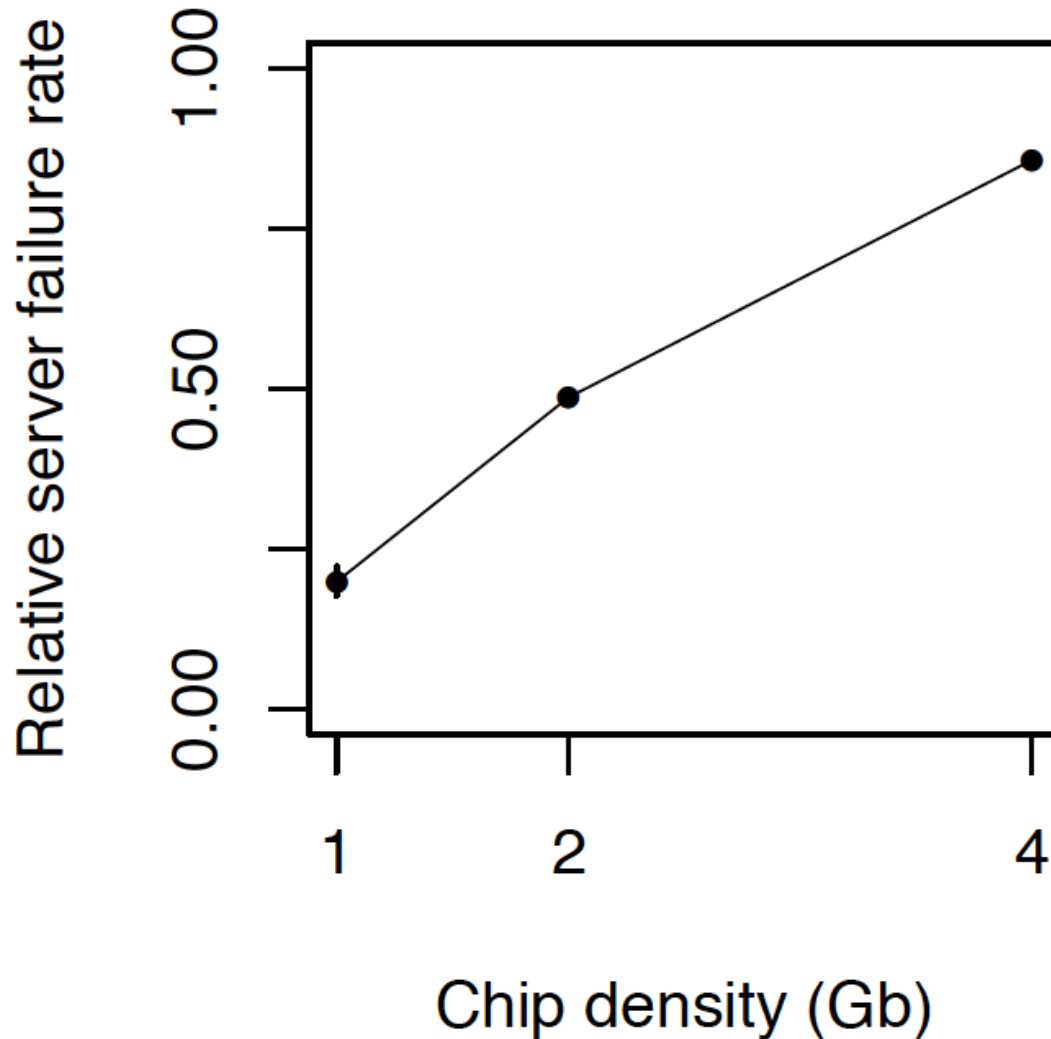
---

- Analysis and modeling of memory errors found in all of Facebook's server fleet
- Justin Meza, Qiang Wu, Sanjeev Kumar, and Onur Mutlu,  
**"Revisiting Memory Errors in Large-Scale Production Data Centers: Analysis and Modeling of New Trends from the Field"**  
*Proceedings of the 45th Annual IEEE/IFIP International Conference on Dependable Systems and Networks (DSN)*, Rio de Janeiro, Brazil, June 2015.  
[[Slides \(pptx\)](#)] [[pdf](#)] [[DRAM Error Model](#)]

## Revisiting Memory Errors in Large-Scale Production Data Centers: Analysis and Modeling of New Trends from the Field

Justin Meza   Qiang Wu\*   Sanjeev Kumar\*   Onur Mutlu  
Carnegie Mellon University   \* Facebook, Inc.

# DRAM Reliability Reducing



*Intuition:  
quadratic  
increase in  
capacity*

# Aside: SSD Error Analysis in the Field

---

- First large-scale field study of flash memory errors
- Justin Meza, Qiang Wu, Sanjeev Kumar, and Onur Mutlu,  
**"A Large-Scale Study of Flash Memory Errors in the Field"**  
*Proceedings of the ACM International Conference on Measurement and Modeling of Computer Systems*  
**(*SIGMETRICS*)**, Portland, OR, June 2015.  
[[Slides \(pptx\)](#)] [[pdf](#)] [[Coverage at ZDNet](#)]

## A Large-Scale Study of Flash Memory Failures in the Field

Justin Meza  
Carnegie Mellon University  
[meza@cmu.edu](mailto:meza@cmu.edu)

Qiang Wu  
Facebook, Inc.  
[qw@fb.com](mailto:qw@fb.com)

Sanjeev Kumar  
Facebook, Inc.  
[skumar@fb.com](mailto:skumar@fb.com)

Onur Mutlu  
Carnegie Mellon University  
[onur@cmu.edu](mailto:onur@cmu.edu)

# Future of Main Memory

---

- DRAM is becoming less reliable → more vulnerable
- Due to difficulties in DRAM scaling, other problems may also appear (or they may be going unnoticed)
- Some errors may already be slipping into the field
  - Read disturb errors (Rowhammer)
  - Retention errors
  - Read errors, write errors
  - ...
- These errors can also pose security vulnerabilities



# DRAM Data Retention Time Failures

---

- Determining the data retention time of a cell/row is getting more difficult
- Retention failures may already be slipping into the field

# Analysis of Data Retention Failures [ISCA'13]

---

- Jamie Liu, Ben Jaiyen, Yoongu Kim, Chris Wilkerson, and Onur Mutlu,  
**"An Experimental Study of Data Retention Behavior in Modern DRAM Devices: Implications for Retention Time Profiling Mechanisms"**  
*Proceedings of the 40th International Symposium on Computer Architecture (ISCA)*, Tel-Aviv, Israel, June 2013. [Slides \(ppt\)](#) [Slides \(pdf\)](#)

## An Experimental Study of Data Retention Behavior in Modern DRAM Devices: Implications for Retention Time Profiling Mechanisms

Jamie Liu<sup>\*</sup>  
Carnegie Mellon University  
5000 Forbes Ave.  
Pittsburgh, PA 15213  
[jamiel@alumni.cmu.edu](mailto:jamiel@alumni.cmu.edu)

Ben Jaiyen<sup>\*</sup>  
Carnegie Mellon University  
5000 Forbes Ave.  
Pittsburgh, PA 15213  
[bjaiyen@alumni.cmu.edu](mailto:bjaiyen@alumni.cmu.edu)

Yoongu Kim  
Carnegie Mellon University  
5000 Forbes Ave.  
Pittsburgh, PA 15213  
[yoonguk@ece.cmu.edu](mailto:yoonguk@ece.cmu.edu)

Chris Wilkerson  
Intel Corporation  
2200 Mission College Blvd.  
Santa Clara, CA 95054  
[chris.wilkerson@intel.com](mailto:chris.wilkerson@intel.com)

Onur Mutlu  
Carnegie Mellon University  
5000 Forbes Ave.  
Pittsburgh, PA 15213  
[onur@cmu.edu](mailto:onur@cmu.edu)

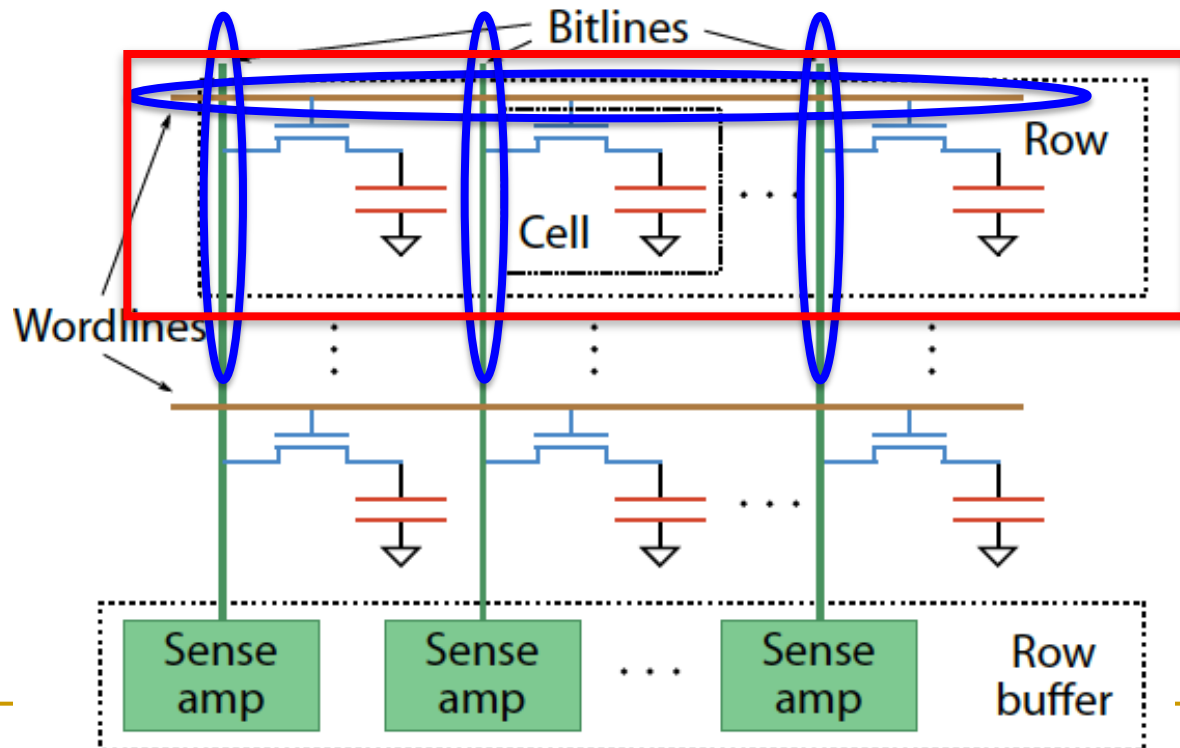
# Two Challenges to Retention Time Profiling

---

- Data Pattern Dependence (DPD) of retention time
- Variable Retention Time (VRT) phenomenon

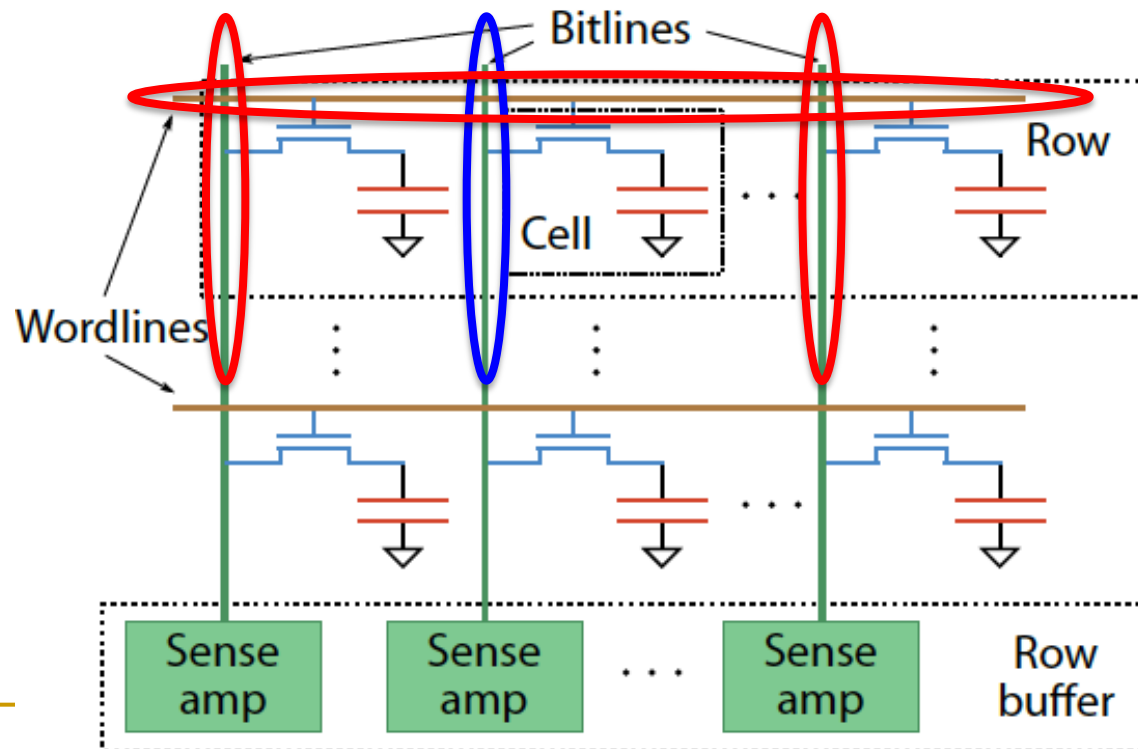
# Two Challenges to Retention Time Profiling

- **Challenge 1: Data Pattern Dependence (DPD)**
  - Retention time of a DRAM cell depends on its value and the values of cells nearby it
  - When a row is activated, all bitlines are perturbed simultaneously



# Data Pattern Dependence

- Electrical noise on the bitline affects reliable sensing of a DRAM cell
- The magnitude of this noise is affected by values of nearby cells via
  - Bitline-bitline coupling → electrical coupling between adjacent bitlines
  - Bitline-wordline coupling → electrical coupling between each bitline and the activated wordline



# Data Pattern Dependence

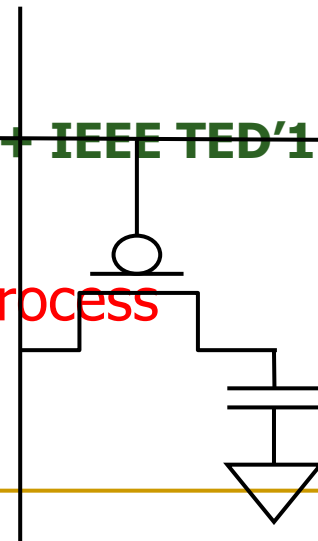
---

- Electrical noise on the bitline affects reliable sensing of a DRAM cell
- The magnitude of this noise is affected by values of nearby cells via
  - Bitline-bitline coupling → electrical coupling between adjacent bitlines
  - Bitline-wordline coupling → electrical coupling between each bitline and the activated wordline
  
- Retention time of a cell depends on data patterns stored in nearby cells
  - need to find the worst data pattern to find worst-case retention time
  - this pattern is location dependent

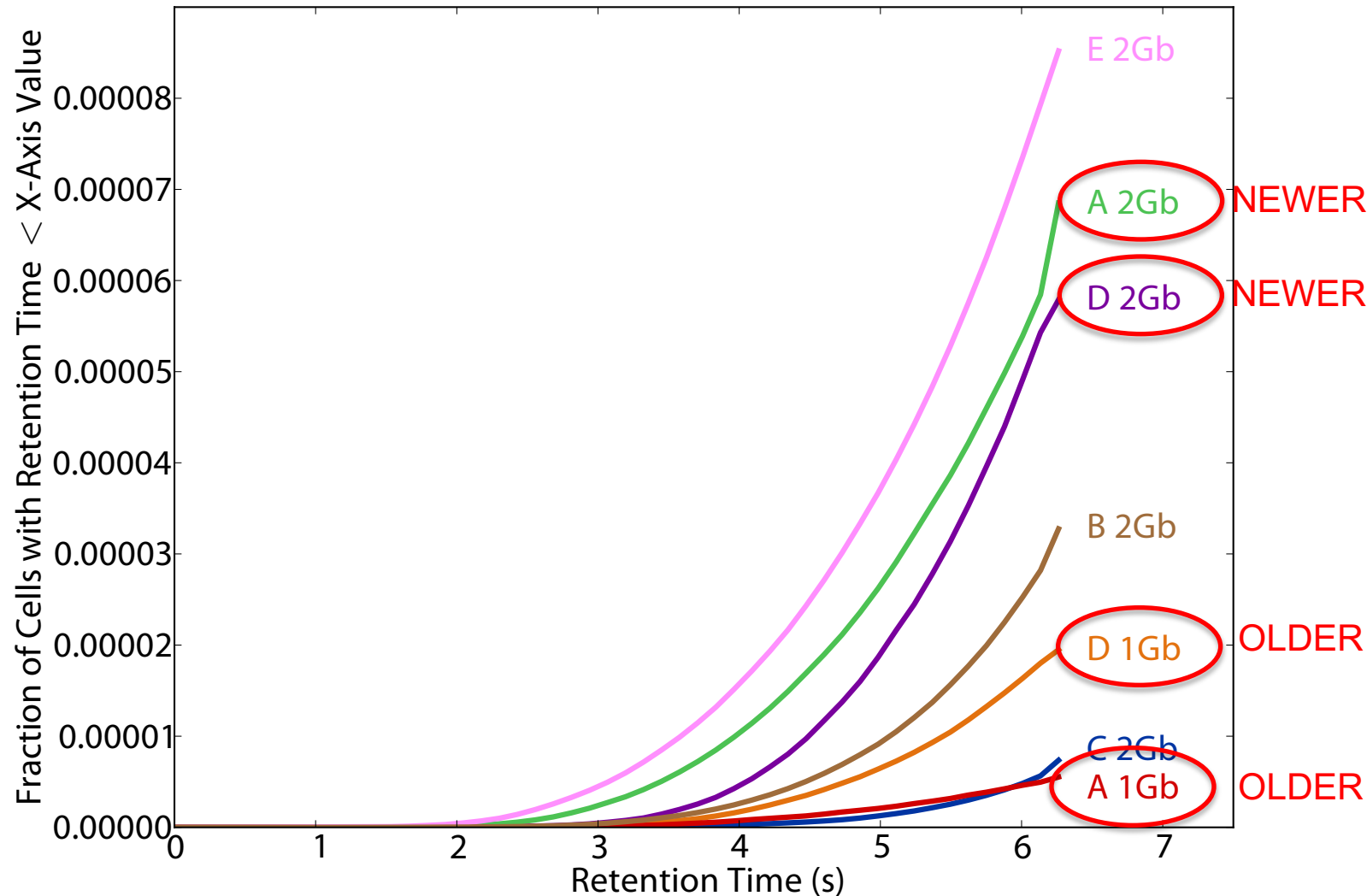
# Two Challenges to Retention Time Profiling

## ■ Challenge 2: Variable Retention Time (VRT)

- ❑ Retention time of a DRAM cell changes randomly over time
  - a cell alternates between multiple retention time states
- ❑ Leakage current of a cell changes sporadically due to a charge trap in the gate oxide of the DRAM cell access transistor
- ❑ When the trap becomes occupied, charge leaks more readily from the transistor's drain, leading to a short retention time
  - Called *Trap-Assisted Gate-Induced Drain Leakage*
- ❑ This process appears to be a random process [Kim+ IEEE TED'11]
- ❑ Worst-case retention time depends on a random process  
→ need to find the worst case despite this



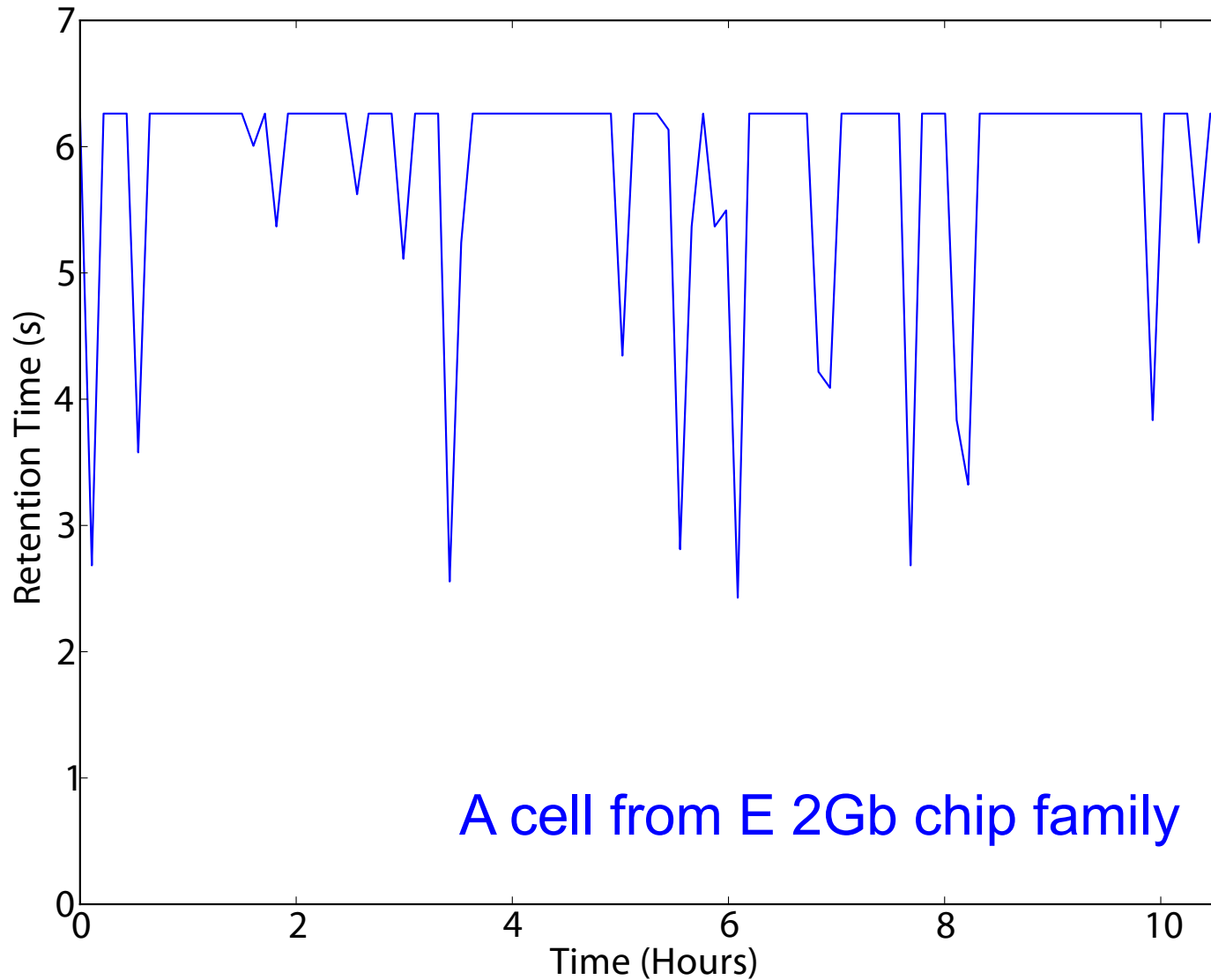
# Modern DRAM Retention Time Distribution



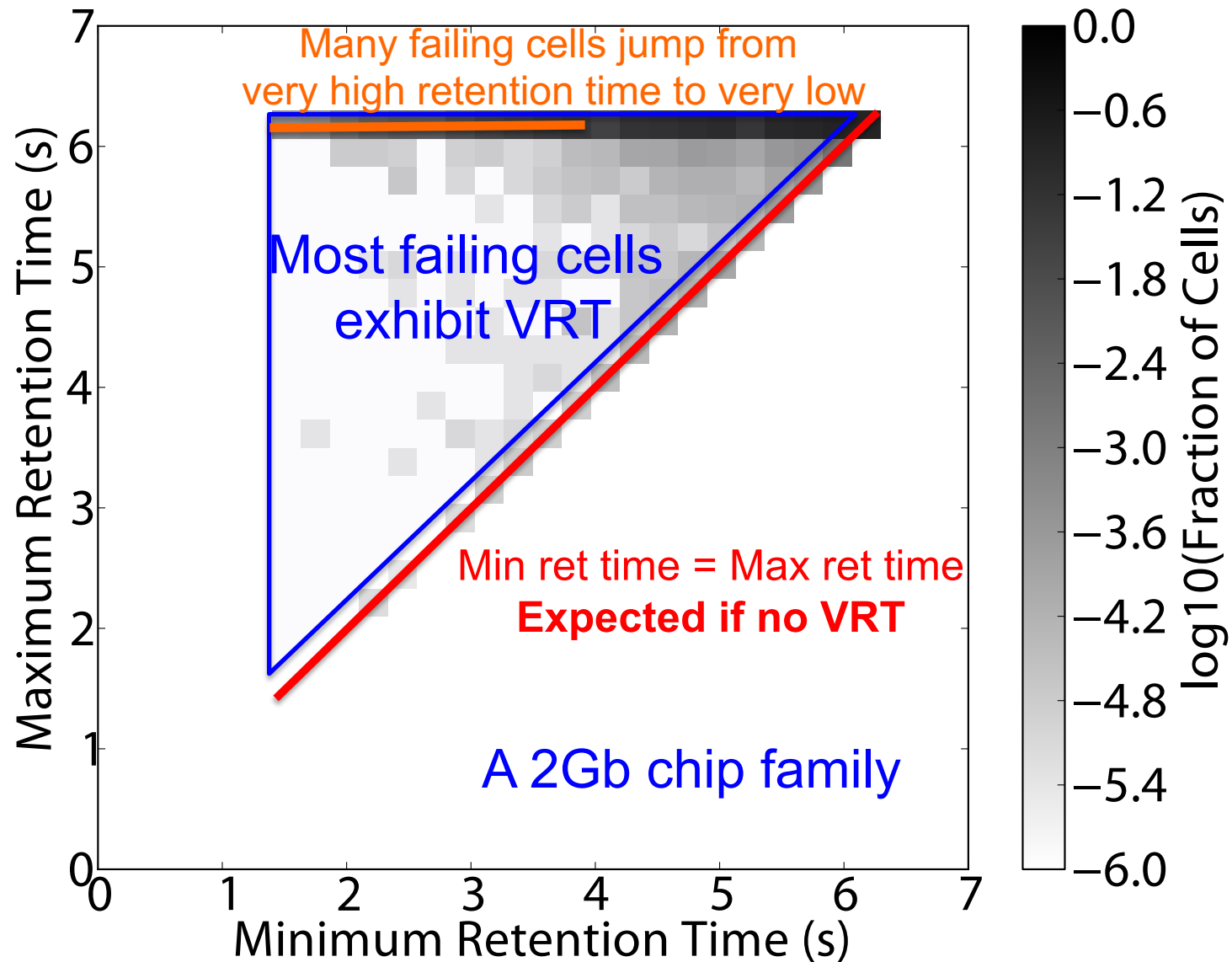
**Newer device families have more weak cells than older ones**  
**Likely a result of technology scaling**



# An Example VRT Cell



# Variable Retention Time



# Industry Is Writing Papers About It, Too

## DRAM Process Scaling Challenges

### ❖ Refresh

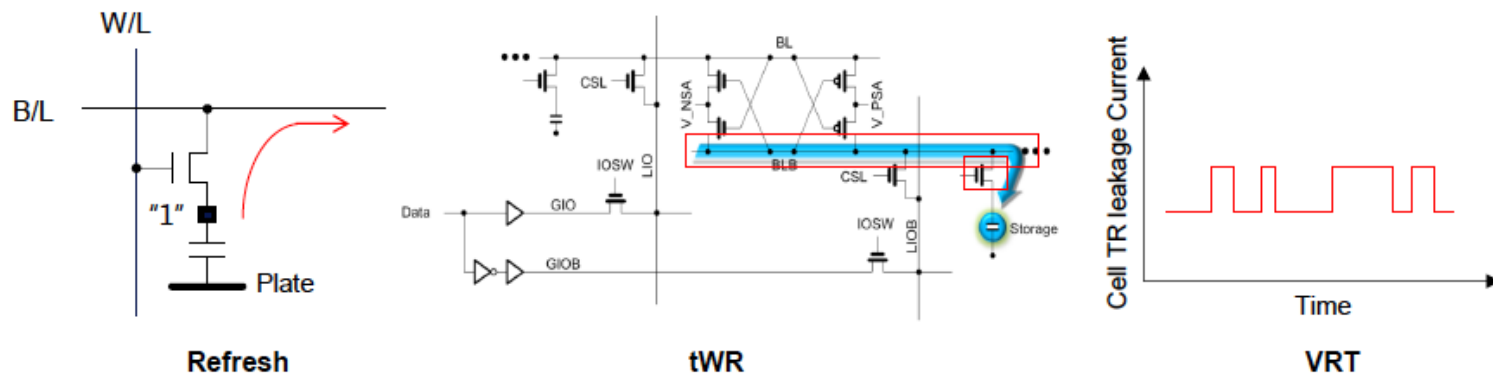
- Difficult to build high-aspect ratio cell capacitors decreasing cell capacitance
- Leakage current of cell access transistors increasing

### ❖ tWR

- Contact resistance between the cell capacitor and access transistor increasing
- On-current of the cell access transistor decreasing
- Bit-line resistance increasing

### ❖ VRT

- Occurring more frequently with cell capacitance decreasing



# Industry Is Writing Papers About It, Too

## DRAM Process Scaling Challenges

### ❖ Refresh

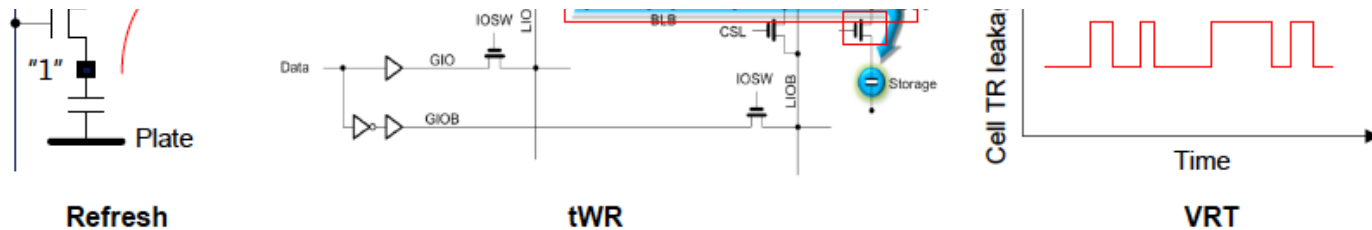
- Difficult to build high-aspect ratio cell capacitors decreasing cell capacitance

THE MEMORY FORUM 2014

## Co-Architecting Controllers and DRAM to Enhance DRAM Process Scaling

Uksong Kang, Hak-soo Yu, Churoo Park, \*Hongzhong Zheng,  
\*\*John Halbert, \*\*Kuljit Bains, SeongJin Jang, and Joo Sun Choi

*Samsung Electronics, Hwasung, Korea / \*Samsung Electronics, San Jose / \*\*Intel*



# Keeping Future Memory Secure

# How Do We Keep Memory Secure?

---

- DRAM
- Flash memory
- Emerging Technologies
  - Phase Change Memory
  - STT-MRAM
  - RRAM, memristors
  - ...

## Fundamentally Secure, Reliable, Safe Computing Architectures

# Solution Direction: Principled Designs

---

Design fundamentally secure  
computing architectures

Predict and prevent  
such safety issues

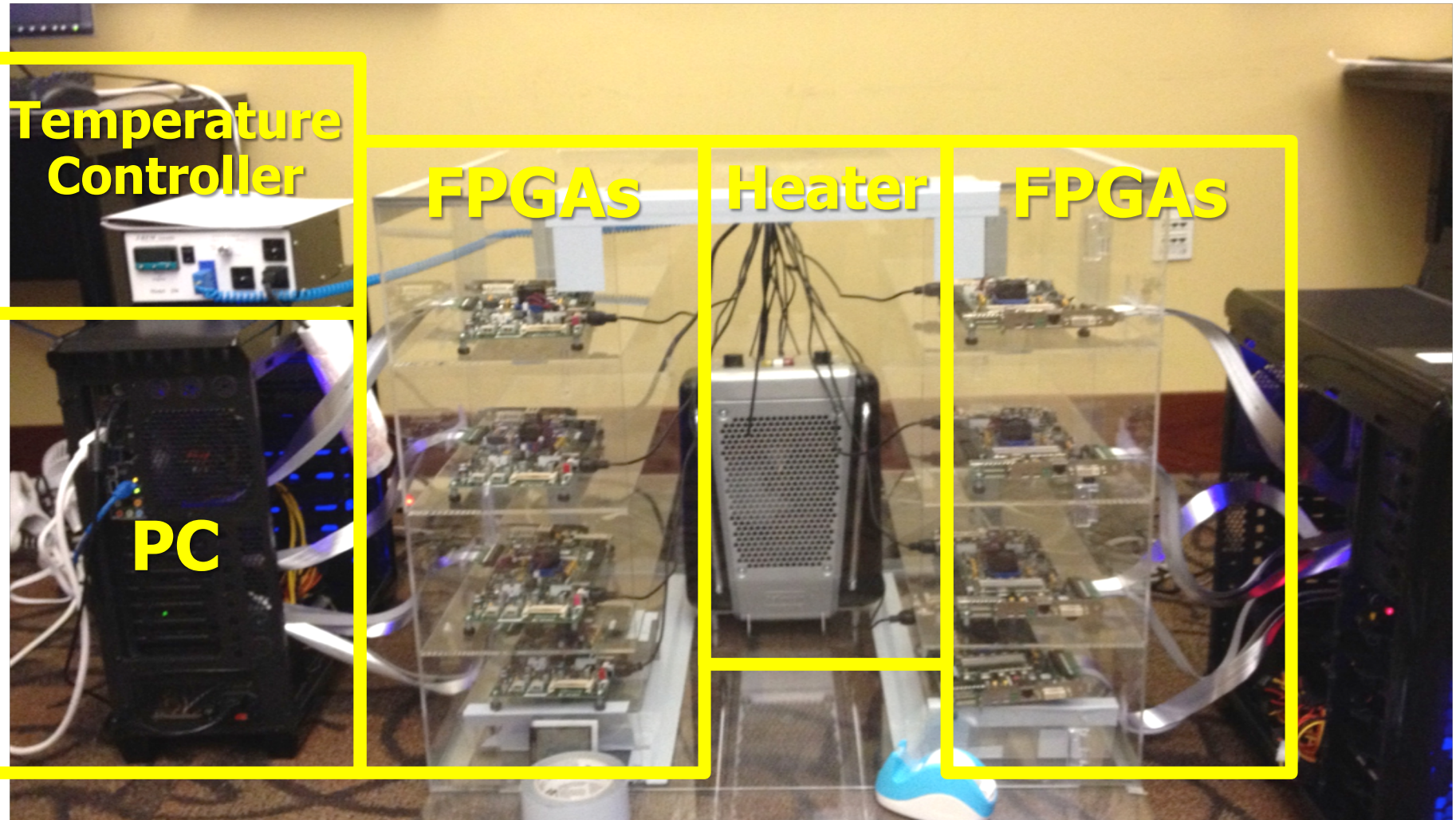


# Architecting for Security

---

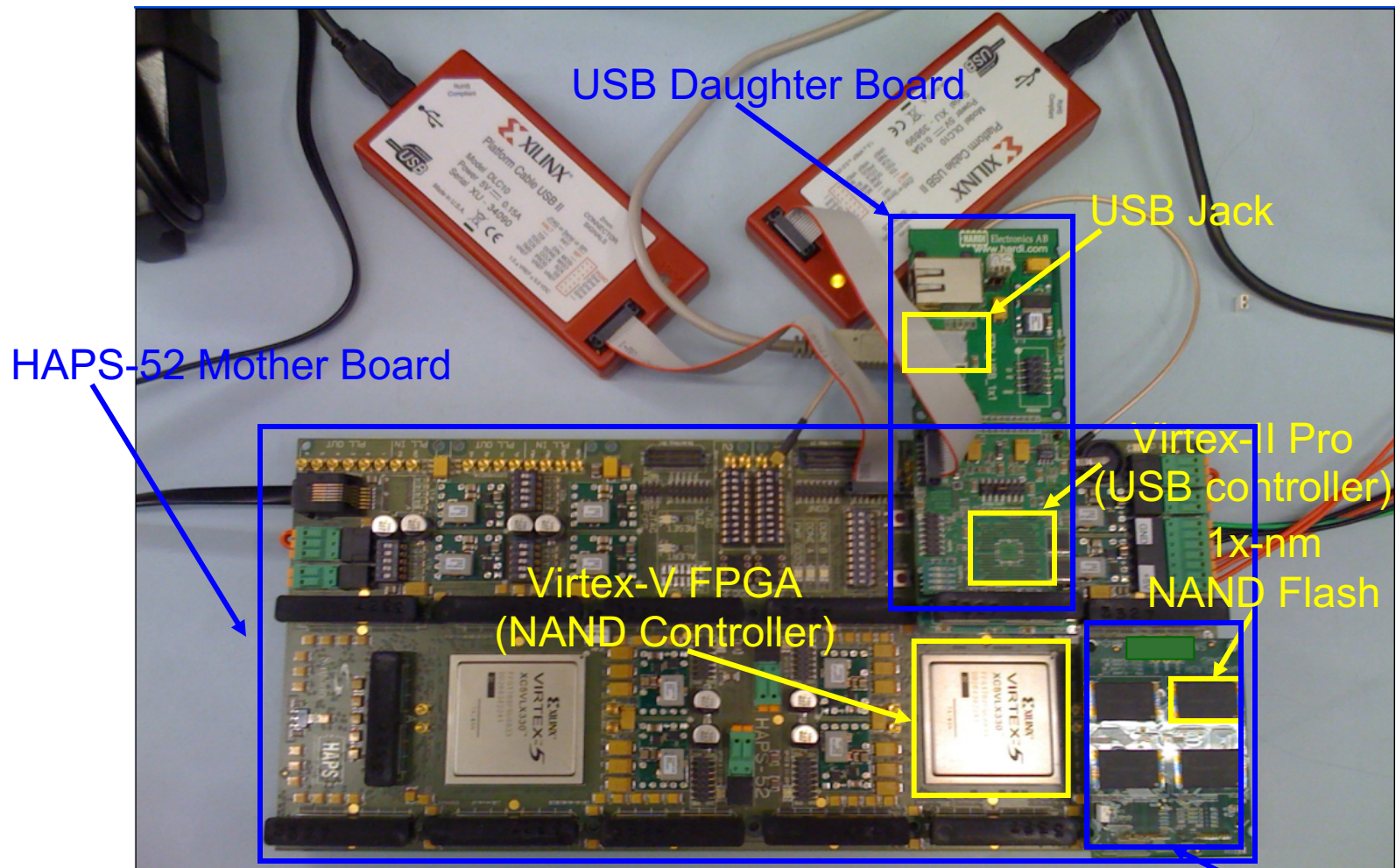
- **Understand:** Methods for vulnerability modeling & discovery
  - ❑ Modeling and prediction based on real (device) data and analysis
  - ❑ Understanding vulnerabilities
  - ❑ Developing reliable metrics
- **Architect:** Principled architectures with security as key concern
  - ❑ Good partitioning of duties across the stack
  - ❑ Cannot give up performance and efficiency
  - ❑ Patch-ability in the field
- **Design & Test:** Principled design, automation, (online) testing
  - ❑ Design for security
  - ❑ High coverage and good interaction with system reliability methods

# Understand and Model with Experiments (DRAM)





# Understand and Model with Experiments (Flash)



[DATE 2012, ICCD 2012, DATE 2013, ITJ 2013, ICCD 2013, SIGMETRICS 2014, HPCA 2015, DSN 2015, MSST 2015, JSAC 2016, HPCA 2017, DFRWS 2017, PIEEE 2017, HPCA 2018, SIGMETRICS 2018]

Cai+, "Error Characterization, Mitigation, and Recovery in Flash Memory Based Solid State Drives," Proc. IEEE 2017.



*Proceedings of the IEEE, Sept. 2017*

## Error Characterization, Mitigation, and Recovery in Flash-Memory-Based Solid-State Drives

*This paper reviews the most recent advances in solid-state drive (SSD) error characterization, mitigation, and data recovery techniques to improve both SSD's reliability and lifetime.*

By YU CAI, SAUGATA GHOSE, ERICH F. HARATSCH, YIXIN LUO, AND ONUR MUTLU

# Understanding Flash Memory Reliability

---

- Justin Meza, Qiang Wu, Sanjeev Kumar, and Onur Mutlu,  
**"A Large-Scale Study of Flash Memory Errors in the Field"**  
*Proceedings of the ACM International Conference on Measurement and Modeling of Computer Systems (**SIGMETRICS**), Portland, OR, June 2015.*  
[[Slides \(pptx\)](#)] [[pdf](#)] [[Coverage at ZDNet](#)] [[Coverage on The Register](#)]  
[[Coverage on TechSpot](#)] [[Coverage on The Tech Report](#)]

## A Large-Scale Study of Flash Memory Failures in the Field

Justin Meza  
Carnegie Mellon University  
[meza@cmu.edu](mailto:meza@cmu.edu)

Qiang Wu  
Facebook, Inc.  
[qw@fb.com](mailto:qw@fb.com)

Sanjeev Kumar  
Facebook, Inc.  
[skumar@fb.com](mailto:skumar@fb.com)

Onur Mutlu  
Carnegie Mellon University  
[onur@cmu.edu](mailto:onur@cmu.edu)



# NAND Flash Vulnerabilities [HPCA'17]

*HPCA, Feb. 2017*

## Vulnerabilities in MLC NAND Flash Memory Programming: Experimental Analysis, Exploits, and Mitigation Techniques

Yu Cai<sup>†</sup>   Saugata Ghose<sup>†</sup>   Yixin Luo<sup>††</sup>   Ken Mai<sup>†</sup>   Onur Mutlu<sup>§†</sup>   Erich F. Haratsch<sup>‡</sup>  
<sup>†</sup>Carnegie Mellon University   <sup>‡</sup>Seagate Technology   <sup>§</sup>ETH Zürich

*Modern NAND flash memory chips provide high density by storing two bits of data in each flash cell, called a multi-level cell (MLC). An MLC partitions the threshold voltage range of a flash cell into four voltage states. When a flash cell is programmed, a high voltage is applied to the cell. Due to parasitic capacitance coupling between flash cells that are physically close to each other, flash cell programming can lead to cell-to-cell program interference, which introduces errors into neighboring flash cells. In order to reduce the impact of cell-to-cell interference on the reliability of MLC NAND flash memory, flash manufacturers adopt a two-step programming method, which programs the MLC in two separate steps. First, the flash memory partially programs the least significant bit of the MLC to some intermediate threshold voltage. Second, it programs the most significant bit to bring the MLC up to its full voltage state.*

*In this paper, we demonstrate that two-step programming exposes new reliability and security vulnerabilities. We expe-*

*belongs to a different flash memory page (the unit of data programmed and read at the same time), which we refer to, respectively, as the least significant bit (LSB) page and the most significant bit (MSB) page [5].*

*A flash cell is programmed by applying a large voltage on the control gate of the transistor, which triggers charge transfer into the floating gate, thereby increasing the threshold voltage. To precisely control the threshold voltage of the cell, the flash memory uses incremental step pulse programming (ISPP) [12, 21, 25, 41]. ISPP applies multiple short pulses of the programming voltage to the control gate, in order to increase the cell threshold voltage by some small voltage amount ( $V_{step}$ ) after each step. Initial MLC designs programmed the threshold voltage in one shot, issuing all of the pulses back-to-back to program both bits of data at the same time. However, as flash memory scales down, the distance between neighboring flash cells decreases, which*

[https://people.inf.ethz.ch/omutlu/pub/flash-memory-programming-vulnerabilities\\_hpca17.pdf](https://people.inf.ethz.ch/omutlu/pub/flash-memory-programming-vulnerabilities_hpca17.pdf)

# 3D NAND Flash Reliability I [HPCA'18]

---

- Yixin Luo, Saugata Ghose, Yu Cai, Erich F. Haratsch, and Onur Mutlu, **"HeatWatch: Improving 3D NAND Flash Memory Device Reliability by Exploiting Self-Recovery and Temperature-Awareness"**

*Proceedings of the 24th International Symposium on High-Performance Computer Architecture (HPCA)*, Vienna, Austria, February 2018.

[[Lightning Talk Video](#)]

[[Slides \(pptx\)](#) ([pdf](#))] [[Lightning Session Slides \(pptx\)](#) ([pdf](#))]

## HeatWatch: Improving 3D NAND Flash Memory Device Reliability by Exploiting Self-Recovery and Temperature Awareness

Yixin Luo<sup>†</sup>      Saugata Ghose<sup>†</sup>      Yu Cai<sup>‡</sup>      Erich F. Haratsch<sup>‡</sup>      Onur Mutlu<sup>§†</sup>  
<sup>†</sup>*Carnegie Mellon University*      <sup>‡</sup>*Seagate Technology*      <sup>§</sup>*ETH Zürich*

# 3D NAND Flash Reliability II [SIGMETRICS'18]

---

- Yixin Luo, Saugata Ghose, Yu Cai, Erich F. Haratsch, and Onur Mutlu,  
**"Improving 3D NAND Flash Memory Lifetime by Tolerating Early Retention Loss and Process Variation"**  
*Proceedings of the ACM International Conference on Measurement and Modeling of Computer Systems (**SIGMETRICS**), Irvine, CA, USA, June 2018.*  
[[Abstract](#)]  
[[POMACS Journal Version \(same content, different format\)](#)]  
[[Slides \(pptx\)](#) ([pdf](#))]

## Improving 3D NAND Flash Memory Lifetime by Tolerating Early Retention Loss and Process Variation

Yixin Luo<sup>†</sup>      Saugata Ghose<sup>†</sup>      Yu Cai<sup>†</sup>      Erich F. Haratsch<sup>‡</sup>      Onur Mutlu<sup>§†</sup>

<sup>†</sup>Carnegie Mellon University

<sup>‡</sup>Seagate Technology

<sup>§</sup>ETH Zürich



# Recall: Collapse of the “Galloping Gertie”

---



# Another Example (1994)





# Yet Another Example (2007)

---



Source: Morry Gash/AP,  
<https://www.npr.org/2017/08/01/540669701/10-years-after-bridge-collapse-america-is-still-crumbling?t=1535427165809>



# A More Recent Example (2018)

---



In-Field Patch-ability  
(Intelligent Memory)  
Can Avoid Such Failures

# Final Thoughts on RowHammer

# Some Thoughts on RowHammer

---

- A simple hardware failure mechanism can create a widespread system security vulnerability
- How to exploit and fix the vulnerability requires a strong understanding across the transformation layers
  - And, a strong understanding of tools available to you
- Fixing needs to happen for two types of chips
  - Existing chips (already in the field)
  - Future chips
- Mechanisms for fixing are different between the two types

# Aside: Byzantine Failures

---

- This class of failures is known as **Byzantine failures**
- Characterized by
  - **Undetected erroneous computation**
  - Opposite of “fail fast (with an error or no result)”
- “erroneous” can be “malicious” (intent is the only distinction)
- Very difficult to detect and confine Byzantine failures
- **Do all you can to avoid them**
- Lamport et al., “The Byzantine Generals Problem,” ACM TOPLAS 1982.



# Aside: Byzantine Generals Problem

---

## The Byzantine Generals Problem

LESLIE LAMPORT, ROBERT SHOSTAK, and MARSHALL PEASE  
SRI International

---

Reliable computer systems must handle malfunctioning components that give conflicting information to different parts of the system. This situation can be expressed abstractly in terms of a group of generals of the Byzantine army camped with their troops around an enemy city. Communicating only by messenger, the generals must agree upon a common battle plan. However, one or more of them may be traitors who will try to confuse the others. The problem is to find an algorithm to ensure that the loyal generals will reach agreement. It is shown that, using only oral messages, this problem is solvable if and only if more than two-thirds of the generals are loyal; so a single traitor can confound two loyal generals. With unforgeable written messages, the problem is solvable for any number of generals and possible traitors. Applications of the solutions to reliable computer systems are then discussed.

Categories and Subject Descriptors: C.2.4. [**Computer-Communication Networks**]: Distributed Systems—*network operating systems*; D.4.4 [**Operating Systems**]: Communications Management—*network communication*; D.4.5 [**Operating Systems**]: Reliability—*fault tolerance*

General Terms: Algorithms, Reliability

Additional Key Words and Phrases: Interactive consistency

# RowHammer, Revisited

- One can **predictably induce bit flips** in commodity DRAM chips
  - >80% of the tested DRAM chips are vulnerable
- First example of how a **simple hardware failure mechanism** can create a **widespread system security vulnerability**

**WIRED**

Forget Software—Now Hackers Are Exploiting Physics

BUSINESS	CULTURE	DESIGN	GEAR	SCIENCE
----------	---------	--------	------	---------

ANDY GREENBERG SECURITY 08.31.16 7:00 AM

SHARE



SHARE  
18276



TWEET

# FORGET SOFTWARE—NOW HACKERS ARE EXPLOITING PHYSICS

# RowHammer: Retrospective

---

- New mindset that has enabled a renewed interest in HW security attack research:
  - ❑ Real (memory) chips are vulnerable, in a simple and widespread manner  
→ this causes real security problems
  - ❑ Hardware reliability → security connection is now mainstream discourse
- Many new RowHammer attacks...
  - ❑ Tens of papers in top security venues
  - ❑ **More to come** as RowHammer is getting worse (DDR4 & beyond)
- Many new RowHammer solutions...
  - ❑ Apple security release; Memtest86 updated
  - ❑ Many solution proposals in top venues (latest in ISCA 2019)
  - ❑ Principled system-DRAM co-design (in original RowHammer paper)
  - ❑ **More to come...**

# Perhaps Most Importantly...

---

- RowHammer enabled a shift of mindset in mainstream security researchers
  - ❑ General-purpose hardware is fallible, in a widespread manner
  - ❑ Its problems are exploitable
- This mindset has enabled many systems security researchers to examine hardware in more depth
  - ❑ And understand HW's inner workings and vulnerabilities
- It is no coincidence that two of the groups that discovered Meltdown and Spectre heavily worked on RowHammer attacks before
  - ❑ **More to come...**

# Summary: RowHammer

---

- DRAM reliability is reducing
- Reliability issues open up security vulnerabilities
  - Very hard to defend against
- **Rowhammer is a prime example**
  - First example of how a simple hardware failure mechanism can create a widespread system security vulnerability
  - Its implications on system security research are tremendous & exciting
- Bad news: RowHammer is getting worse.
- **Good news: We have a lot more to do.**
  - We are now fully aware hardware is easily fallible.
  - We are developing both attacks and solutions.
  - We are developing principled models, methodologies, solutions.

# For More on RowHammer...

---

- Onur Mutlu and Jeremie Kim,  
**"RowHammer: A Retrospective"**  
*IEEE Transactions on Computer-Aided Design of Integrated Circuits and Systems (TCAD) Special Issue on Top Picks in Hardware and Embedded Security*, 2019.  
[[Preliminary arXiv version](#)]

## RowHammer: A Retrospective

Onur Mutlu<sup>§‡</sup>      Jeremie S. Kim<sup>‡§</sup>  
§ETH Zürich      ‡Carnegie Mellon University



# Computer Architecture

## Lecture 6a: RowHammer II

Prof. Onur Mutlu

ETH Zürich

Fall 2019

4 October 2019



# Future Memory Reliability/Security Challenges

# Finding DRAM Retention Failures

---

- How can we reliably find the retention time of all DRAM cells?
- Goals: so that we can
  - Make DRAM reliable and secure
  - Make techniques like RAIDR work
    - improve performance and energy

# Mitigation of Retention Issues [SIGMETRICS'14]

---

- Samira Khan, Donghyuk Lee, Yoongu Kim, Alaa Alameldeen, Chris Wilkerson, and Onur Mutlu,  
**"The Efficacy of Error Mitigation Techniques for DRAM Retention Failures: A Comparative Experimental Study"**  
*Proceedings of the ACM International Conference on Measurement and Modeling of Computer Systems (**SIGMETRICS**), Austin, TX, June 2014. [[Slides \(pptx\)](#)] [[pdf](#)] [[Poster \(pptx\)](#)] [[pdf](#)] [[Full data sets](#)]*

## The Efficacy of Error Mitigation Techniques for DRAM Retention Failures: A Comparative Experimental Study

Samira Khan<sup>†\*</sup>  
samirakhan@cmu.edu

Donghyuk Lee<sup>†</sup>  
donghyuk1@cmu.edu

Yoongu Kim<sup>†</sup>  
yoongukim@cmu.edu

Alaa R. Alameldeen<sup>\*</sup>  
alaa.r.alameldeen@intel.com

Chris Wilkerson<sup>\*</sup>  
chris.wilkerson@intel.com

Onur Mutlu<sup>†</sup>  
onur@cmu.edu

<sup>†</sup>Carnegie Mellon University

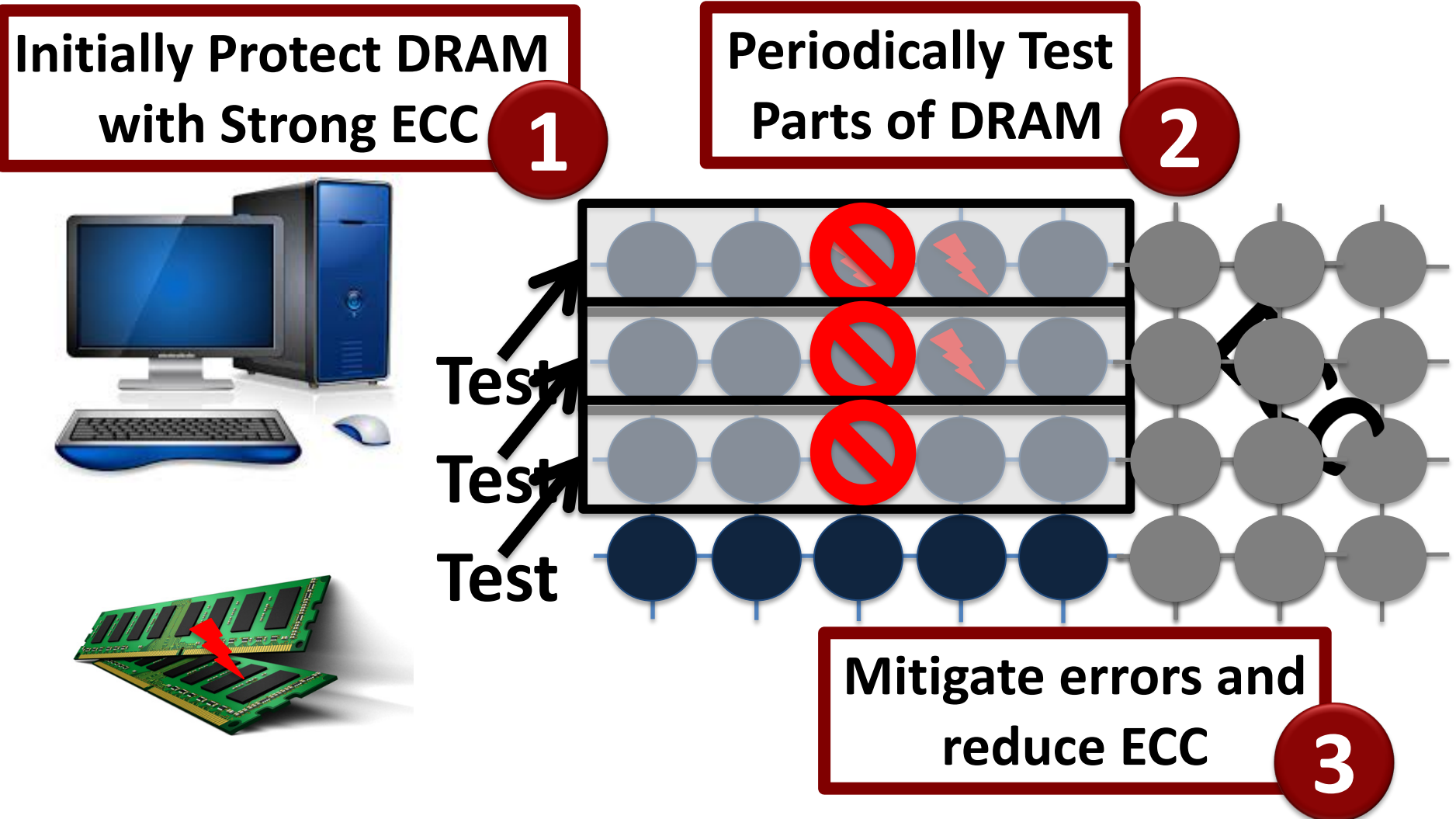
<sup>\*</sup>Intel Labs

# Towards an Online Profiling System

## *Key Observations:*

- **Testing** alone **cannot detect** all possible failures
- **Combination** of ECC and other mitigation techniques is much more **effective**
  - But degrades performance
- **Testing** can help to reduce the **ECC strength**
  - Even when starting with a **higher strength ECC**

# Towards an Online Profiling System



**Run tests periodically after a short interval  
at smaller regions of memory**

# Handling Variable Retention Time [DSN'15]

---

- Moinuddin Qureshi, Dae Hyun Kim, Samira Khan, Prashant Nair, and Onur Mutlu, **"AVATAR: A Variable-Retention-Time (VRT) Aware Refresh for DRAM Systems"**

*Proceedings of the 45th Annual IEEE/IFIP International Conference on Dependable Systems and Networks (DSN)*, Rio de Janeiro, Brazil, June 2015.  
[[Slides \(pptx\)](#)] [[pdf](#)]

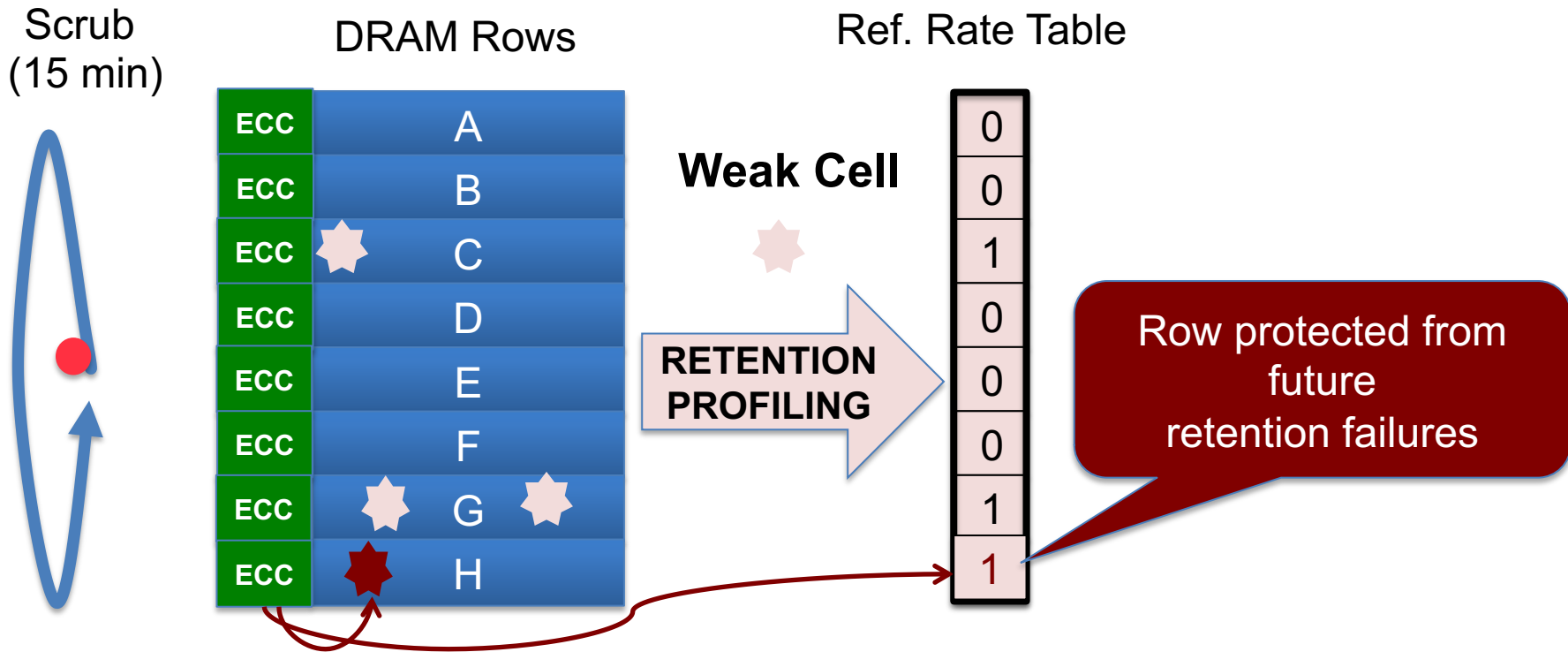
## AVATAR: A Variable-Retention-Time (VRT) Aware Refresh for DRAM Systems

Moinuddin K. Qureshi <sup>†</sup>	Dae-Hyun Kim <sup>†</sup>	Samira Khan <sup>‡</sup>	Prashant J. Nair <sup>†</sup>	Onur Mutlu <sup>‡</sup>
<sup>†</sup> Georgia Institute of Technology {moin, dhkim, pnair6}@ece.gatech.edu			<sup>‡</sup> Carnegie Mellon University {samirakhan, onur}@cmu.edu	

# AVATAR

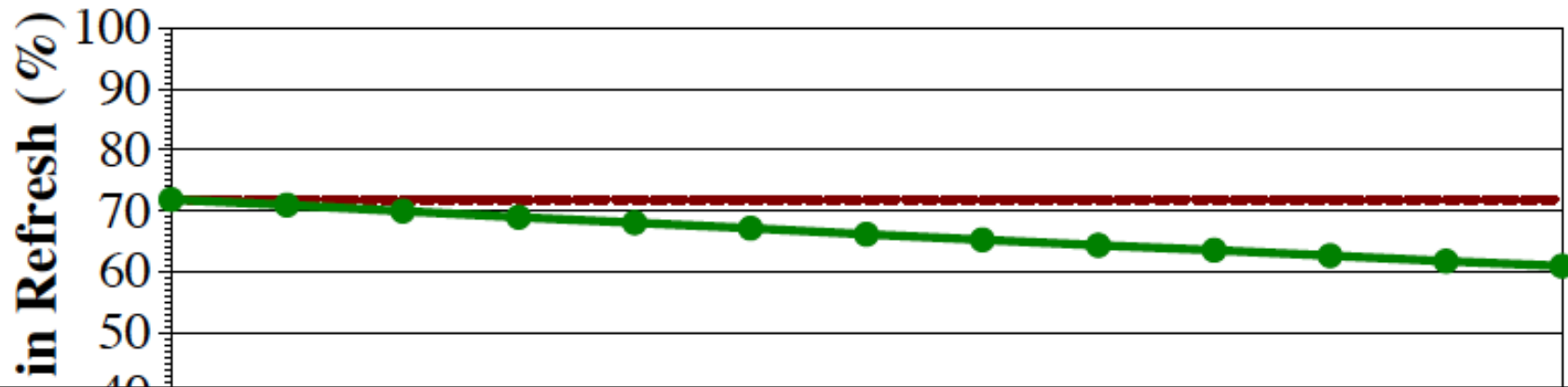
**Insight:** Avoid retention failures → Upgrade row on ECC error

**Observation:** Rate of VRT >> Rate of soft error (50x-2500x)

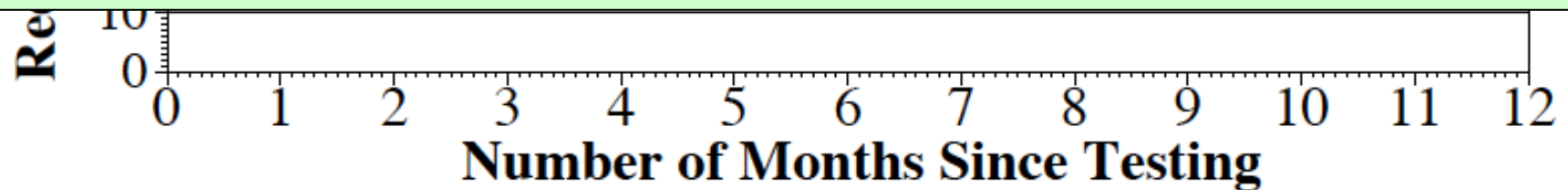


**AVATAR mitigates VRT by increasing refresh rate on error**

# RESULTS: REFRESH SAVINGS



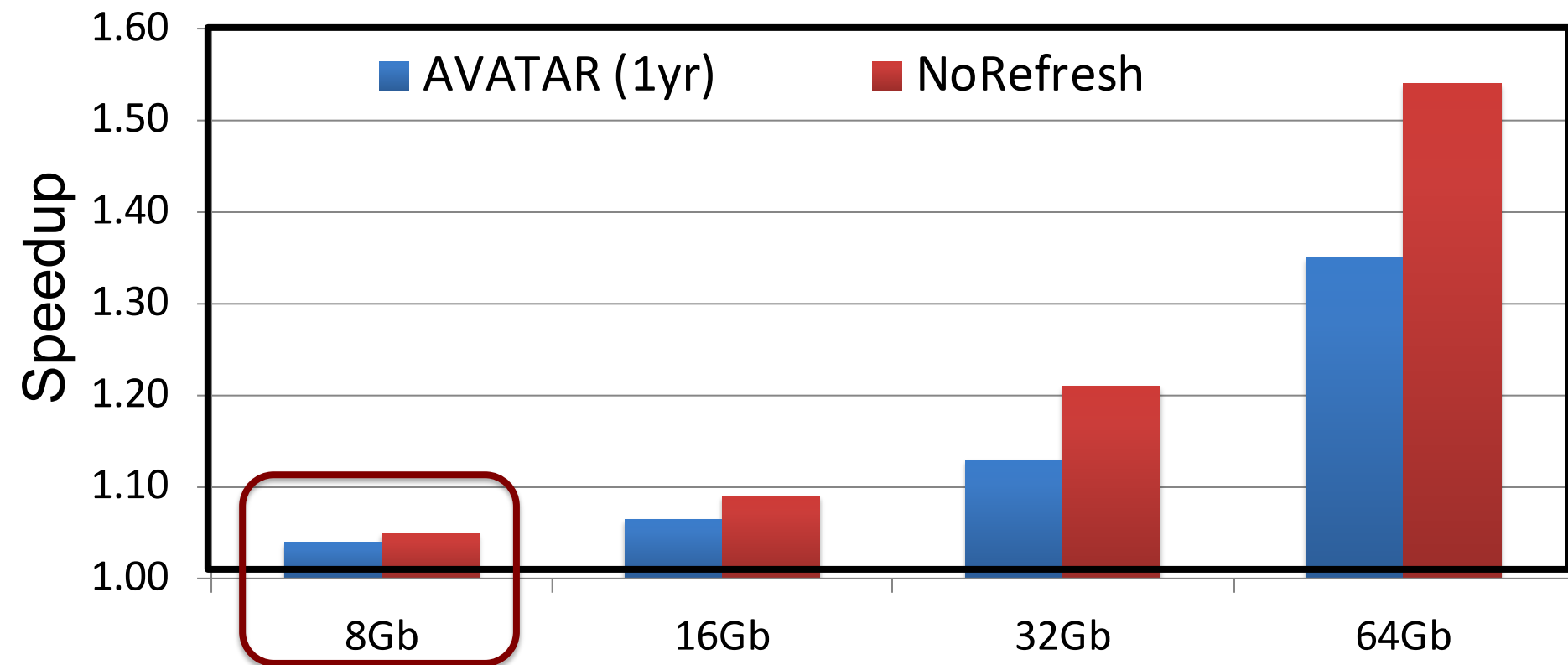
**Retention Testing Once a Year can revert refresh saving from 60% to 70%**



**AVATAR reduces refresh by 60%-70%, similar to multi rate refresh but with VRT tolerance**

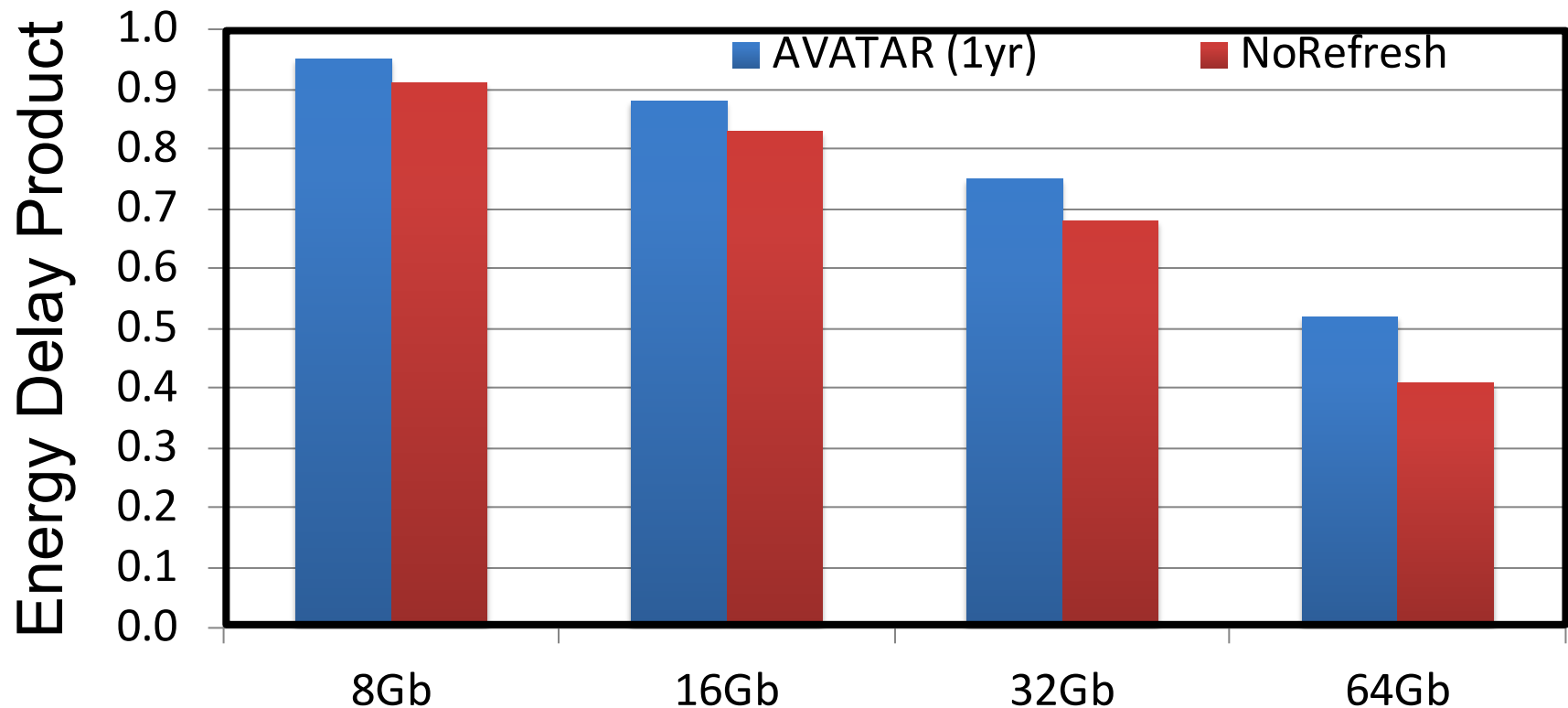


# SPEEDUP



**AVATAR gets 2/3<sup>rd</sup> the performance of NoRefresh. More gains at higher capacity nodes**

# ENERGY DELAY PRODUCT



**AVATAR reduces EDP,  
Significant reduction at higher capacity nodes**

# Handling Data-Dependent Failures [DSN'16]

---

- Samira Khan, Donghyuk Lee, and Onur Mutlu,  
**"PARBOR: An Efficient System-Level Technique to Detect Data-Dependent Failures in DRAM"**  
*Proceedings of the 45th Annual IEEE/IFIP International Conference on Dependable Systems and Networks (DSN)*, Toulouse, France, June 2016.  
[[Slides \(pptx\)](#)] [[pdf](#)]

## PARBOR: An Efficient System-Level Technique to Detect Data-Dependent Failures in DRAM

Samira Khan<sup>\*</sup>

<sup>\*</sup>University of Virginia

Donghyuk Lee<sup>†‡</sup>

<sup>†</sup>Carnegie Mellon University

Onur Mutlu<sup>\*†</sup>

<sup>‡</sup>Nvidia

<sup>\*</sup>ETH Zürich

# Handling Data-Dependent Failures [MICRO'17]

---

- Samira Khan, Chris Wilkerson, Zhe Wang, Alaa R. Alameldeen, Donghyuk Lee, and Onur Mutlu,  
**"Detecting and Mitigating Data-Dependent DRAM Failures by Exploiting Current Memory Content"**  
*Proceedings of the 50th International Symposium on Microarchitecture (MICRO)*, Boston, MA, USA, October 2017.  
[\[Slides \(pptx\) \(pdf\)\]](#) [\[Lightning Session Slides \(pptx\) \(pdf\)\]](#) [\[Poster \(pptx\) \(pdf\)\]](#)

## Detecting and Mitigating Data-Dependent DRAM Failures by Exploiting Current Memory Content

Samira Khan<sup>\*</sup> Chris Wilkerson<sup>†</sup> Zhe Wang<sup>†</sup> Alaa R. Alameldeen<sup>†</sup> Donghyuk Lee<sup>‡</sup> Onur Mutlu<sup>\*</sup>  
<sup>\*</sup>University of Virginia    <sup>†</sup>Intel Labs    <sup>‡</sup>Nvidia Research    <sup>\*</sup>ETH Zürich

# Handling Both DPD and VRT [ISCA'17]

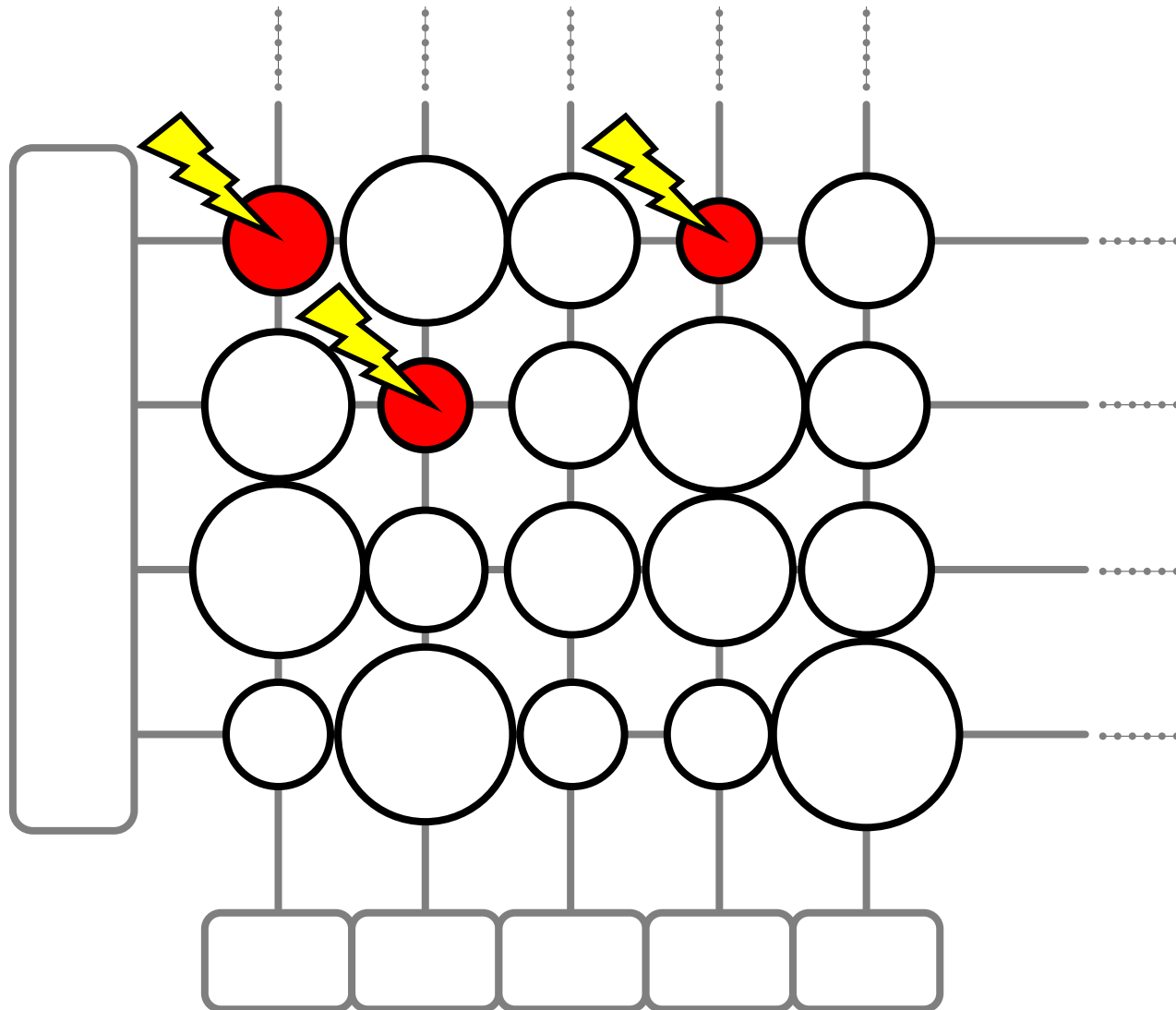
---

- Minesh Patel, Jeremie S. Kim, and Onur Mutlu,  
**"The Reach Profiler (REAPER): Enabling the Mitigation of DRAM Retention Failures via Profiling at Aggressive Conditions"**  
*Proceedings of the 44th International Symposium on Computer Architecture (ISCA)*, Toronto, Canada, June 2017.  
[[Slides \(pptx\)](#)] [[pdf](#)]  
[[Lightning Session Slides \(pptx\)](#)] [[pdf](#)]
- First experimental analysis of (mobile) LPDDR4 chips
- Analyzes the complex tradeoff space of retention time profiling
- Idea: enable fast and robust profiling at higher refresh intervals & temperatures

## The Reach Profiler (REAPER): Enabling the Mitigation of DRAM Retention Failures via Profiling at Aggressive Conditions

Minesh Patel<sup>§‡</sup>   Jeremie S. Kim<sup>‡§</sup>   Onur Mutlu<sup>§‡</sup>  
<sup>§</sup>ETH Zürich   <sup>‡</sup>Carnegie Mellon University

**Goal:** find *all* retention failures for a refresh interval  $T > \text{default (64ms)}$





**Process, voltage, temperature**

**Variable retention time**

**Data pattern dependence**

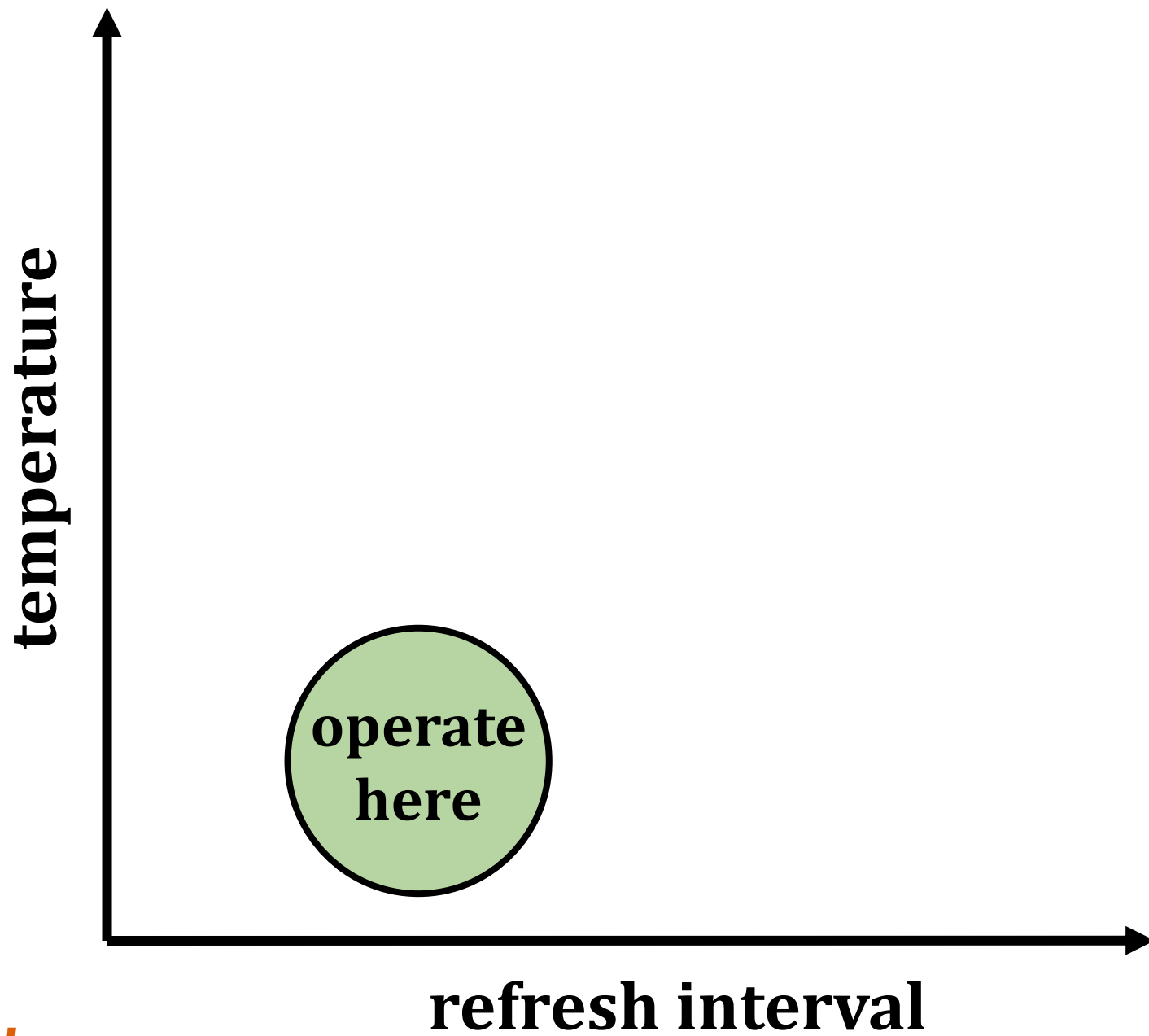
# Characterization of 368 LPDDR4 DRAM Chips

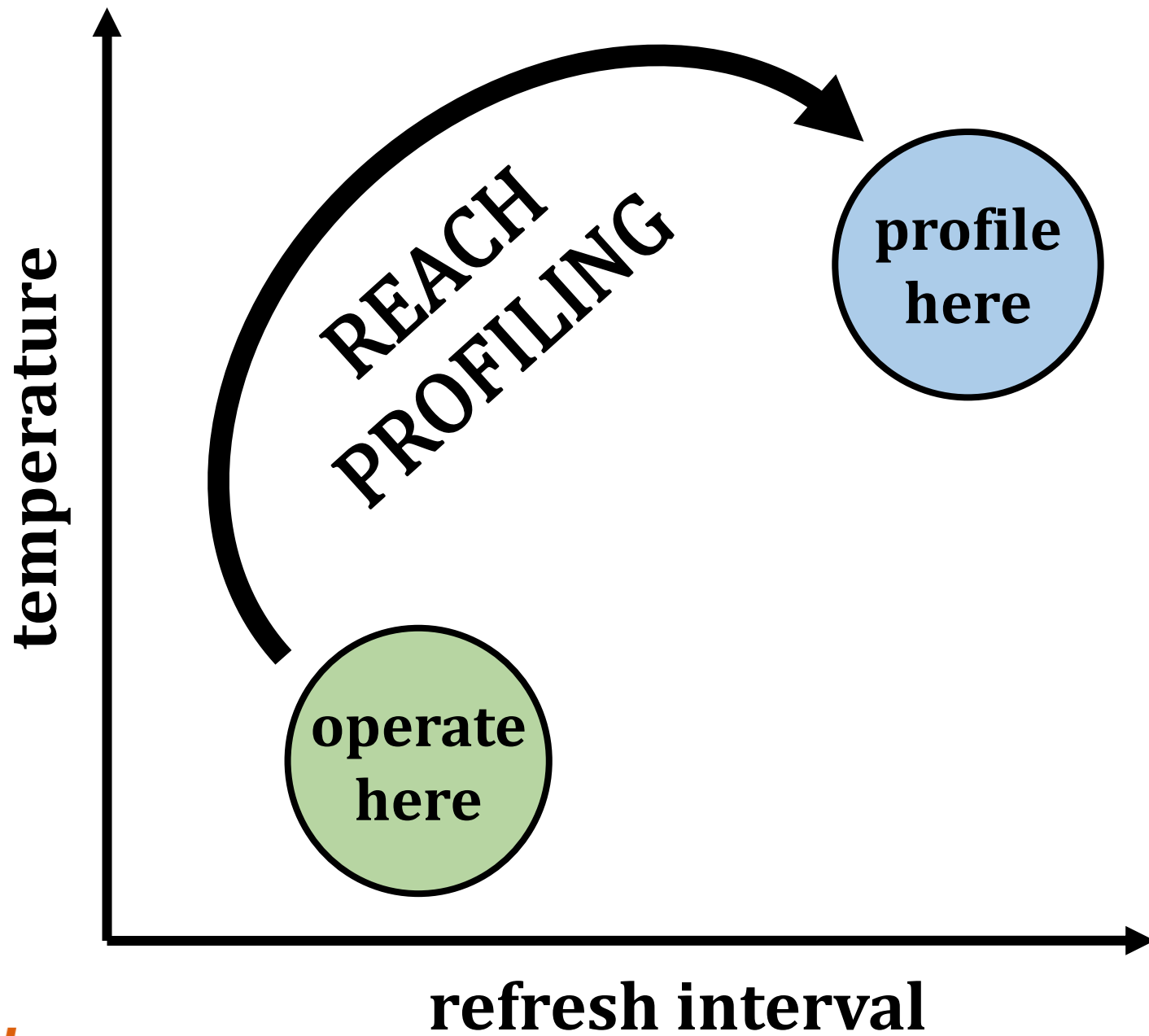
①

Cells are **more likely to fail** at an **increased (refresh interval | temperature)**

②

**Complex tradeoff space** between profiling  
**(speed & coverage & false positives)**





# Reach Profiling

**A new DRAM retention failure  
profiling methodology**

+ **Faster** and **more reliable**  
than current approaches

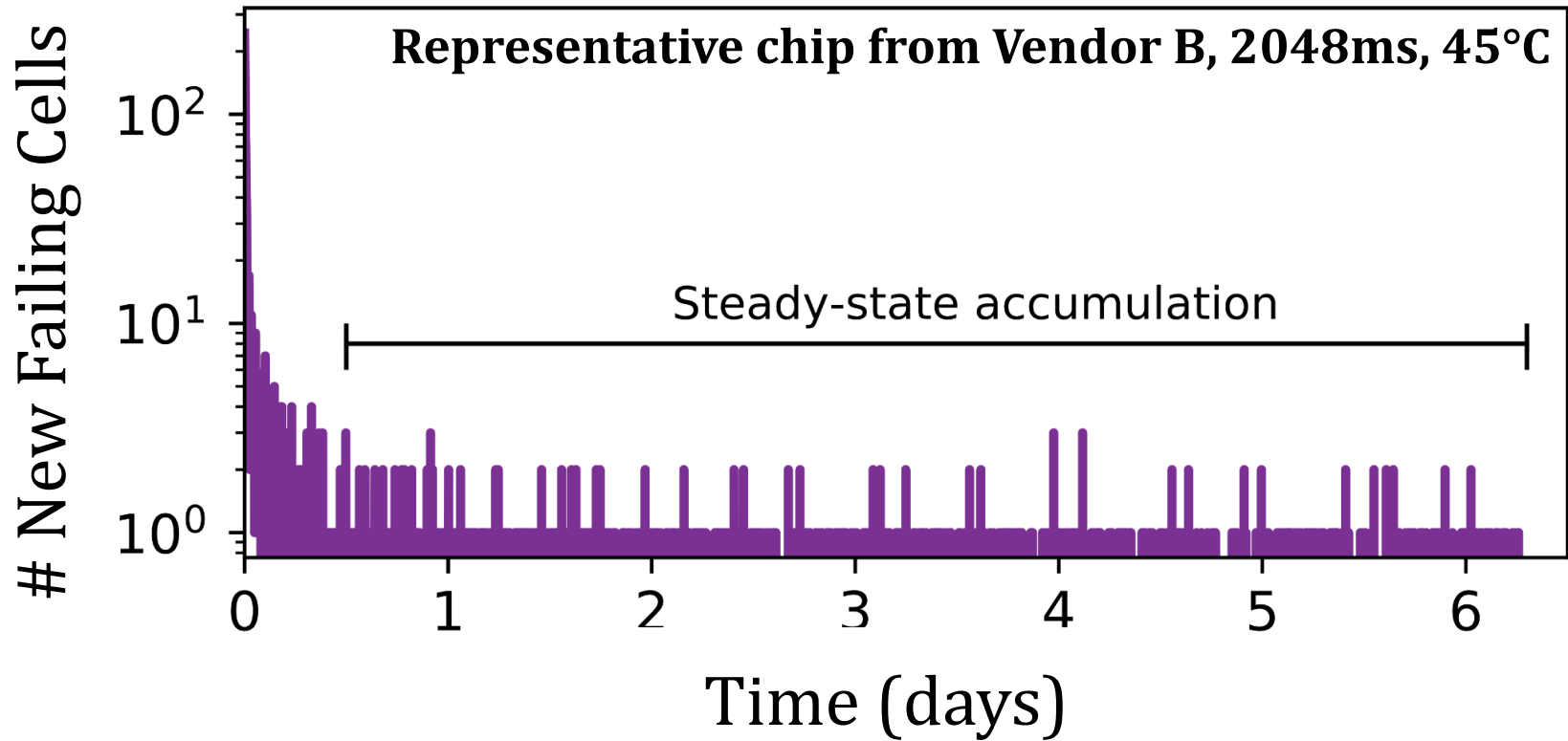
+ Enables **longer refresh intervals**

# LPDDR4 Studies

1. Temperature
2. Data Pattern Dependence
3. Retention Time Distributions
- 4. Variable Retention Time**
- 5. Individual Cell Characterization**



# Long-term Continuous Profiling



- New failing cells continue to appear over time
  - Attributed to **variable retention time (VRT)**
- The set of failing cells changes over time

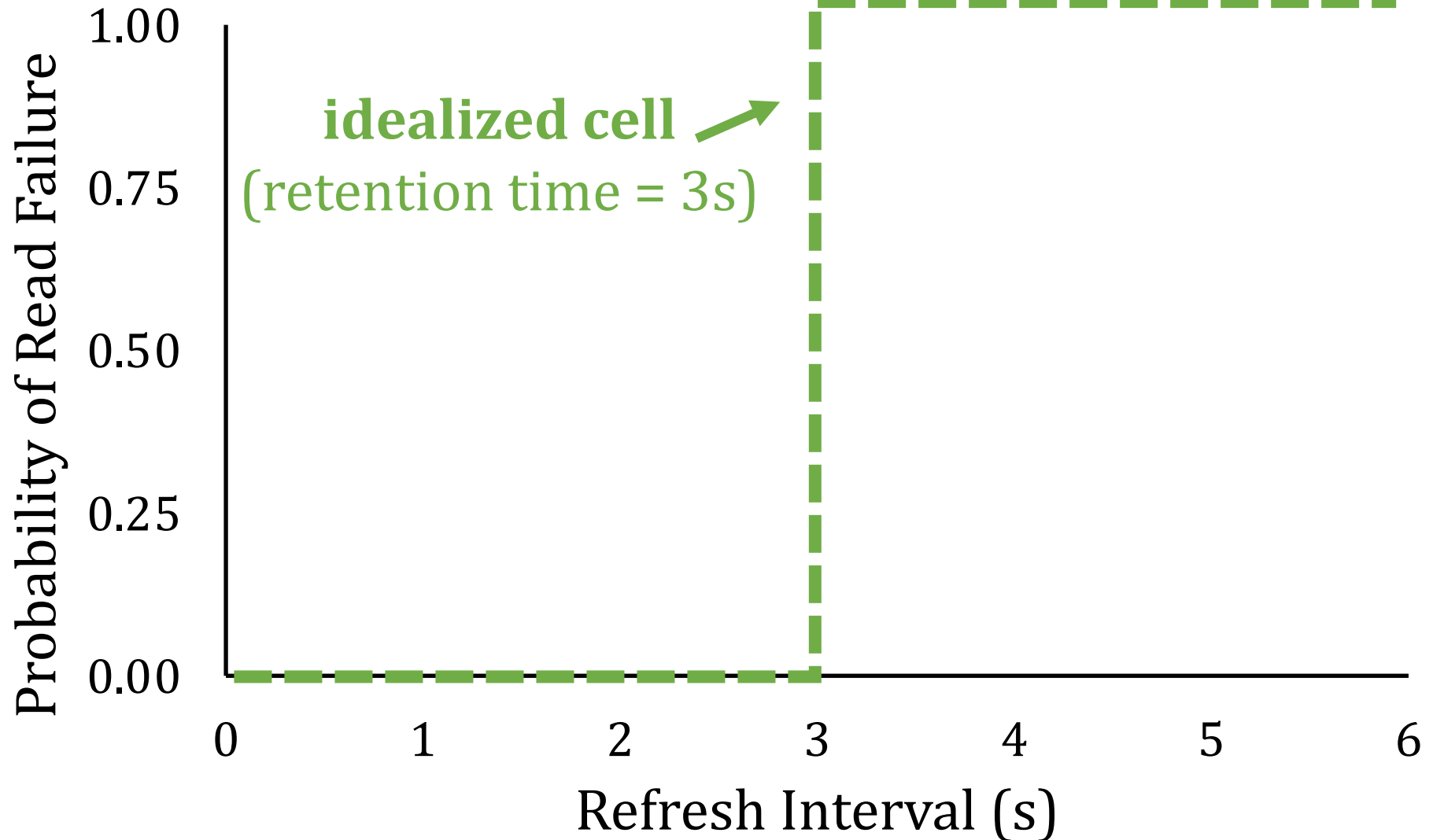
# Long-term Continuous Profiling



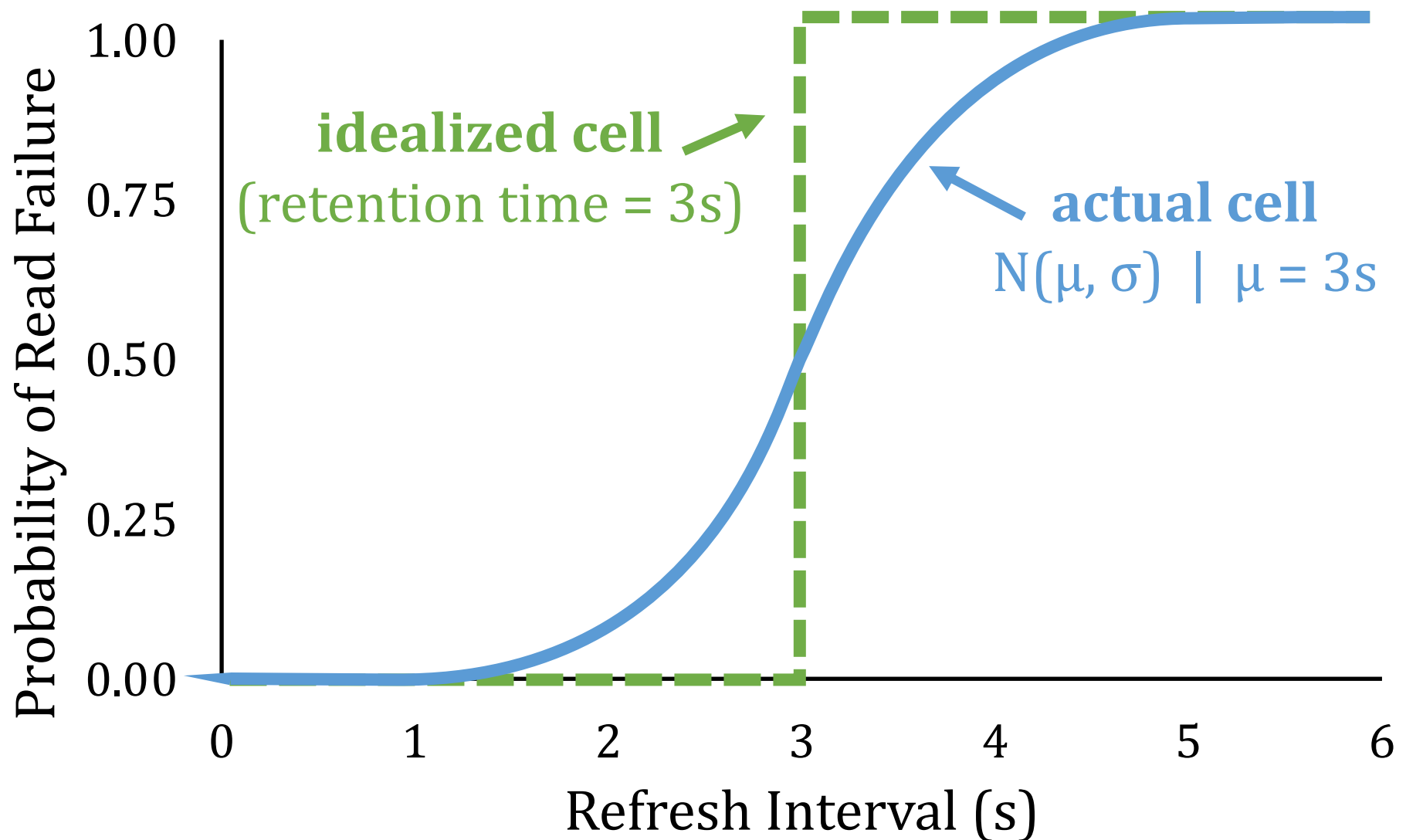
**Error correction codes (ECC)**  
**and online profiling are *necessary***  
**to manage new failing cells**

- New failing cells continue to appear over time
  - Attributed to **variable retention time (VRT)**
- The set of failing cells changes over time

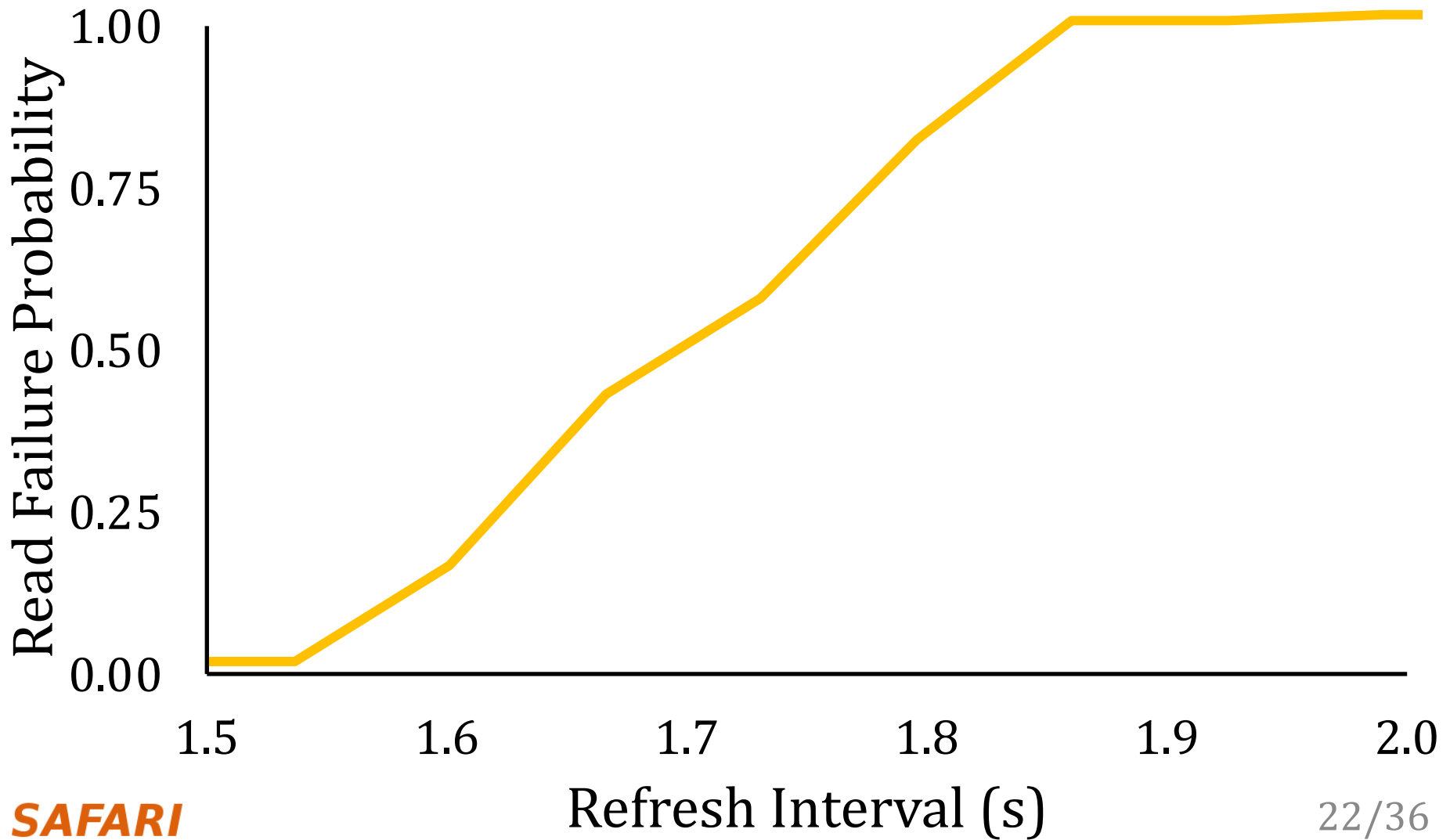
# Single-cell Failure Probability (Cartoon)



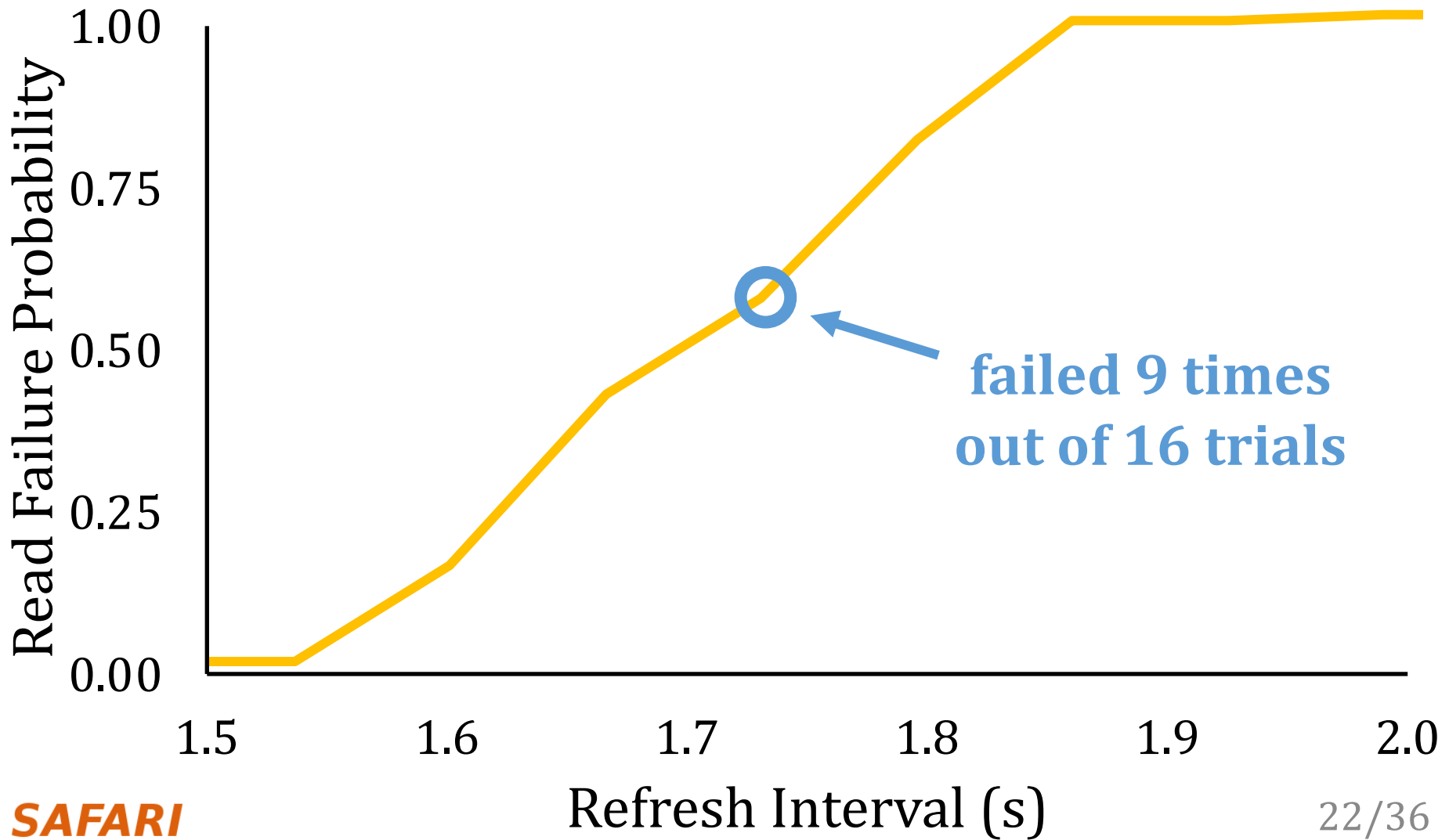
# Single-cell Failure Probability (Cartoon)



# Single-cell Failure Probability (Real)

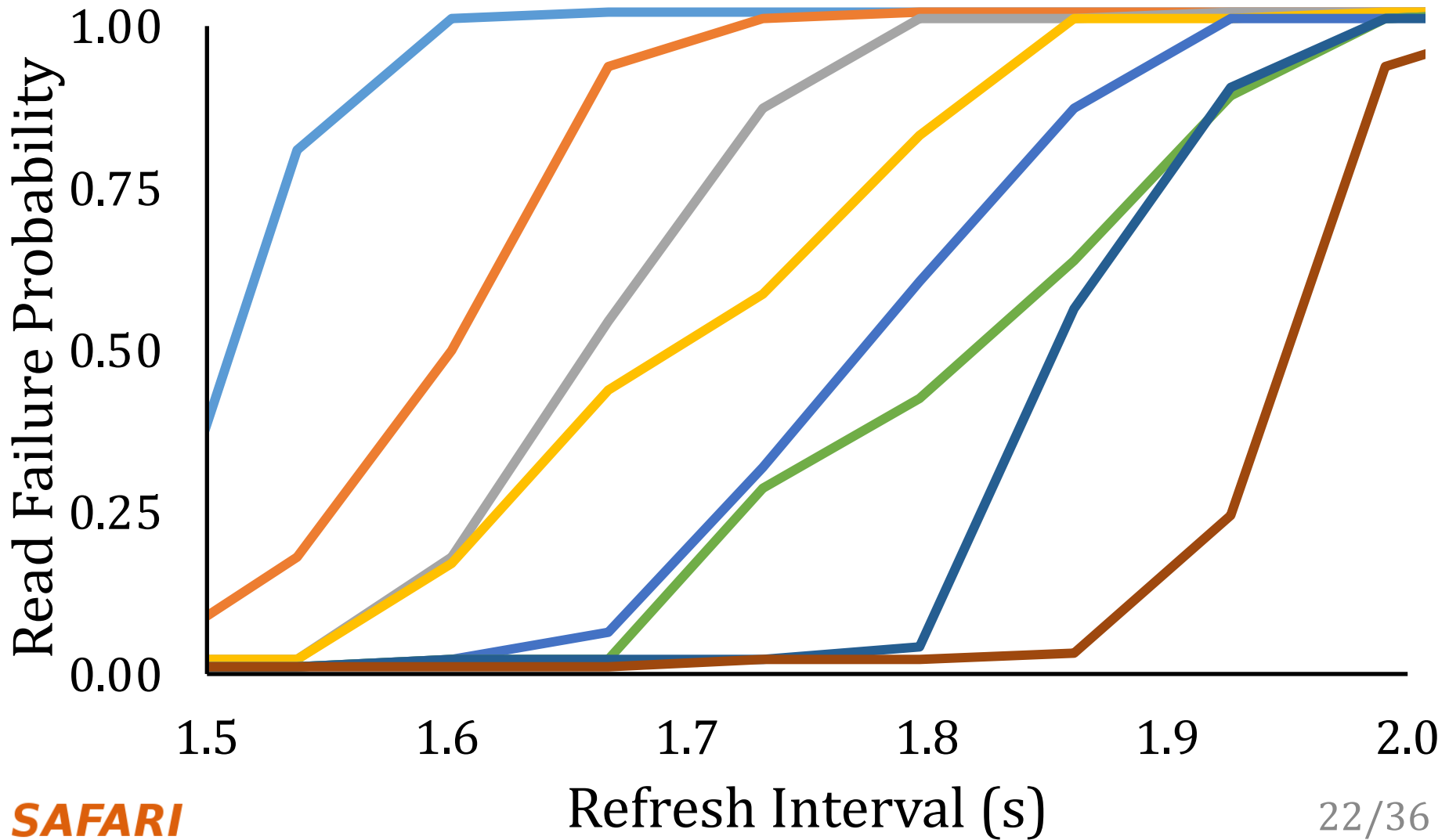


# Single-cell Failure Probability (Real)

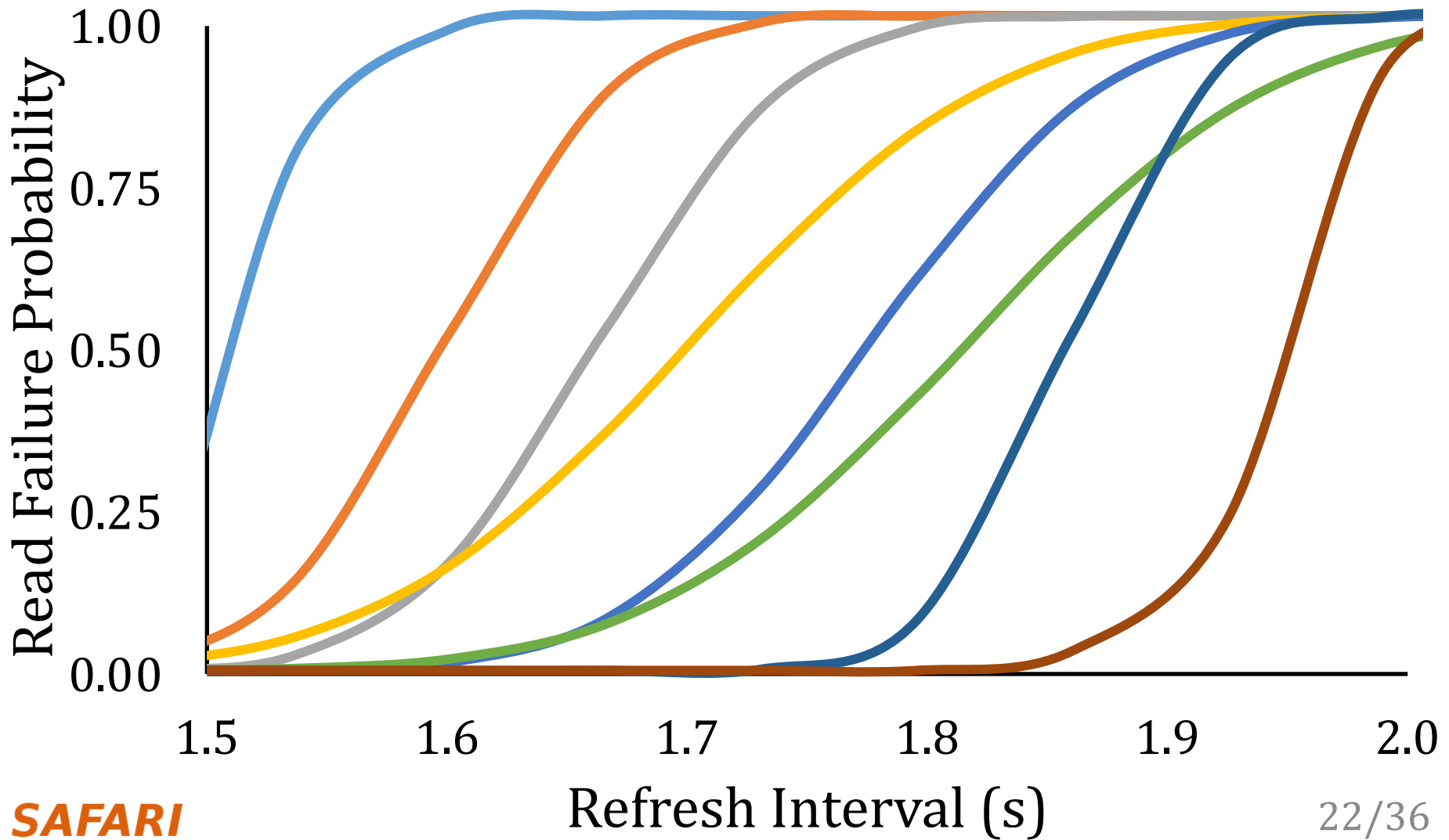




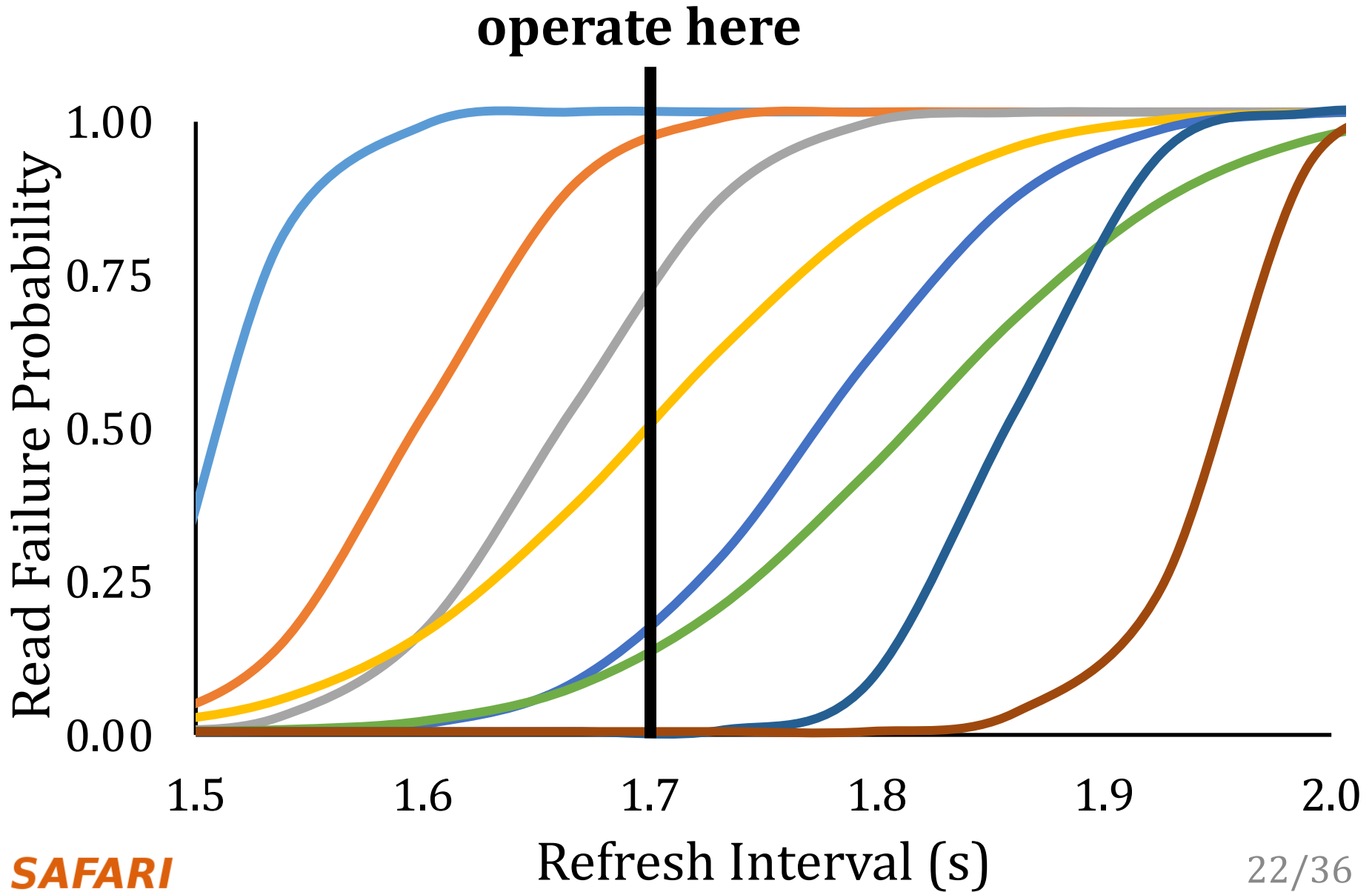
# Single-cell Failure Probability (Real)



# Single-cell Failure Probability (Real)

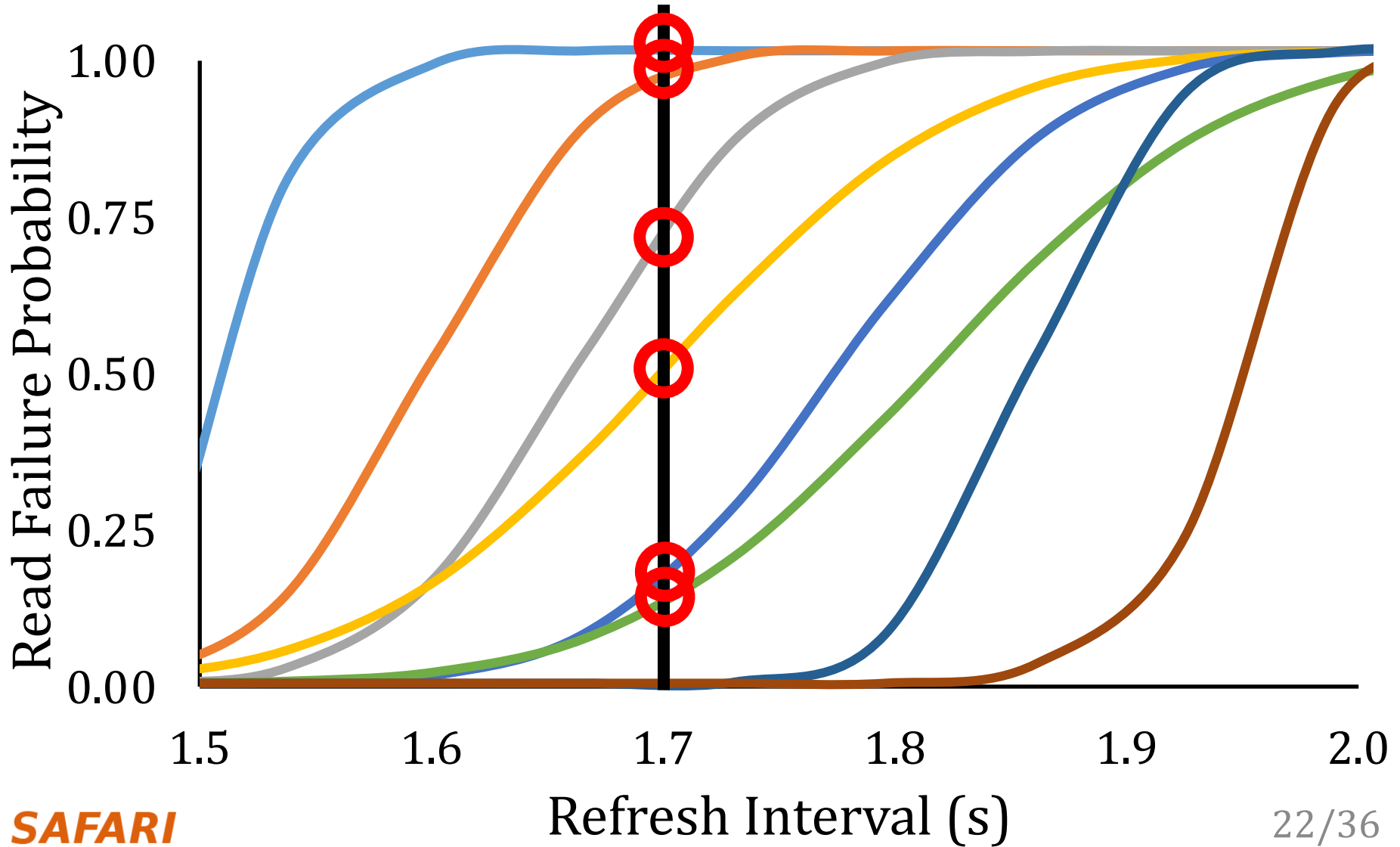


# Single-cell Failure Probability (Real)

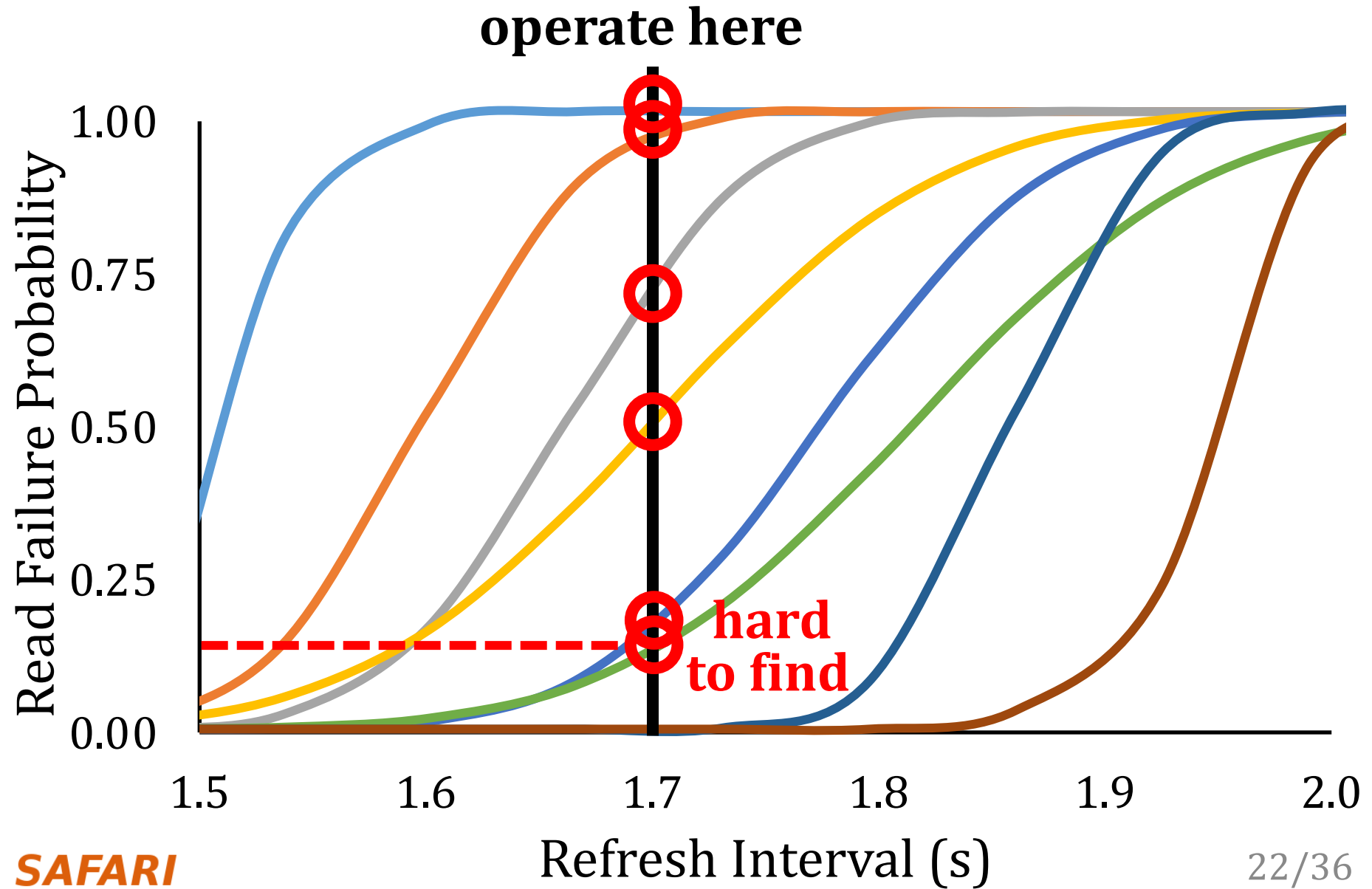


# Single-cell Failure Probability (Real)

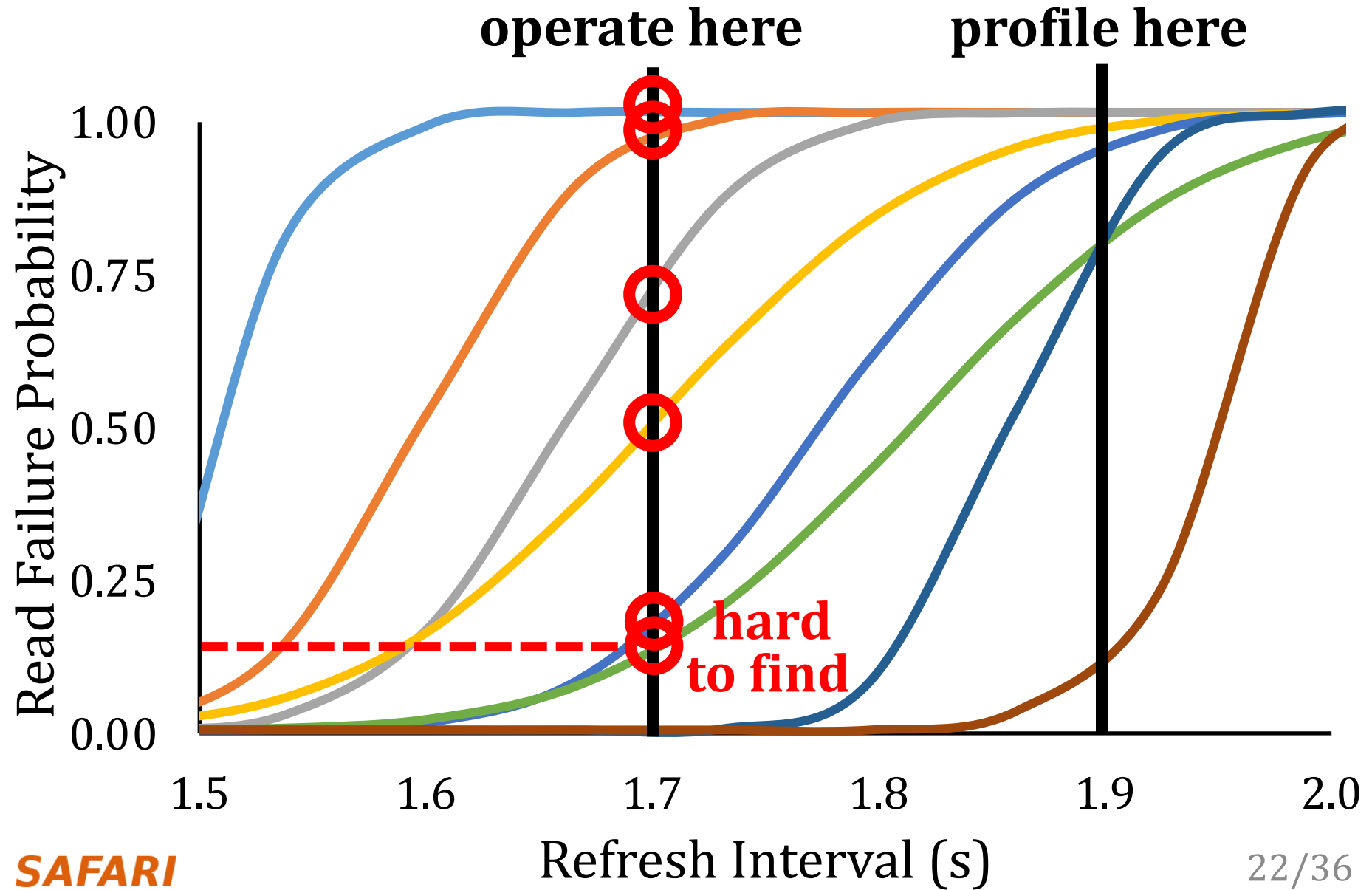
# operate here



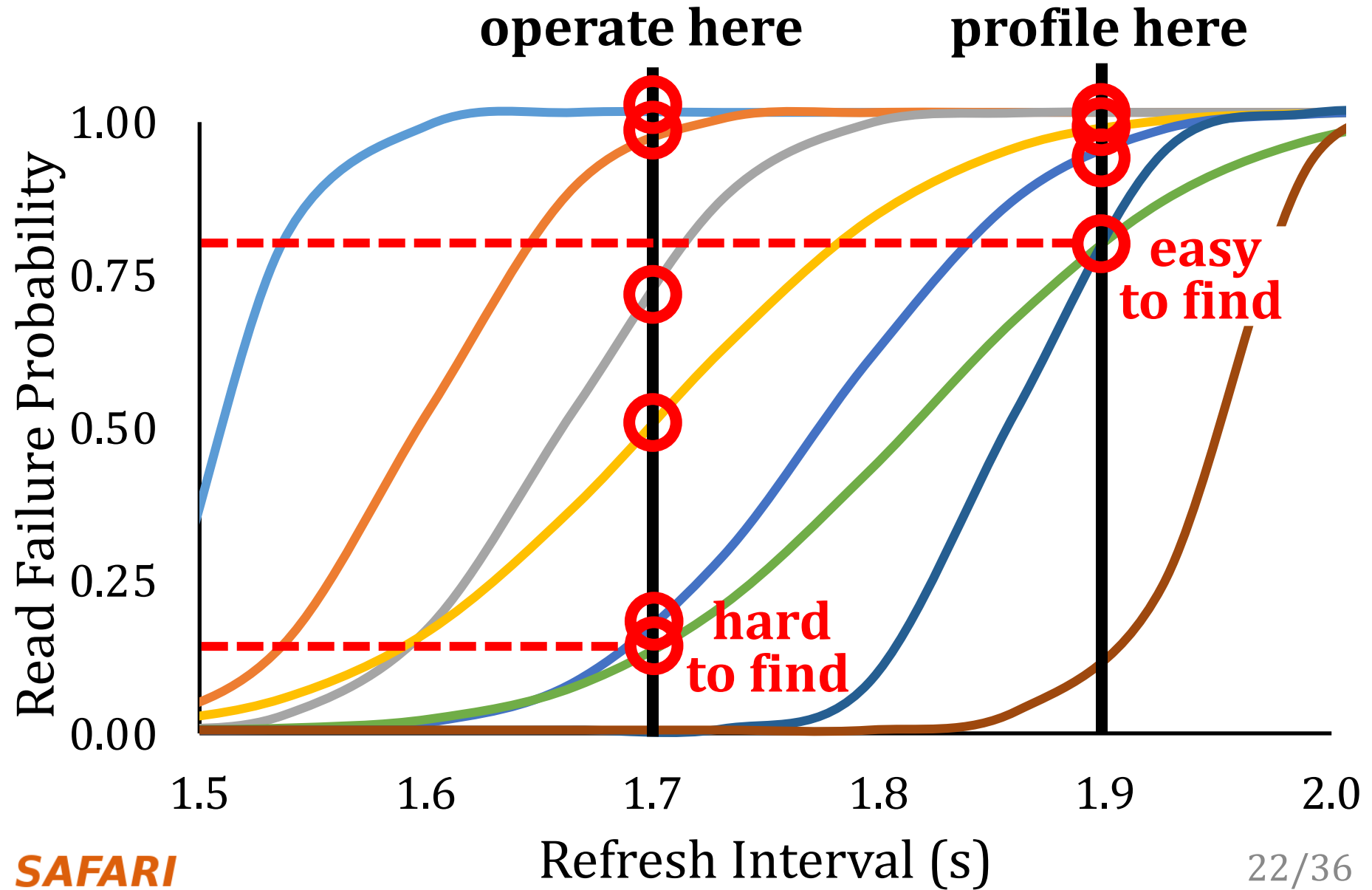
# Single-cell Failure Probability (Real)



# Single-cell Failure Probability (Real)

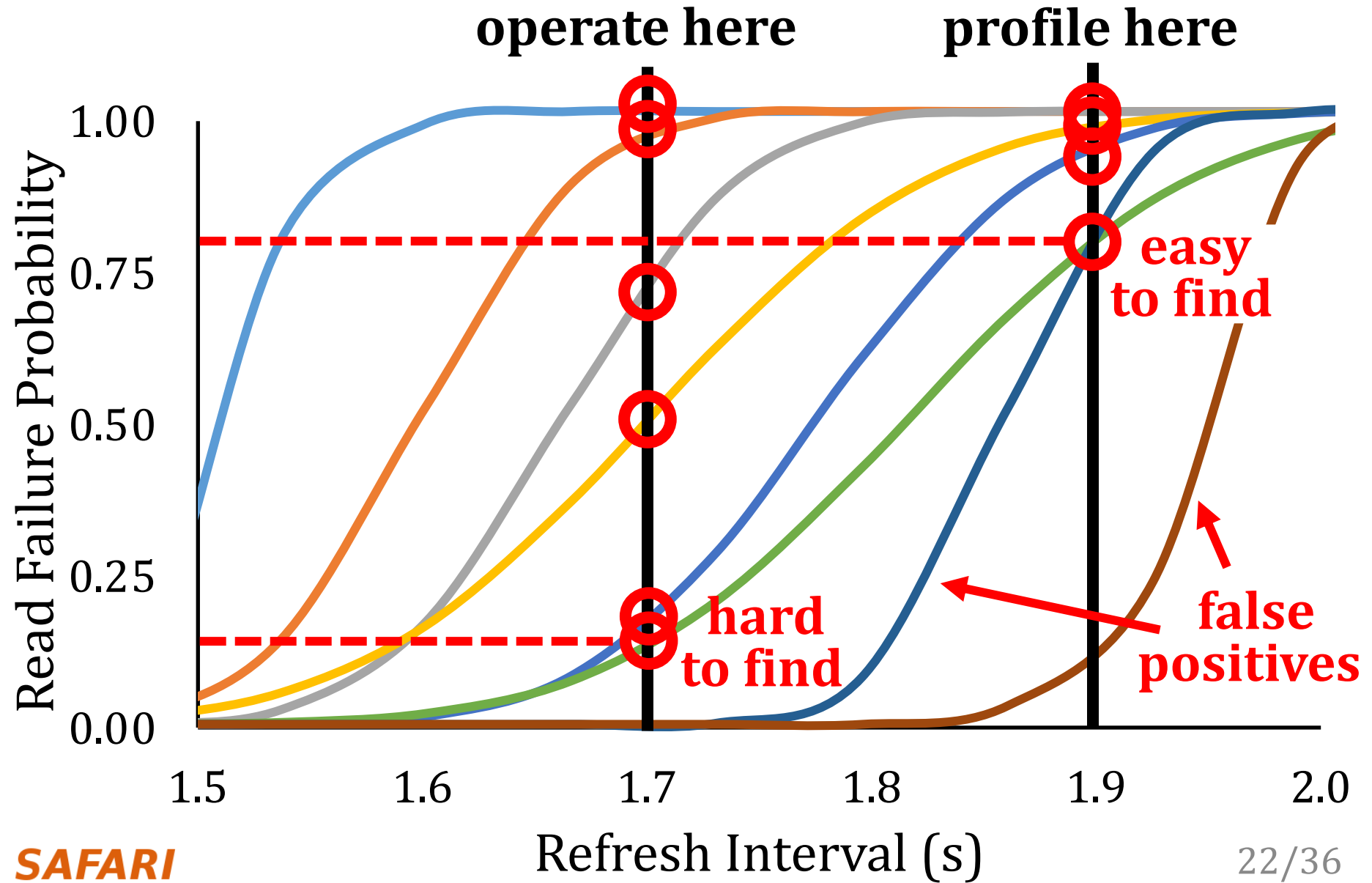


# Single-cell Failure Probability (Real)

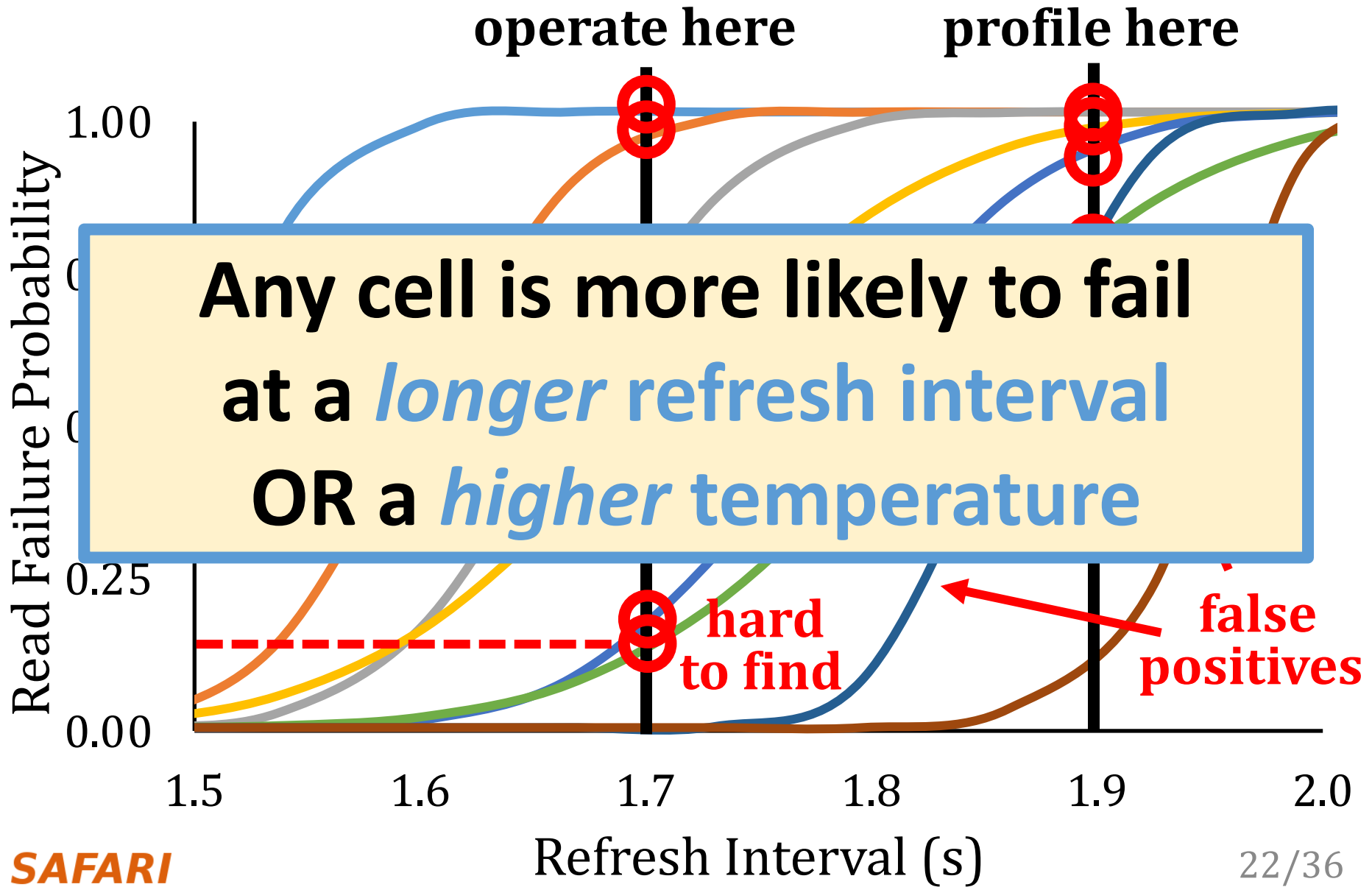




# Single-cell Failure Probability (Real)



# Single-cell Failure Probability (Real)



# Reach Profiling

**Key idea:** profile at a *longer refresh interval* and/or a *higher temperature*

- **Pros**

- **Fast + Reliable:** reach profiling searches for cells *where they are most likely to fail*

- **Cons**

- **False Positives:** profiler may identify cells that fail under profiling conditions, but not under operating conditions

# Towards an Implementation

Reach profiling is a **general methodology**

3 key questions for an implementation:

**What are desirable profiling conditions?**

**How often should the system profile?**

**What information does the profiler need?**

# Three Key Profiling Metrics

- 1. Runtime:** how long profiling takes
- 2. Coverage:** portion of all possible failures discovered by profiling
- 3. False positives:** number of cells observed to fail during profiling but never during actual operation

# Three Key Profiling Metrics

- 1. Runtime:** how long profiling takes
- 2. Coverage:** portion of all possible failures discovered by profiling

We explore how these three metrics change under **many** different profiling conditions

# Simulated End-to-end Performance

 Brute-force profiling    REAPER    Ideal profiling

**On average, REAPER enables:**

**16.3% system performance improvement**

**36.4% DRAM power reduction**



**REAPER enables longer refresh intervals,  
which are unreasonable  
using brute-force profiling**



# REAPER Summary

## Problem:

- DRAM refresh performance and energy overhead is high
- Current approaches to retention failure profiling are slow or unreliable

## Goals:

1. Thoroughly analyze profiling tradeoffs
2. Develop a **fast** and **reliable** profiling mechanism

## Key Contributions:

1. **First** detailed characterization of 368 LPDDR4 DRAM chips
2. **Reach profiling:** Profile at a **longer refresh interval** or **higher temperature** than target conditions, where cells are more likely to fail

## Evaluation:

- **2.5x** faster profiling with **99%** coverage and **50%** false positives
- REAPER enables **16.3% system performance improvement** and **36.4% DRAM power reduction**
- Enables longer refresh intervals that were previously unreasonable

**Main Memory Needs**  
**Intelligent Controllers**  
**for Reliability & Security**

# Understanding In-DRAM ECC

---

- Minesh Patel, Jeremie S. Kim, Hasan Hassan, and Onur Mutlu,  
**"Understanding and Modeling On-Die Error Correction in Modern DRAM: An Experimental Study Using Real Devices"**  
*Proceedings of the 49th Annual IEEE/IFIP International Conference on Dependable Systems and Networks (DSN)*, Portland, OR, USA, June 2019.  
[[Source Code for EINSim, the Error Inference Simulator](#)]  
***Best paper session.***

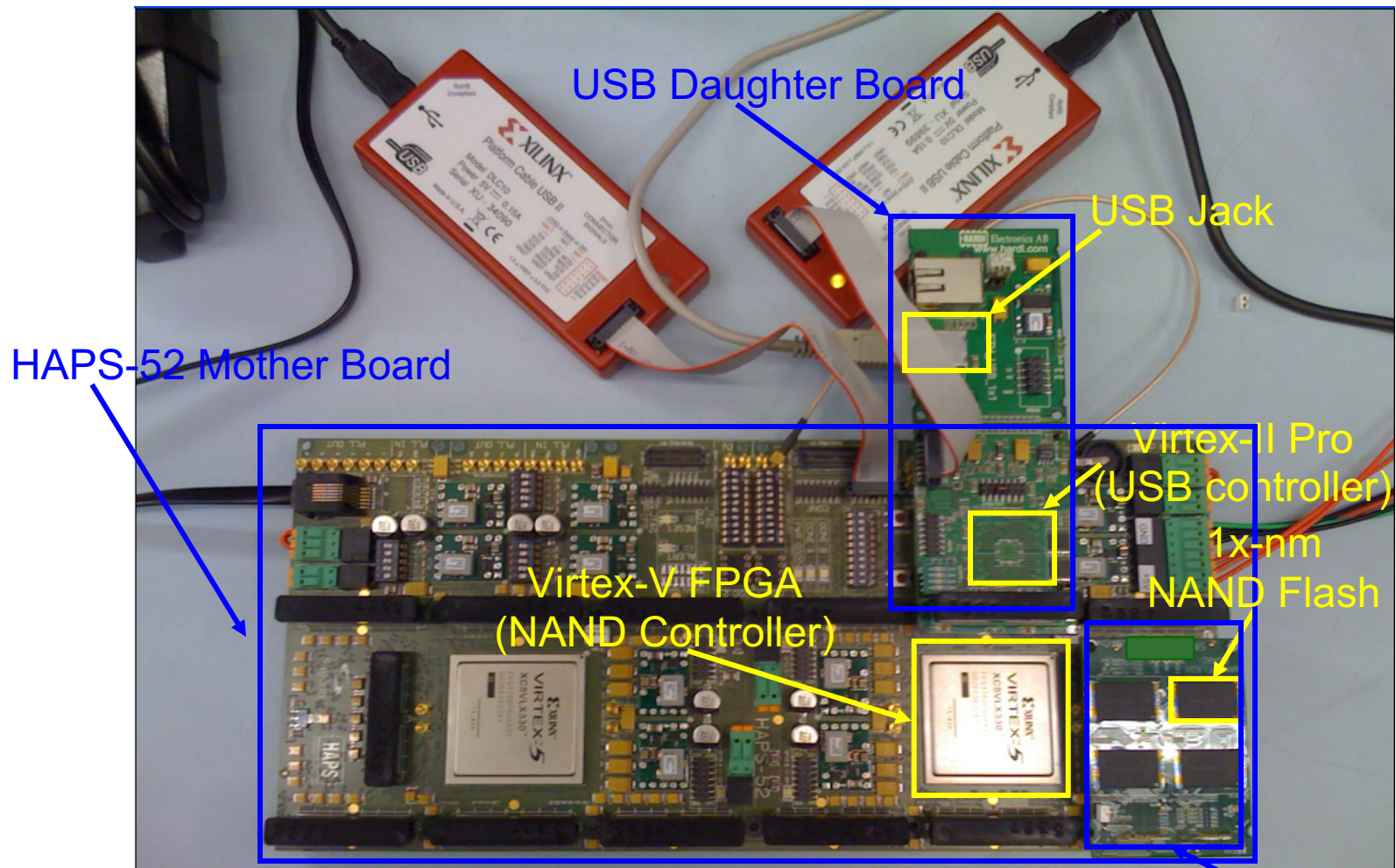
## Understanding and Modeling On-Die Error Correction in Modern DRAM: An Experimental Study Using Real Devices

Minesh Patel<sup>†</sup>   Jeremie S. Kim<sup>‡†</sup>   Hasan Hassan<sup>†</sup>   Onur Mutlu<sup>†‡</sup>

<sup>†</sup>*ETH Zürich*   <sup>‡</sup>*Carnegie Mellon University*

# Understanding Flash Memory Vulnerabilities

# Understand and Model with Experiments (Flash)



[DATE 2012, ICCD 2012, DATE 2013, ITJ 2013, ICCD 2013, SIGMETRICS 2014, HPCA 2015, DSN 2015, MSST 2015, JSAC 2016, HPCA 2017, DFRWS 2017, PIEEE 2017, HPCA 2018, SIGMETRICS 2018]

Cai+, "Error Characterization, Mitigation, and Recovery in Flash Memory Based Solid State Drives," Proc. IEEE 2017.



*Proceedings of the IEEE, Sept. 2017*

## Error Characterization, Mitigation, and Recovery in Flash-Memory-Based Solid-State Drives

*This paper reviews the most recent advances in solid-state drive (SSD) error characterization, mitigation, and data recovery techniques to improve both SSD's reliability and lifetime.*

By YU CAI, SAUGATA GHOSE, ERICH F. HARATSCH, YIXIN LUO, AND ONUR MUTLU



# Understanding Flash Memory Reliability

---

- Justin Meza, Qiang Wu, Sanjeev Kumar, and Onur Mutlu,  
**"A Large-Scale Study of Flash Memory Errors in the Field"**  
*Proceedings of the ACM International Conference on Measurement and Modeling of Computer Systems (**SIGMETRICS**), Portland, OR, June 2015.*  
[[Slides \(pptx\)](#)] [[pdf](#)] [[Coverage at ZDNet](#)] [[Coverage on The Register](#)]  
[[Coverage on TechSpot](#)] [[Coverage on The Tech Report](#)]

## A Large-Scale Study of Flash Memory Failures in the Field

Justin Meza  
Carnegie Mellon University  
[meza@cmu.edu](mailto:meza@cmu.edu)

Qiang Wu  
Facebook, Inc.  
[qw@fb.com](mailto:qw@fb.com)

Sanjeev Kumar  
Facebook, Inc.  
[skumar@fb.com](mailto:skumar@fb.com)

Onur Mutlu  
Carnegie Mellon University  
[onur@cmu.edu](mailto:onur@cmu.edu)

# NAND Flash Vulnerabilities [HPCA'17]

*HPCA, Feb. 2017*

## Vulnerabilities in MLC NAND Flash Memory Programming: Experimental Analysis, Exploits, and Mitigation Techniques

Yu Cai<sup>†</sup>   Saugata Ghose<sup>†</sup>   Yixin Luo<sup>††</sup>   Ken Mai<sup>†</sup>   Onur Mutlu<sup>§†</sup>   Erich F. Haratsch<sup>‡</sup>  
<sup>†</sup>Carnegie Mellon University   <sup>‡</sup>Seagate Technology   <sup>§</sup>ETH Zürich

*Modern NAND flash memory chips provide high density by storing two bits of data in each flash cell, called a multi-level cell (MLC). An MLC partitions the threshold voltage range of a flash cell into four voltage states. When a flash cell is programmed, a high voltage is applied to the cell. Due to parasitic capacitance coupling between flash cells that are physically close to each other, flash cell programming can lead to cell-to-cell program interference, which introduces errors into neighboring flash cells. In order to reduce the impact of cell-to-cell interference on the reliability of MLC NAND flash memory, flash manufacturers adopt a two-step programming method, which programs the MLC in two separate steps. First, the flash memory partially programs the least significant bit of the MLC to some intermediate threshold voltage. Second, it programs the most significant bit to bring the MLC up to its full voltage state.*

*In this paper, we demonstrate that two-step programming exposes new reliability and security vulnerabilities. We expe-*

*belongs to a different flash memory page (the unit of data programmed and read at the same time), which we refer to, respectively, as the least significant bit (LSB) page and the most significant bit (MSB) page [5].*

*A flash cell is programmed by applying a large voltage on the control gate of the transistor, which triggers charge transfer into the floating gate, thereby increasing the threshold voltage. To precisely control the threshold voltage of the cell, the flash memory uses incremental step pulse programming (ISPP) [12, 21, 25, 41]. ISPP applies multiple short pulses of the programming voltage to the control gate, in order to increase the cell threshold voltage by some small voltage amount ( $V_{step}$ ) after each step. Initial MLC designs programmed the threshold voltage in one shot, issuing all of the pulses back-to-back to program both bits of data at the same time. However, as flash memory scales down, the distance between neighboring flash cells decreases, which*

[https://people.inf.ethz.ch/omutlu/pub/flash-memory-programming-vulnerabilities\\_hpca17.pdf](https://people.inf.ethz.ch/omutlu/pub/flash-memory-programming-vulnerabilities_hpca17.pdf)



# 3D NAND Flash Reliability I [HPCA'18]

---

- Yixin Luo, Saugata Ghose, Yu Cai, Erich F. Haratsch, and Onur Mutlu, **"HeatWatch: Improving 3D NAND Flash Memory Device Reliability by Exploiting Self-Recovery and Temperature-Awareness"**

*Proceedings of the 24th International Symposium on High-Performance Computer Architecture (HPCA)*, Vienna, Austria, February 2018.

[[Lightning Talk Video](#)]

[[Slides \(pptx\)](#)] [[pdf](#)] [[Lightning Session Slides \(pptx\)](#)] [[pdf](#)]

## HeatWatch: Improving 3D NAND Flash Memory Device Reliability by Exploiting Self-Recovery and Temperature Awareness

Yixin Luo<sup>†</sup>      Saugata Ghose<sup>†</sup>      Yu Cai<sup>‡</sup>      Erich F. Haratsch<sup>‡</sup>      Onur Mutlu<sup>§†</sup>  
<sup>†</sup>*Carnegie Mellon University*      <sup>‡</sup>*Seagate Technology*      <sup>§</sup>*ETH Zürich*

# 3D NAND Flash Reliability II [SIGMETRICS'18]

---

- Yixin Luo, Saugata Ghose, Yu Cai, Erich F. Haratsch, and Onur Mutlu,  
**"Improving 3D NAND Flash Memory Lifetime by Tolerating Early Retention Loss and Process Variation"**  
*Proceedings of the ACM International Conference on Measurement and Modeling of Computer Systems (**SIGMETRICS**), Irvine, CA, USA, June 2018.*  
[Abstract]  
[POMACS Journal Version (same content, different format)]  
[Slides (pptx) (pdf)]

## Improving 3D NAND Flash Memory Lifetime by Tolerating Early Retention Loss and Process Variation

Yixin Luo<sup>†</sup>      Saugata Ghose<sup>†</sup>      Yu Cai<sup>†</sup>      Erich F. Haratsch<sup>‡</sup>      Onur Mutlu<sup>§†</sup>

<sup>†</sup>Carnegie Mellon University

<sup>‡</sup>Seagate Technology

<sup>§</sup>ETH Zürich

# Another Talk: NAND Flash Memory Robustness

---

- Yu Cai, Saugata Ghose, Erich F. Haratsch, Yixin Luo, and Onur Mutlu,  
**"Error Characterization, Mitigation, and Recovery in Flash Memory Based Solid State Drives"**

*to appear in Proceedings of the IEEE, 2017.*

Cai+, "Error Patterns in MLC NAND Flash Memory: Measurement, Characterization, and Analysis," DATE 2012.

Cai+, "Flash Correct-and-Refresh: Retention-Aware Error Management for Increased Flash Memory Lifetime," ICCD 2012.

Cai+, "Threshold Voltage Distribution in MLC NAND Flash Memory: Characterization, Analysis and Modeling," DATE 2013.

Cai+, "Error Analysis and Retention-Aware Error Management for NAND Flash Memory," Intel Technology Journal 2013.

Cai+, "Program Interference in MLC NAND Flash Memory: Characterization, Modeling, and Mitigation," ICCD 2013.

Cai+, "Neighbor-Cell Assisted Error Correction for MLC NAND Flash Memories," SIGMETRICS 2014.

Cai+, "Data Retention in MLC NAND Flash Memory: Characterization, Optimization and Recovery," HPCA 2015.

Cai+, "Read Disturb Errors in MLC NAND Flash Memory: Characterization and Mitigation," DSN 2015.

Luo+, "WARM: Improving NAND Flash Memory Lifetime with Write-hotness Aware Retention Management," MSST 2015.

Meza+, "A Large-Scale Study of Flash Memory Errors in the Field," SIGMETRICS 2015.

Luo+, "Enabling Accurate and Practical Online Flash Channel Modeling for Modern MLC NAND Flash Memory," IEEE JSAC 2016.

Cai+, "Vulnerabilities in MLC NAND Flash Memory Programming: Experimental Analysis, Exploits, and Mitigation Techniques," HPCA 2017.

Fukami+, "Improving the Reliability of Chip-Off Forensic Analysis of NAND Flash Memory Devices," DFRWS EU 2017.

Luo+, "HeatWatch: Improving 3D NAND Flash Memory Device Reliability by Exploiting Self-Recovery and Temperature-Awareness," HPCA 2018.

Luo+, "Improving 3D NAND Flash Memory Lifetime by Tolerating Early Retention Loss and Process Variation," SIGMETRICS 2018.

---

Cai+, "Error Characterization, Mitigation, and Recovery in Flash Memory Based Solid State Drives," Proc. IEEE 2017.

# Two Other Solution Directions

# There are Two Other Solution Directions

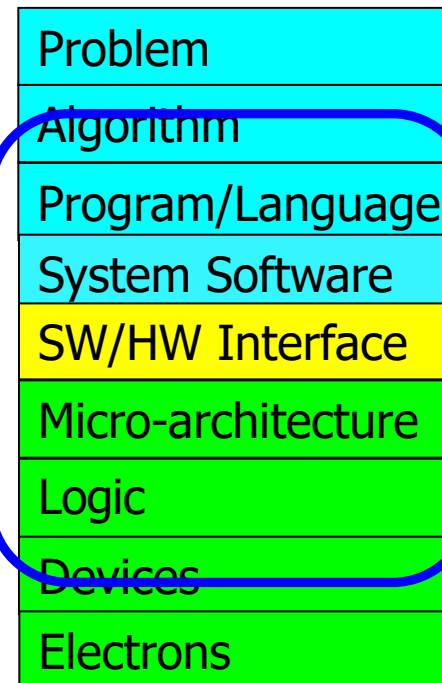
- **New Technologies:** Replace or (more likely) augment DRAM with a different technology

- ❑ Non-volatile memories

- **Embracing Un-reliability:**

Design memories with different reliability and store data intelligently across them

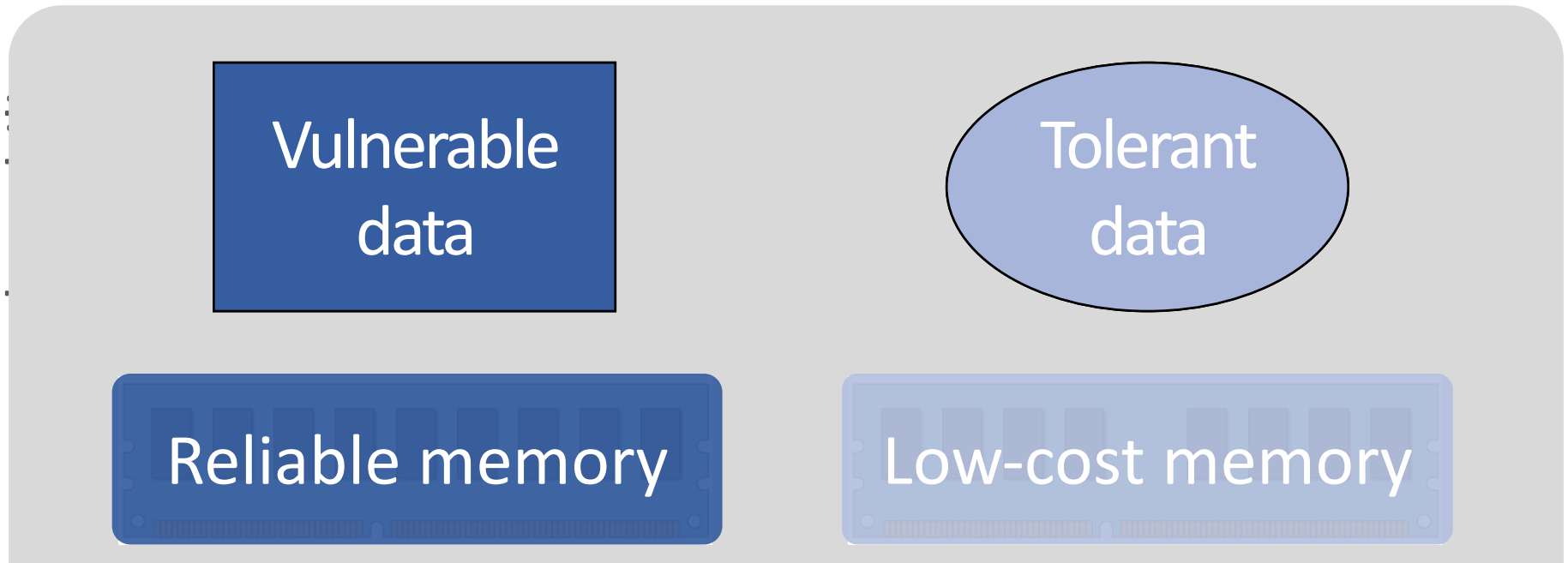
**[Luo+ DSN 2014]**



- ...

**Fundamental solutions to security  
require co-design across the hierarchy**

# Exploiting Memory Error Tolerance with Hybrid Memory Systems



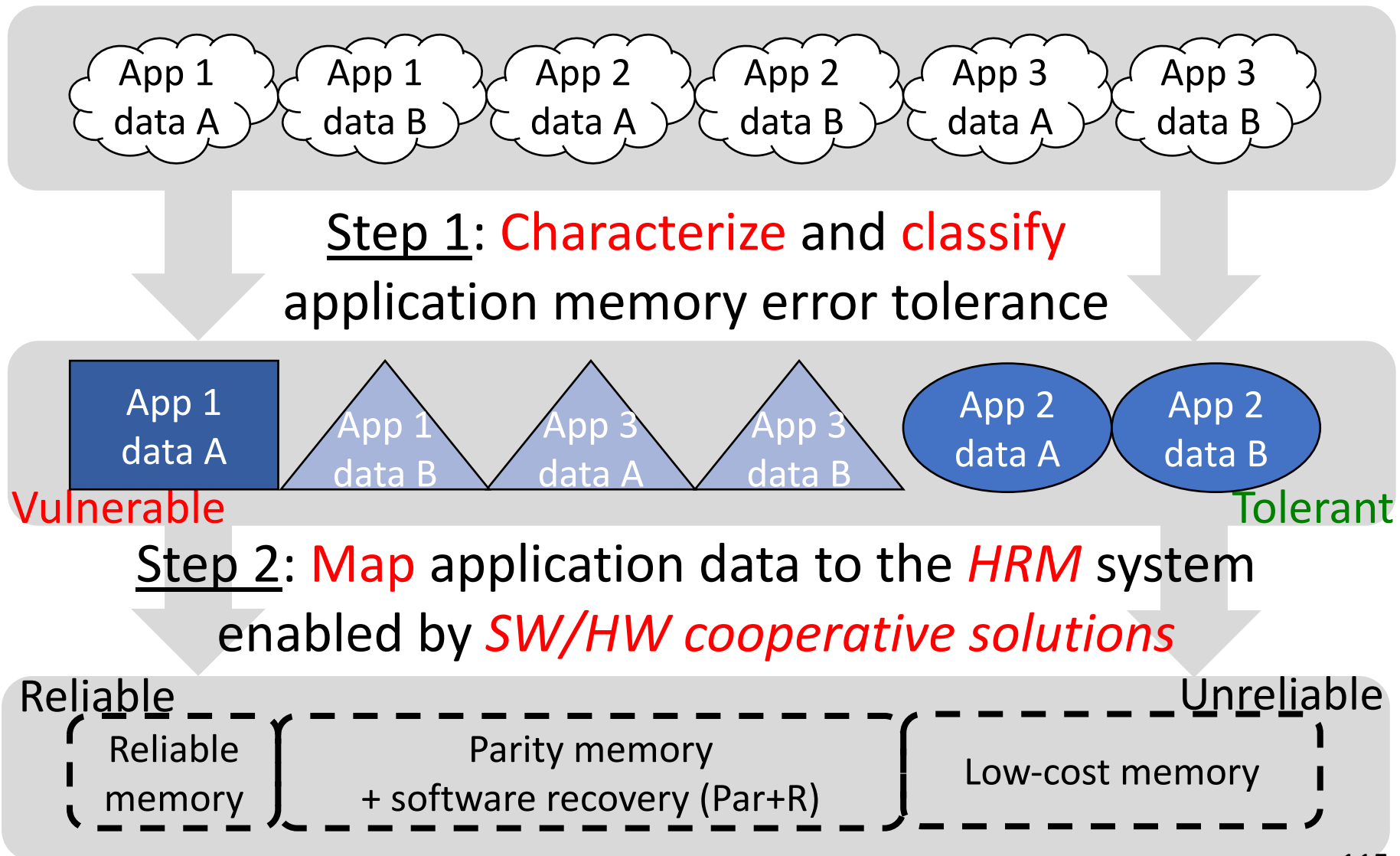
On Microsoft's Web Search workload

Reduces server hardware **cost** by **4.7 %**

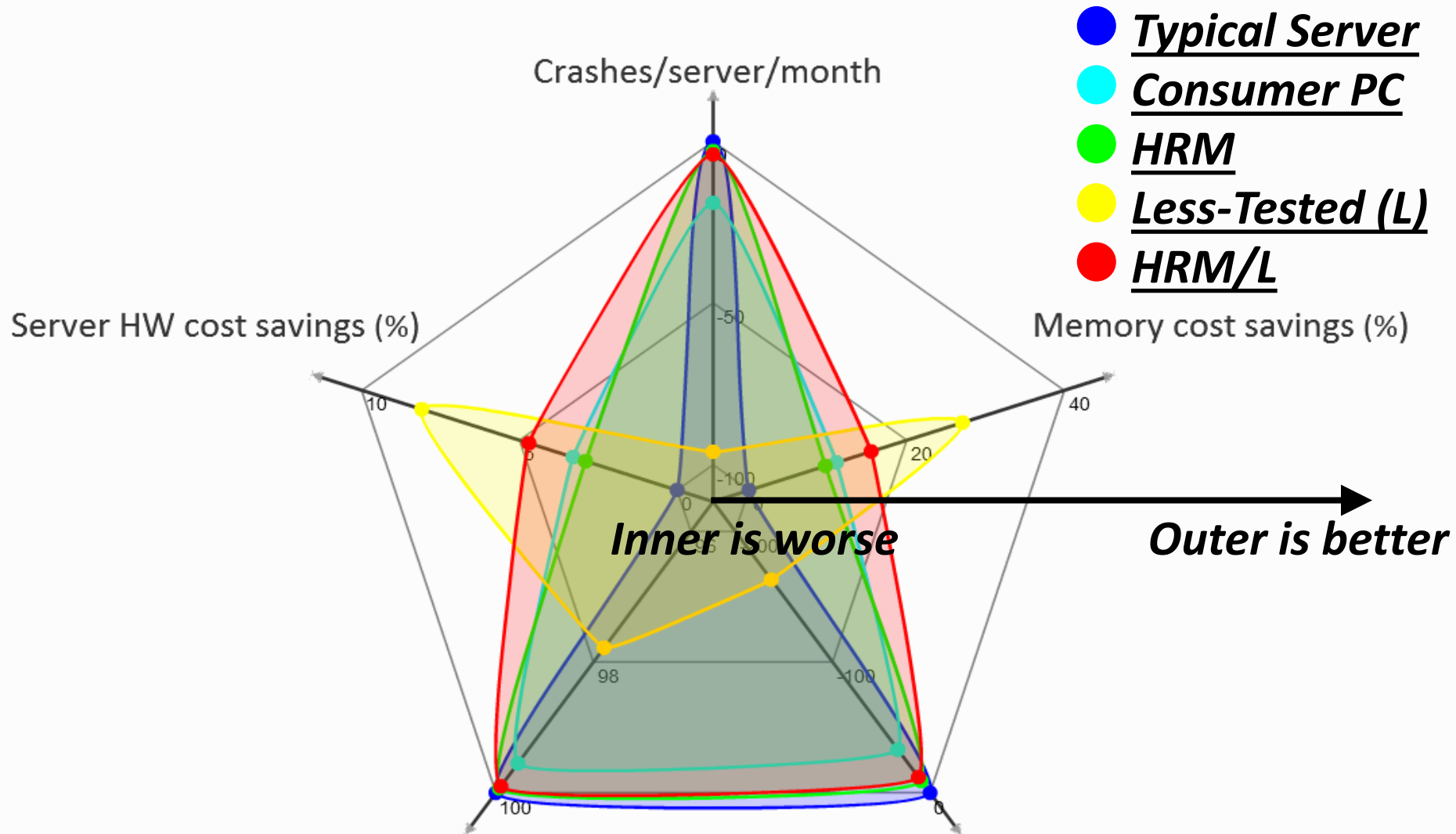
Achieves single server **availability** target of **99.90 %**

**Heterogeneous-Reliability Memory** [DSN 2014]

# Heterogeneous-Reliability Memory



# Evaluation Results



● ● Bigger area means better tradeoff



# More on Heterogeneous-Reliability Memory

---

- Yixin Luo, Sriram Govindan, Bikash Sharma, Mark Santaniello, Justin Meza, Aman Kansal, Jie Liu, Badriddine Khessib, Kushagra Vaid, and Onur Mutlu,  
**"Characterizing Application Memory Error Vulnerability to Optimize Data Center Cost via Heterogeneous-Reliability Memory"**  
*Proceedings of the 44th Annual IEEE/IFIP International Conference on Dependable Systems and Networks (DSN)*, Atlanta, GA, June 2014. [[Summary](#)]  
[[Slides \(pptx\)](#)] [[pdf](#)] [[Coverage on ZDNet](#)]

## Characterizing Application Memory Error Vulnerability to Optimize Datacenter Cost via Heterogeneous-Reliability Memory

Yixin Luo   Sriram Govindan\*   Bikash Sharma\*   Mark Santaniello\*   Justin Meza  
Aman Kansal\*   Jie Liu\*   Badriddine Khessib\*   Kushagra Vaid\*   Onur Mutlu

Carnegie Mellon University, yixinluo@cs.cmu.edu, {meza, onur}@cmu.edu

\*Microsoft Corporation, {srgovin, bsharma, marksan, kansal, jie.liu, bknessib, kvaid}@microsoft.com