Computer Architecture Lecture 2a: Memory Systems: Challenges and Opportunities

Prof. Onur Mutlu ETH Zürich Fall 2022 30 September 2022

Four Key Directions

Fundamentally Secure/Reliable/Safe Architectures

Fundamentally Energy-Efficient Architectures
 Memory-centric (Data-centric) Architectures

Fundamentally Low-Latency and Predictable Architectures

Architectures for AI/ML, Genomics, Medicine, Health

Memory & Storage

Why Is Memory So Important? (Especially Today)

Importance of Main Memory

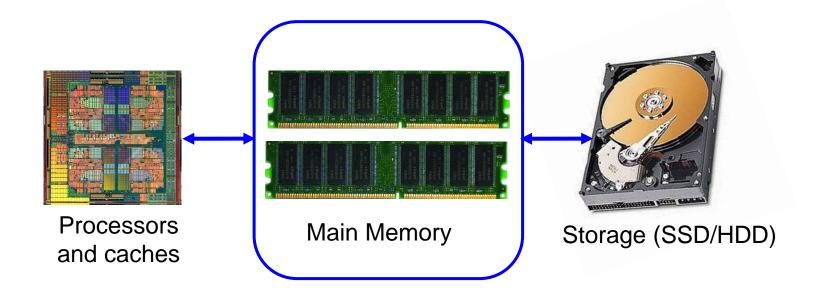
The Performance Perspective

The Energy Perspective

The Scaling/Reliability/Security Perspective

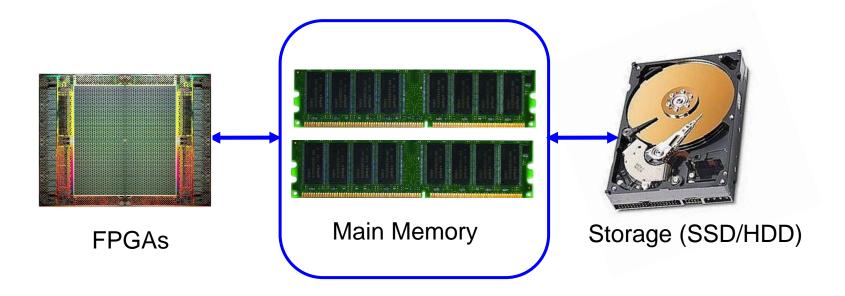
Trends/Challenges/Opportunities in Main Memory

The Main Memory System



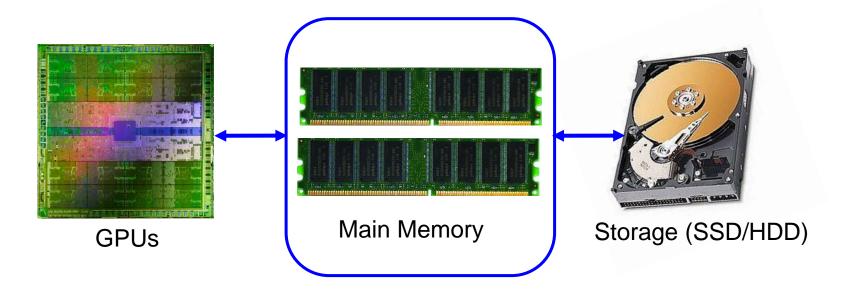
- Main memory is a critical component of all computing systems: server, mobile, embedded, desktop, sensor
- Main memory system must scale (in size, technology, efficiency, cost, and management algorithms) to maintain performance growth and technology scaling benefits

The Main Memory System



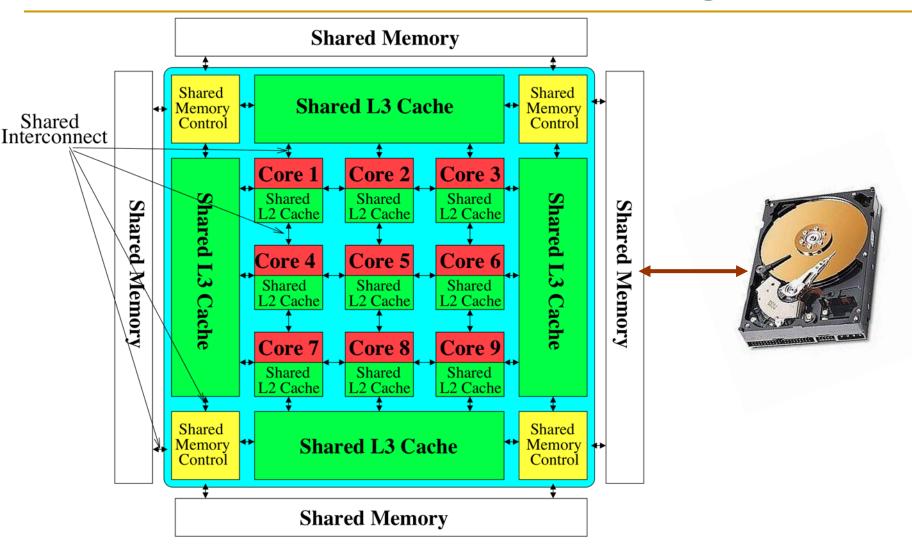
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Perils of Processor-Centric Design



Most of the system is dedicated to storing and moving data

SAFARI Yet, system is still bottlenecked by memory

State of the Main Memory System

- Recent technology, architecture, and application trends
 - lead to new requirements
 - exacerbate old requirements
- DRAM and memory controllers, as we know them today, are (will be) unlikely to satisfy all requirements
- Some emerging non-volatile memory technologies (e.g., PCM) enable new opportunities: memory+storage merging
- We need to rethink the main memory system
 to fix DRAM issues and enable emerging technologies
 to satisfy all requirements

Major Trends Affecting Main Memory (I)

Need for main memory capacity, bandwidth, QoS increasing

Main memory energy/power is a key system design concern

DRAM technology scaling is ending

Major Trends Affecting Main Memory (II)

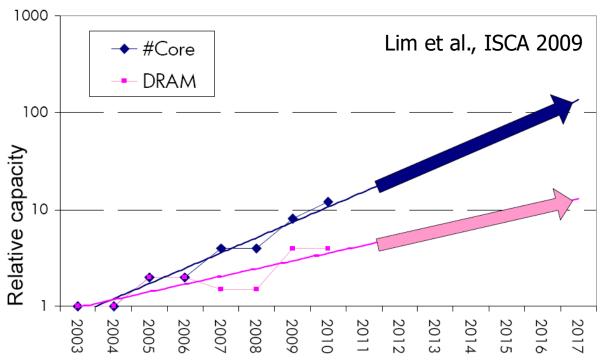
- Need for main memory capacity, bandwidth, QoS increasing
 - Multi-core: increasing number of cores/agents
 - Data-intensive applications: increasing demand/hunger for data
 - Consolidation: cloud computing, GPUs, mobile, heterogeneity

• Main memory energy/power is a key system design concern

DRAM technology scaling is ending

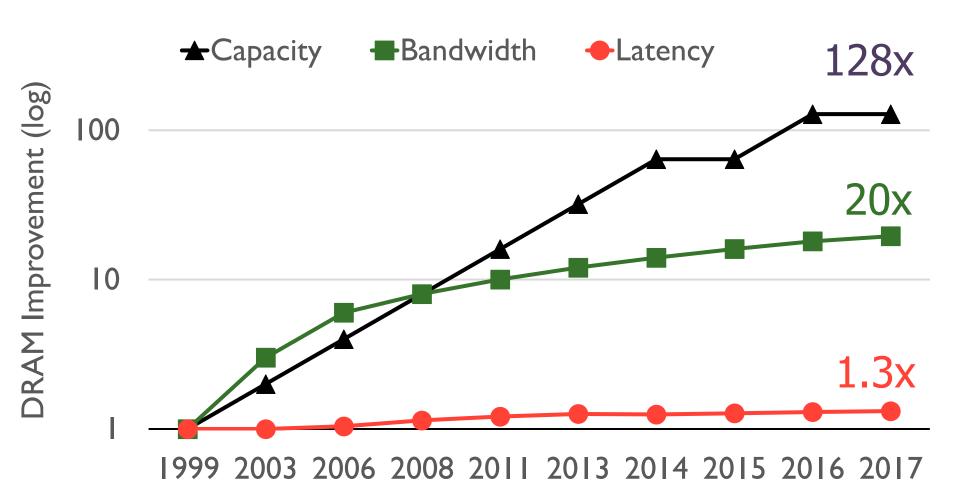
Consequence: The Memory Capacity Gap

Core count doubling ~ every 2 years DRAM DIMM capacity doubling ~ every 3 years



Memory capacity per core expected to drop by 30% every two years
Trends worse for *memory bandwidth per core!*

DRAM Capacity, Bandwidth & Latency





In-memory Databases

[Mao+, EuroSys'12; Clapp+ (**Intel**), IISWC'15]

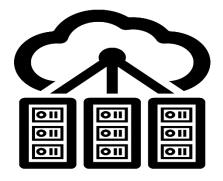


In-Memory Data Analytics

[Clapp+ (**Intel**), IISWC'15; Awan+, BDCloud'15]



Graph/Tree Processing [Xu+, IISWC'12; Umuroglu+, FPL'15]



Datacenter Workloads [Kanev+ (**Google**), ISCA'15]





In-memory Databases

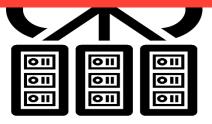
Graph/Tree Processing

Memory → performance bottleneck



In-Memory Data Analytics

[Clapp+ (**Intel**), IISWC'15; Awan+, BDCloud'15]



Datacenter Workloads [Kanev+ (**Google**), ISCA'I 5]



Chrome

Google's web browser



TensorFlow Mobile

Google's machine learning framework



Google's video codec





Memory → performance bottleneck



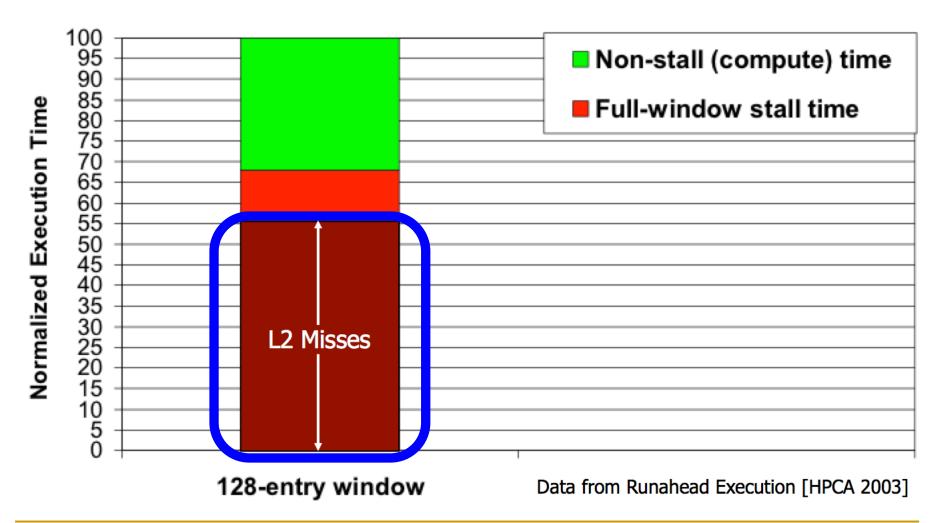
Google's video codec



Memory Bottleneck

I expect that over the coming decade memory subsystem design will be the *only* important design issue for microprocessors.

"It's the Memory, Stupid!" (Richard Sites, MPR, 1996)



Mutlu+, "Runahead Execution: An Alternative to Very Large Instruction Windows for Out-of-Order Processors," HPCA 2003.

The Memory Bottleneck

Onur Mutlu, Jared Stark, Chris Wilkerson, and Yale N. Patt,
 "Runahead Execution: An Alternative to Very Large Instruction Windows for Out-of-order Processors"
 Proceedings of the <u>9th International Symposium on High-Performance Computer Architecture</u> (HPCA), pages 129-140, Anaheim, CA, February 2003.
 [Talk Slides (pdf)]
 [Lecture Slides (pptx) (pdf)]
 [Lecture Video (1 hr 54 mins)]
 [Retrospective HPCA Test of Time Award Talk Slides (pptx) (pdf)]
 [Retrospective HPCA Test of Time Award Talk Video (14 minutes)]
 One of the 15 computer architecture papers of 2003 selected as Top Picks by IEEE Micro.
 HPCA Test of Time Award (awarded in 2021).

Runahead Execution: An Alternative to Very Large Instruction Windows for Out-of-order Processors

Onur Mutlu § Jared Stark † Chris Wilkerson ‡ Yale N. Patt §

§ECE Department The University of Texas at Austin {onur,patt}@ece.utexas.edu †Microprocessor Research Intel Labs jared.w.stark@intel.com

Desktop Platforms Group Intel Corporation chris.wilkerson@intel.com

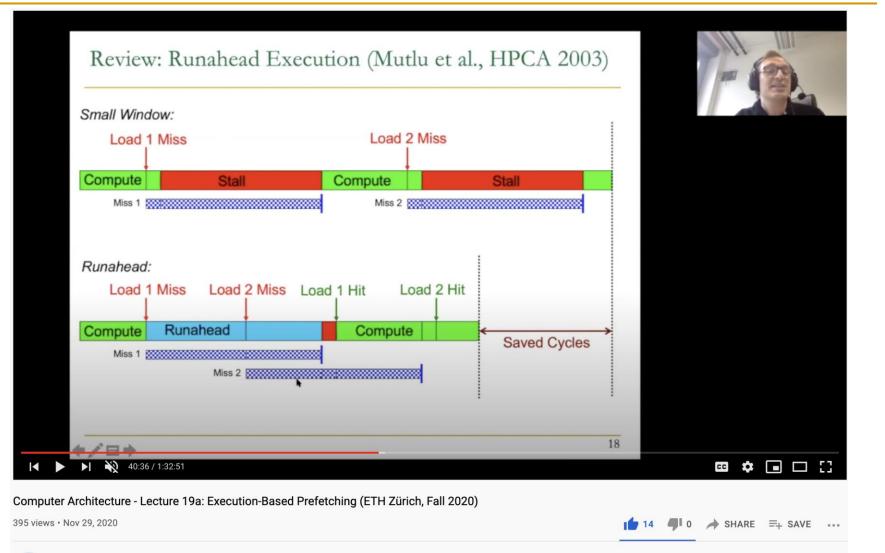
The Memory Bottleneck

 Onur Mutlu, Jared Stark, Chris Wilkerson, and Yale N. Patt, "Runahead Execution: An Effective Alternative to Large <u>Instruction Windows</u>"

IEEE Micro, Special Issue: Micro's Top Picks from Microarchitecture Conferences (*MICRO TOP PICKS*), Vol. 23, No. 6, pages 20-25, November/December 2003.

RUNAHEAD EXECUTION: AN EFFECTIVE ALTERNATIVE TO LARGE INSTRUCTION WINDOWS

More on Runahead Execution (I)



Onur Mutlu Lectures

16.5K subscribers

EDIT VIDEO

ANALYTICS

More on Runahead Execution (II)

Runahead Execution in NVIDIA Denver

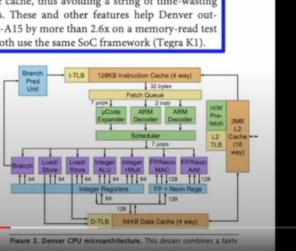
Reducing the effects of long cache-miss penalties has been a major focus of the microarchitecture, using techniques like prefetching and run-ahead. An aggressive hardware prefetcher implementation detects L2 cache requests and tracks up to 32 streams, each with complex stride patterns.

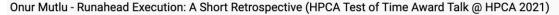
Run-ahead uses the idle time that a CPU spends waiting on a long latency operation to discover cache and DTLB misses further down the instruction stream and generates prefetch requests for these misses.¹ These prefetch requests warm up the data cache and DTLB well before the actual execution of the instructions that require the data. Runahead complements the hardware prefetcher because it's better at prefetching nonstrided streams, and it trains the hardware prefetcher faster than normal execution to yield a combined benefit of 13 percent on SPECint2000 and up to 60 percent on SPECfp2000.

Boggs+, "Denver: NVIDIA's First 64-Bit ARM Processor,"

Gwennap, "NVIDIA's First CPU is a Winner," MPR 2014.

The core includes a hardware prefetch unit that Boggs describes as "aggressive" in preloading the data cache but less aggressive in preloading the instruction cache. It also implements a "run-ahead" feature that continues to execute microcode speculatively after a data-cache miss; this execution can trigger additional cache misses that resolve in the shadow of the first miss. Once the data from the original miss returns, the results of this speculative execution are discarded and execution restarts with the bundle containing the original miss, but run-ahead can preload subsequent data into the cache, thus avoiding a string of time-wasting cache misses. These and other features help Denver outscore Cortex-A15 by more than 2.6x on a memory-read test even when both use the same SoC framework (Tegra K1).







IEEE Micro 2015.



Onur Mutlu Lectures 16.5K subscribers

▶ 💦 6:18 / 14:27

https://www.youtube.com/watch?v=KFCOecRQTIc

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EDIT VIDEO

→ SHARE =+ SAVE

ANALYTICS

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It's the Memory, Stupid!

RICHARD SITES

It's the Memory, Stupid!

When we started the Alpha architecture design in 1988, we estimated a 25-year lifetime and a relatively modest 32% per year compounded performance improvement of implementations over that lifetime (1,000× total). We guestimated about 10× would come from CPU clock improvement, 10× from multiple instruction issue, and 10× from multiple processors.

5, 1996 🏈 MICROPROCESSOR REPORT

An Informal Interview on Memory

Madeleine Gray and Onur Mutlu, "It's the memory, stupid': A conversation with Onur Mutlu" *HiPEAC info 55*, *HiPEAC Newsletter*, October 2018. [Shorter Version in Newsletter] [Longer Online Version with References]

'It's the memory, stupid': A conversation with Onur Mutlu

'We're beyond computation; we know how to do computation really well, we can optimize it, we can build all sorts of accelerators ... but the memory – how to feed the data, how to get the data into the accelerators – is a huge problem.'

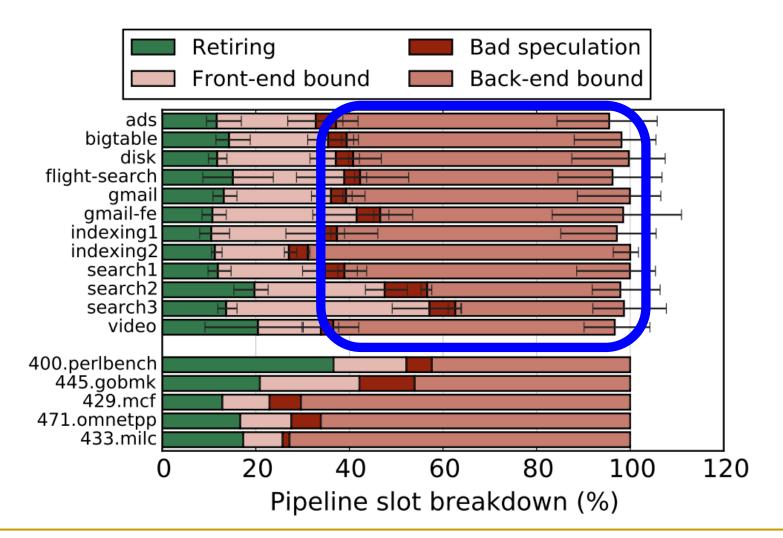
This was how ETH Zürich and Carnegie Mellon Professor Onur Mutlu opened his course on memory systems and memory-centric computing systems at HiPEAC's summer school, ACACES18. A prolific publisher – he recently bagged the top spot on the International Symposium on Computer Architecture (ISCA) hall of fame – Onur is passionate about computation and communication that are efficient and secure by design. In advance of our Computing Systems Week focusing on data centres, storage, and networking, which takes place



next week in Heraklion, HiPEAC picked his brains on all things data-based.

The Memory Bottleneck

All of Google's Data Center Workloads (2015):



Kanev+, "Profiling a Warehouse-Scale Computer," ISCA 2015.

The Memory Bottleneck

All of Google's Data Center Workloads (2015):

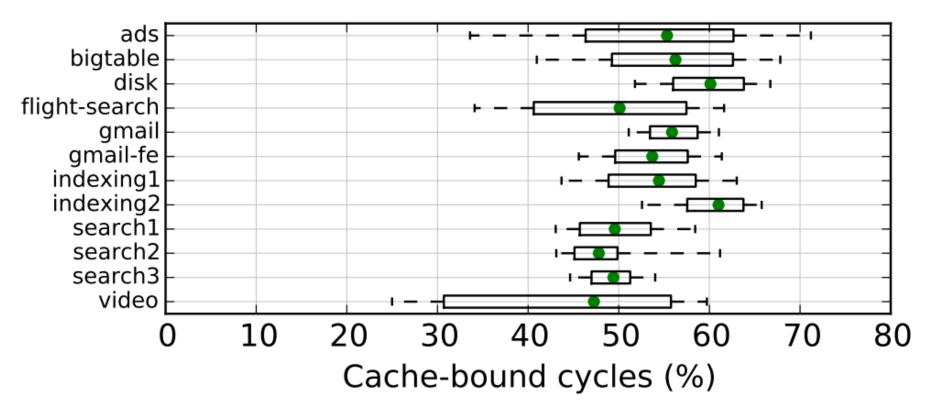
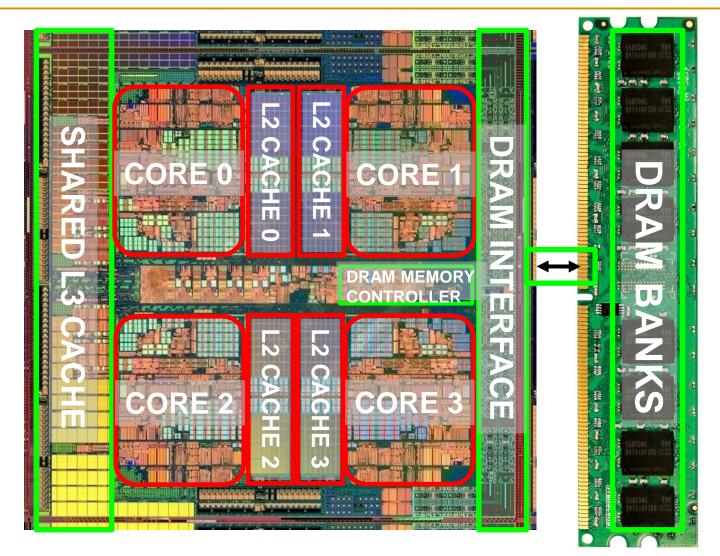


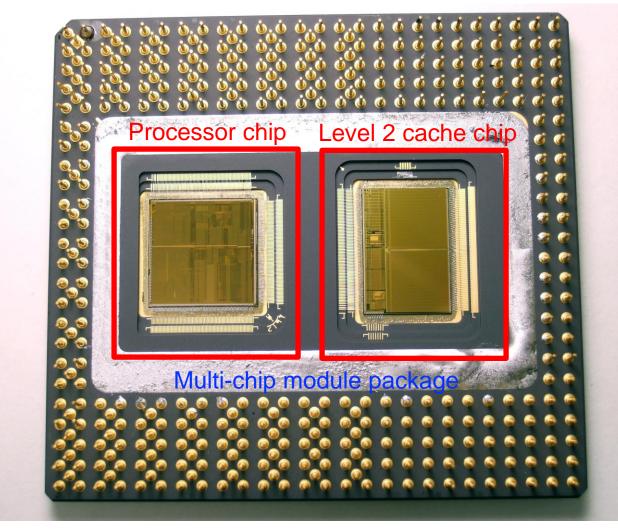
Figure 11: Half of cycles are spent stalled on caches.

Memory in a Modern System



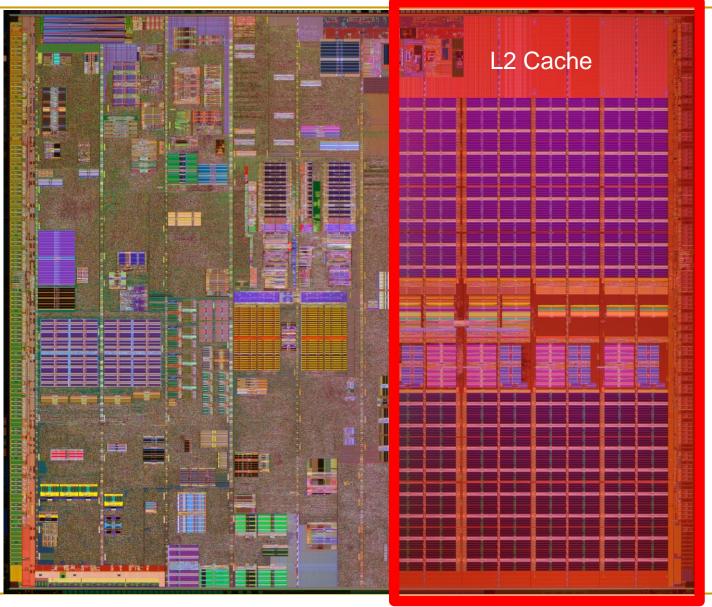
AMD Barcelona, 2006

A Large Fraction of Modern Systems is Memory



Intel Pentium Pro, 1995

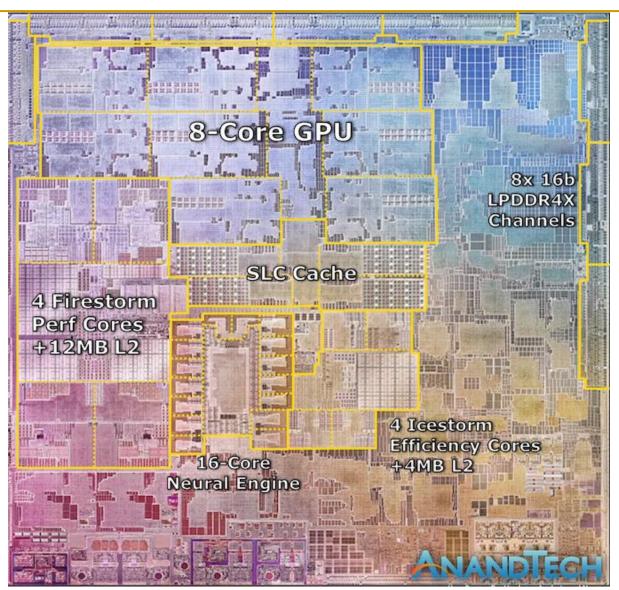
A Large Fraction of Modern Systems is Memory



https://download.intel.com/newsroom/kits/40thanniversary/gallery/images/Pentium_4_6xx-die.jpg

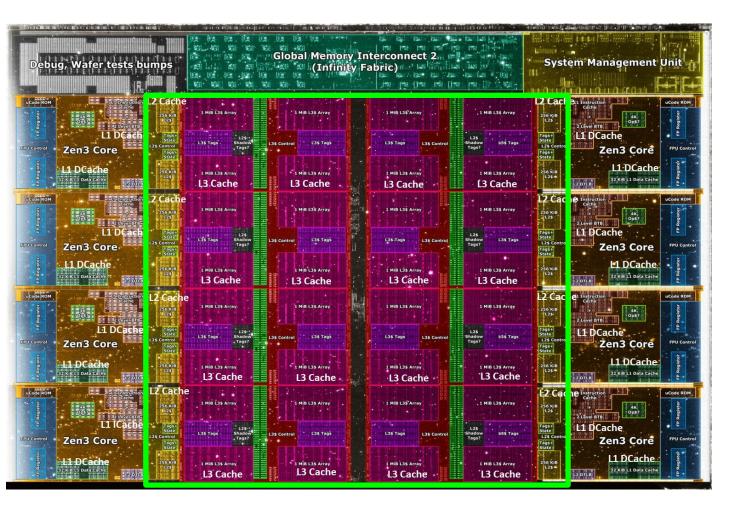
Intel Pentium 4, 2000

30



Apple M1, 2021

Source: https://www.anandtech.com/show/16252/mac-mini-apple-m1-tested



Core Count: 8 cores/16 threads

L1 Caches: 32 KB per core

L2 Caches: 512 KB per core

L3 Cache: 32 MB shared

AMD Ryzen 5000, 2020

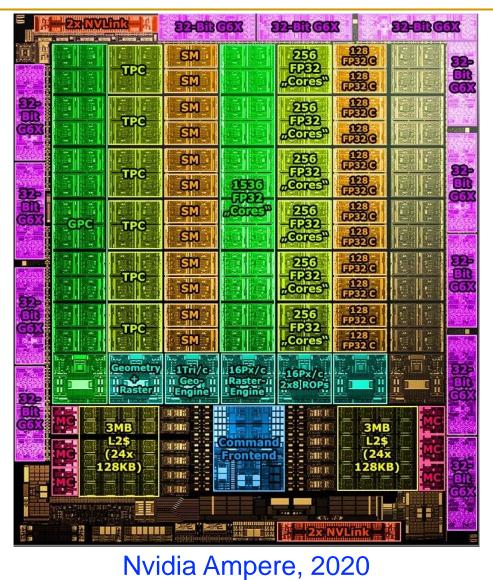


IBM POWER10, 2020

Cores: 15-16 cores, 8 threads/core

L2 Caches: 2 MB per core

L3 Cache: 120 MB shared



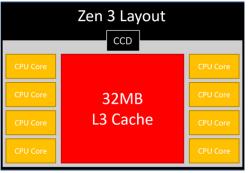
Cores:

128 Streaming Multiprocessors

L1 Cache or Scratchpad: 192KB per SM Can be used as L1 Cache and/or Scratchpad

L2 Cache: 40 MB shared

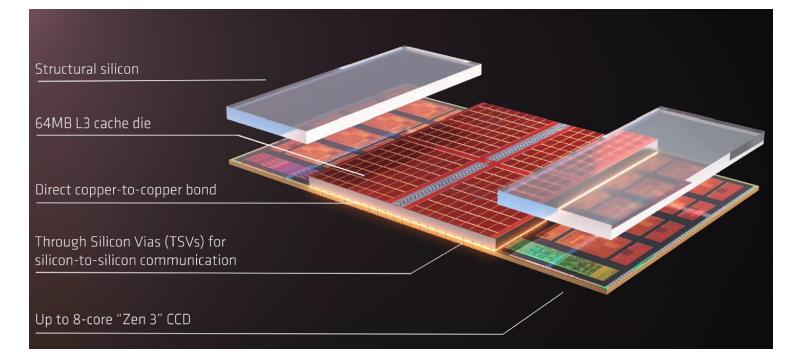
AMD's 3D Last Level Cache (2021)



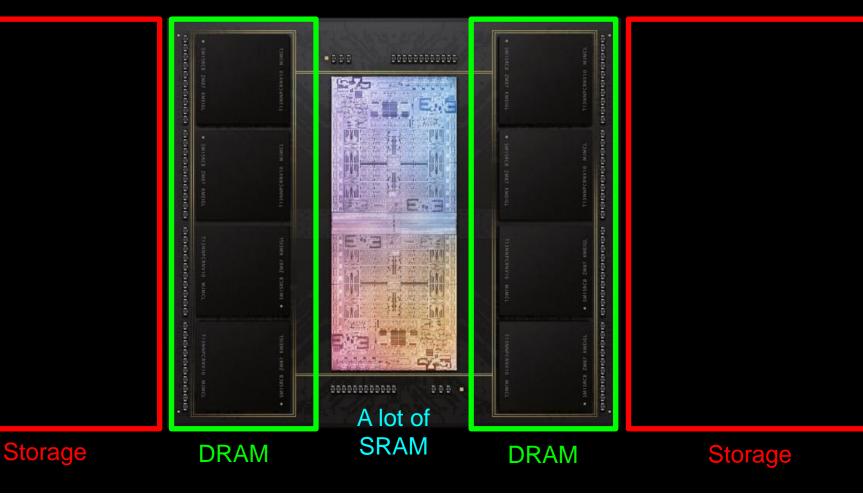
https://community.microcenter.com/discussion/5 134/comparing-zen-3-to-zen-2 AMD increases the L3 size of their 8-core Zen 3 processors from 32 MB to 96 MB

Additional 64 MB L3 cache die stacked on top of the processor die

- Connected using Through Silicon Vias (TSVs)
- Total of 96 MB L3 cache

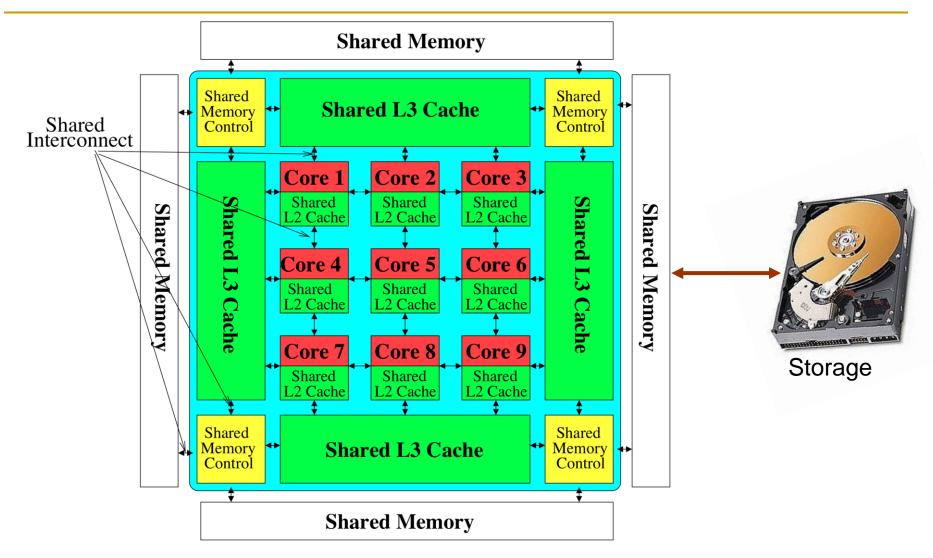


Deeper and Larger Memory Hierarchies



Apple M1 Ultra System (2022)

Memory System: Most of the Platform



Most of the system is dedicated to storing and moving data

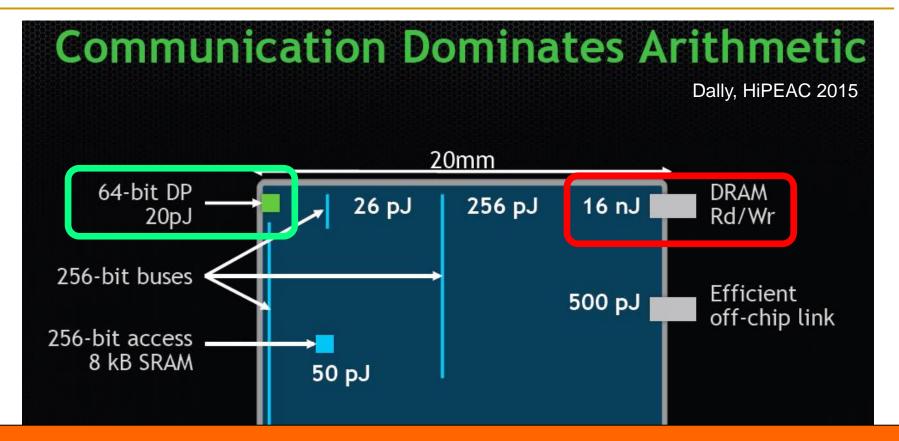
SAFARI Yet, system is still bottlenecked by memory

Major Trends Affecting Main Memory (III)

Need for main memory capacity, bandwidth, QoS increasing

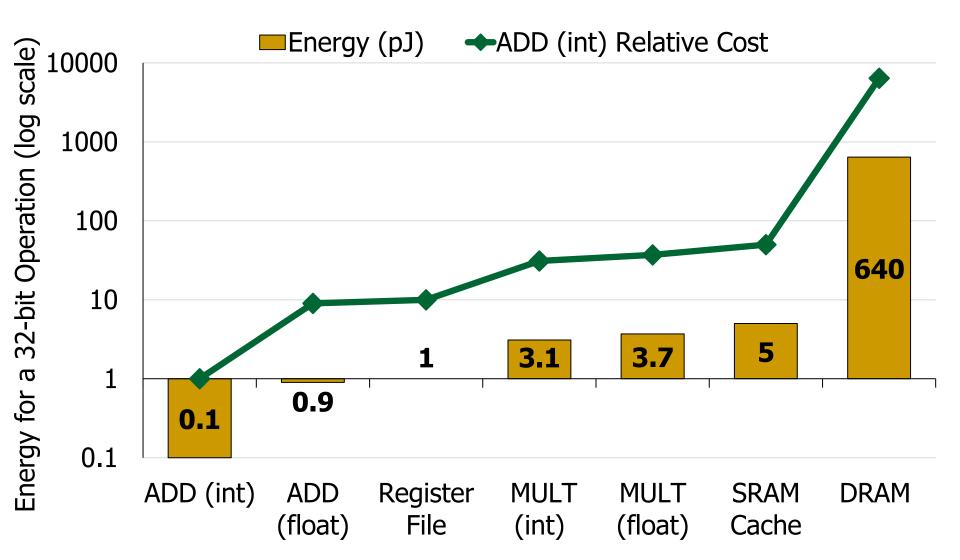
- Main memory energy/power is a key system design concern
 - ~40-50% energy spent in off-chip memory hierarchy [Lefurgy, IEEE Computer'03] >40% power in DRAM [Ware, HPCA'10][Paul,ISCA'15]
 - DRAM consumes power even when not used (periodic refresh)
- DRAM technology scaling is ending

Data Movement vs. Computation Energy

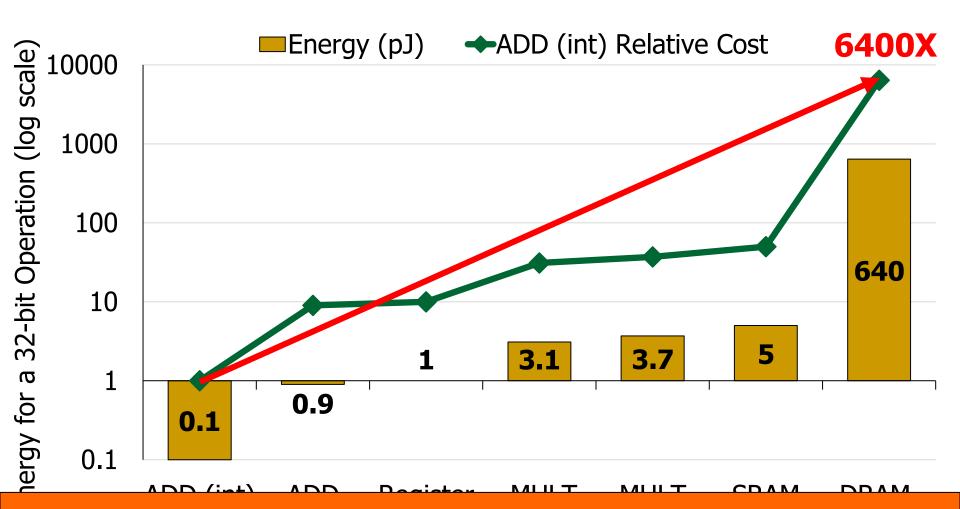


A memory access consumes ~100-1000X the energy of a complex addition

Data Movement vs. Computation Energy



Data Movement vs. Computation Energy



A memory access consumes 6400X the energy of a simple integer addition

Energy Waste in Mobile Devices

 Amirali Boroumand, Saugata Ghose, Youngsok Kim, Rachata Ausavarungnirun, Eric Shiu, Rahul Thakur, Daehyun Kim, Aki Kuusela, Allan Knies, Parthasarathy Ranganathan, and Onur Mutlu, "Google Workloads for Consumer Devices: Mitigating Data Movement Bottlenecks" Proceedings of the <u>23rd International Conference on Architectural Support for Programming</u> <u>Languages and Operating Systems</u> (ASPLOS), Williamsburg, VA, USA, March 2018.

62.7% of the total system energy is spent on data movement

Google Workloads for Consumer Devices: Mitigating Data Movement Bottlenecks

Amirali Boroumand¹Saugata Ghose¹Youngsok Kim²Rachata Ausavarungnirun¹Eric Shiu³Rahul Thakur³Daehyun Kim^{4,3}Aki Kuusela³Allan Knies³Parthasarathy Ranganathan³Onur Mutlu^{5,1}42

Major Trends Affecting Main Memory (IV)

Need for main memory capacity, bandwidth, QoS increasing

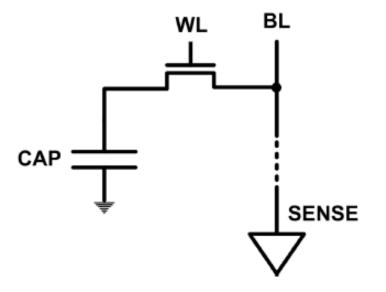
Main memory energy/power is a key system design concern

DRAM technology scaling is ending

- ITRS projects DRAM will not scale easily below X nm
- Scaling has provided many benefits:
 - higher capacity (density), lower cost, lower energy

The DRAM Scaling Problem

- DRAM stores charge in a capacitor (charge-based memory)
 - Capacitor must be large enough for reliable sensing
 - Access transistor should be large enough for low leakage and high retention time
 - □ Scaling beyond 40-35nm (2013) is challenging [ITRS, 2009]

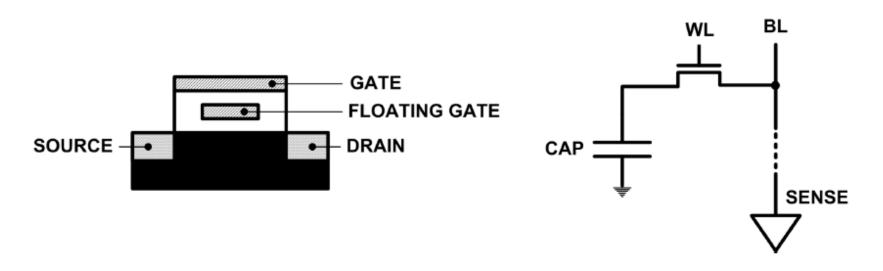


DRAM capacity, cost, and energy/power hard to scale

SAFARI

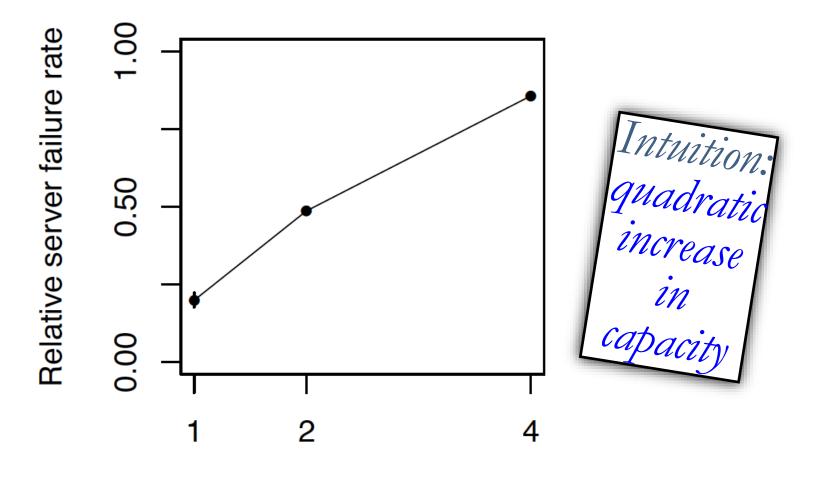
Limits of Charge Memory

- Difficult charge placement and control
 - Flash: floating gate charge
 - DRAM: capacitor charge, transistor leakage
- Data retention and reliable sensing becomes difficult as charge storage unit size reduces



As Memory Scales, It Becomes Unreliable

- Data from all of Facebook's servers worldwide
- Meza+, "Revisiting Memory Errors in Large-Scale Production Data Centers," DSN'15.



Chip density (Gb)

Large-Scale Failure Analysis of DRAM Chips

- Analysis and modeling of memory errors found in all of Facebook's server fleet
- Justin Meza, Qiang Wu, Sanjeev Kumar, and Onur Mutlu, <u>"Revisiting Memory Errors in Large-Scale Production Data</u> <u>Centers: Analysis and Modeling of New Trends from the Field"</u> *Proceedings of the <u>45th Annual IEEE/IFIP International Conference on</u> <u>Dependable Systems and Networks</u> (DSN), Rio de Janeiro, Brazil, June 2015. [Slides (pptx) (pdf)] [DRAM Error Model]*

Revisiting Memory Errors in Large-Scale Production Data Centers: Analysis and Modeling of New Trends from the Field

Justin Meza Qiang Wu* Sanjeev Kumar* Onur Mutlu

Carnegie Mellon University * Facebook, Inc.

Infrastructures to Understand Such Issues

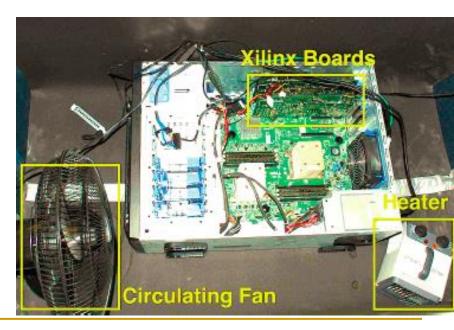


<u>Flipping Bits in Memory Without Accessing</u> <u>Them: An Experimental Study of DRAM</u> <u>Disturbance Errors</u> (Kim et al., ISCA 2014)

Adaptive-Latency DRAM: Optimizing DRAM Timing for the Common-Case (Lee et al., HPCA 2015)

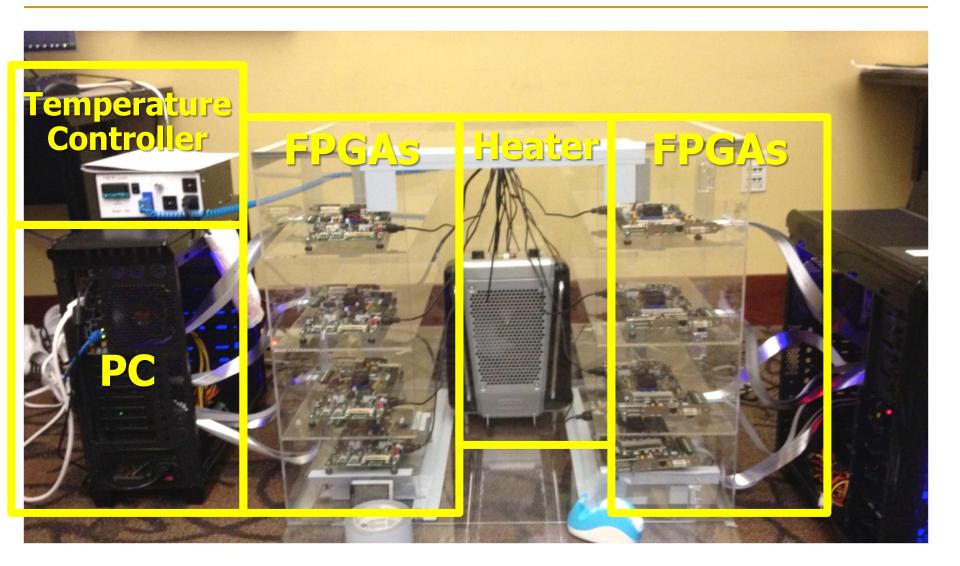
AVATAR: A Variable-Retention-Time (VRT) Aware Refresh for DRAM Systems (Qureshi et al., DSN 2015) An Experimental Study of Data Retention Behavior in Modern DRAM Devices: Implications for Retention Time Profiling Mechanisms (Liu et al., ISCA 2013)

The Efficacy of Error Mitigation Techniques for DRAM Retention Failures: A Comparative Experimental Study (Khan et al., SIGMETRICS 2014)



SAFARI

Infrastructures to Understand Such Issues



SAFARI

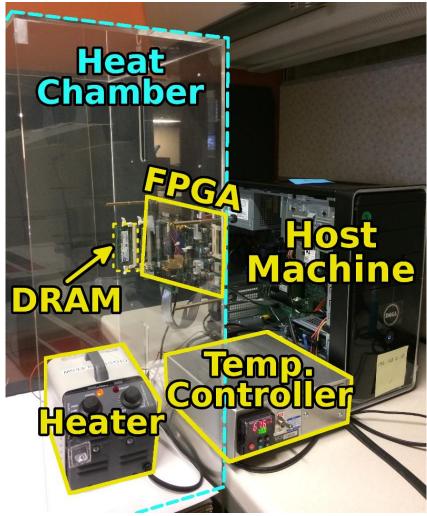
Kim+, "Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors," ISCA 2014.

SoftMC: Open Source DRAM Infrastructure

 Hasan Hassan et al., "<u>SoftMC: A</u> <u>Flexible and Practical Open-</u> <u>Source Infrastructure for</u> <u>Enabling Experimental DRAM</u> <u>Studies</u>," HPCA 2017.

- Flexible
- Easy to Use (C++ API)
- Open-source

github.com/CMU-SAFARI/SoftMC



SoftMC: Open Source DRAM Infrastructure

<u>https://github.com/CMU-SAFARI/SoftMC</u>

SoftMC: A Flexible and Practical Open-Source Infrastructure for Enabling Experimental DRAM Studies

Hasan Hassan^{1,2,3} Nandita Vijaykumar³ Samira Khan^{4,3} Saugata Ghose³ Kevin Chang³ Gennady Pekhimenko^{5,3} Donghyuk Lee^{6,3} Oguz Ergin² Onur Mutlu^{1,3}

¹ETH Zürich ²TOBB University of Economics & Technology ³Carnegie Mellon University ⁴University of Virginia ⁵Microsoft Research ⁶NVIDIA Research A Curious Discovery [Kim et al., ISCA 2014]

One can predictably induce errors in most DRAM memory chips

A simple hardware failure mechanism can create a widespread system security vulnerability



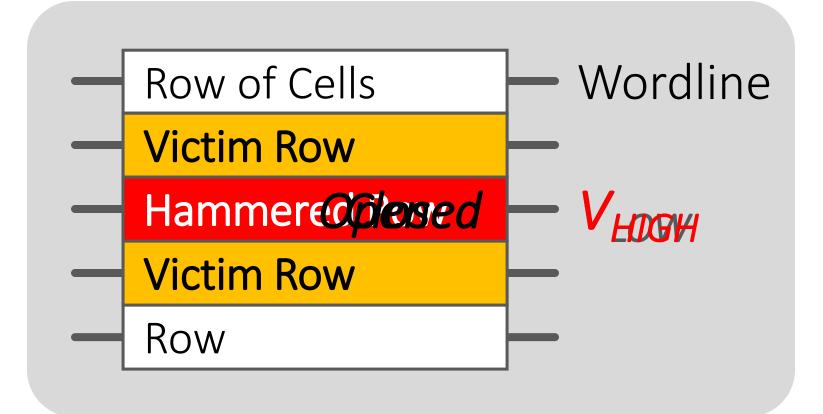
First RowHammer Analysis

Yoongu Kim, Ross Daly, Jeremie Kim, Chris Fallin, Ji Hye Lee, Donghyuk Lee, Chris Wilkerson, Konrad Lai, and Onur Mutlu,
 "Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors"
 Proceedings of the <u>41st International Symposium on Computer Architecture</u> (ISCA), Minneapolis, MN, June 2014.
 [Slides (pptx) (pdf)] [Lightning Session Slides (pptx) (pdf)] [Source Code and Data] [Lecture Video (1 hr 49 mins), 25 September 2020]
 One of the 7 papers of 2012-2017 selected as Top Picks in Hardware and Embedded Security for IEEE TCAD (link).

Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors

Yoongu Kim¹ Ross Daly^{*} Jeremie Kim¹ Chris Fallin^{*} Ji Hye Lee¹ Donghyuk Lee¹ Chris Wilkerson² Konrad Lai Onur Mutlu¹ ¹Carnegie Mellon University ²Intel Labs

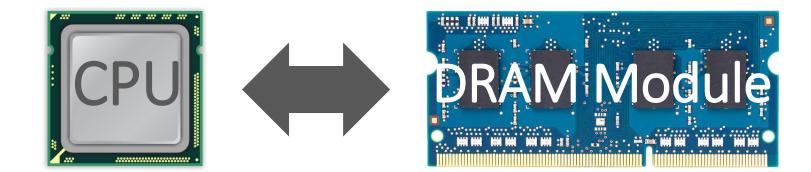
Modern DRAM is Prone to Disturbance Errors



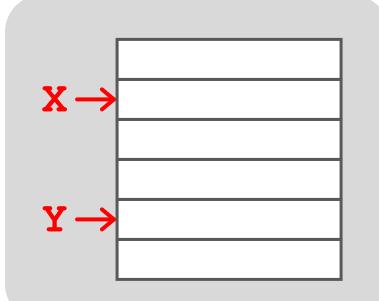
Repeatedly reading a row enough times (before memory gets refreshed) induces disturbance errors in adjacent rows in most real DRAM chips you can buy today

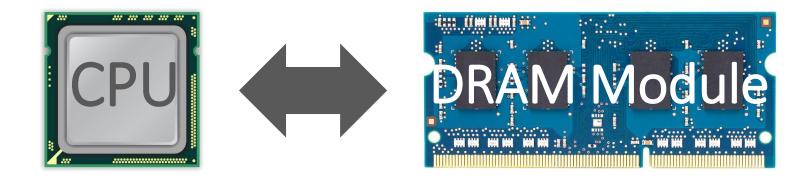
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Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors, (Kim et al., ISCA 2014)

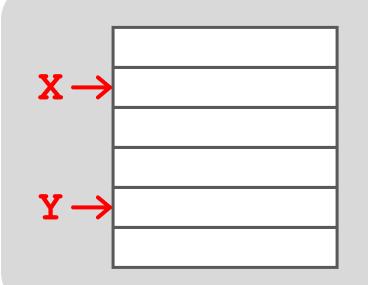


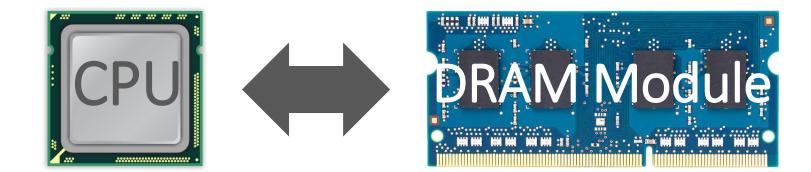
loop: mov (X), %eax mov (Y), %ebx clflush (X) clflush (Y) mfence jmp loop



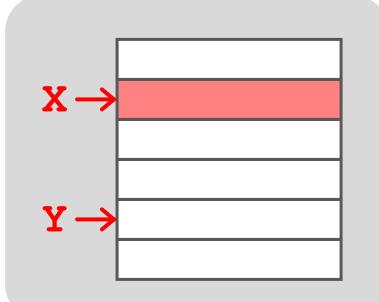


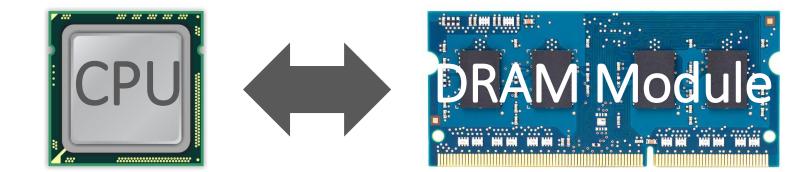
- Avoid *cache hits* Flush X from cache
- Avoid *row hits* to X
 Read Y in another row



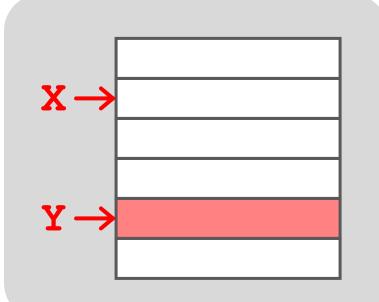


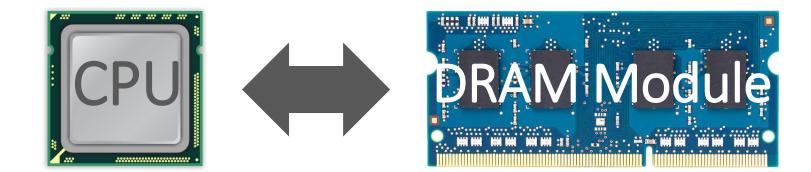
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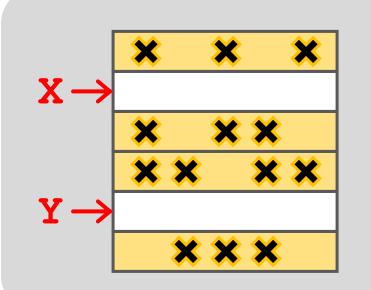


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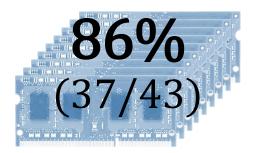


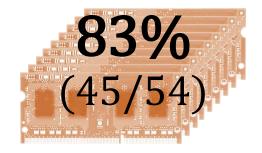
loop: mov (X), %eax mov (Y), %ebx clflush (X) clflush (Y) mfence jmp loop

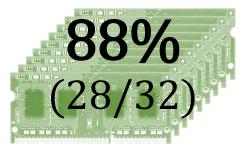


Most DRAM Modules Are Vulnerable





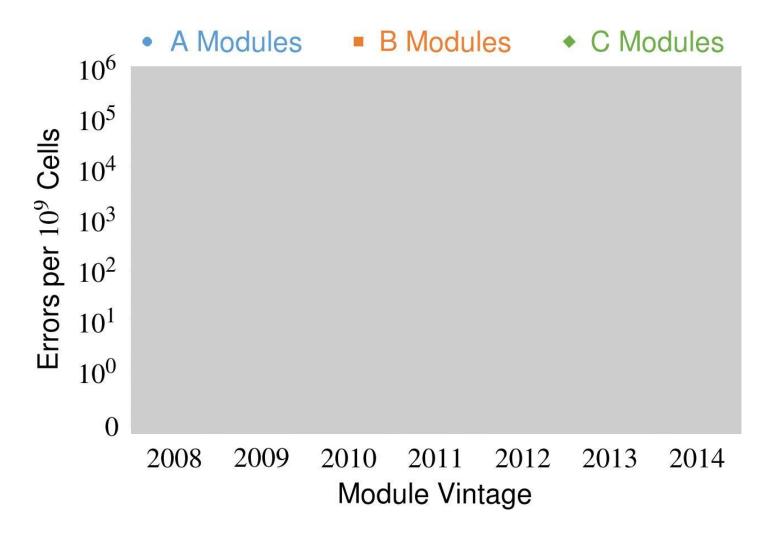




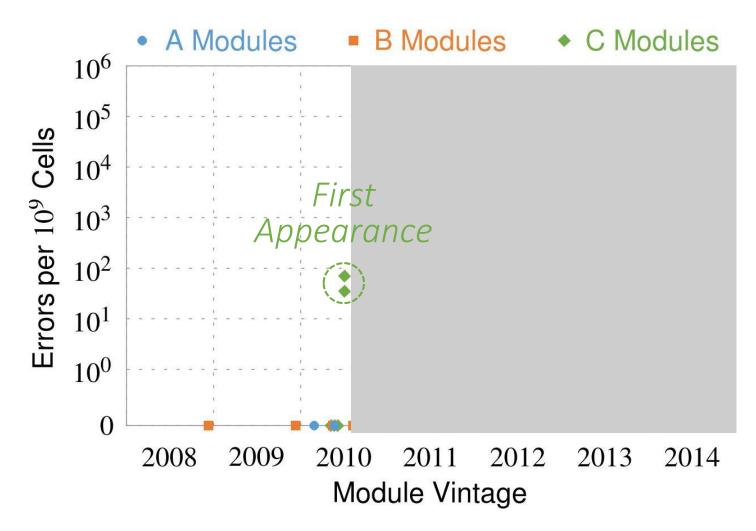
Up to 1.0×10⁷	Up to 2.7×10⁶	Up to 3.3×10⁵

<u>Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM</u> <u>Disturbance Errors</u>, (Kim et al., ISCA 2014)

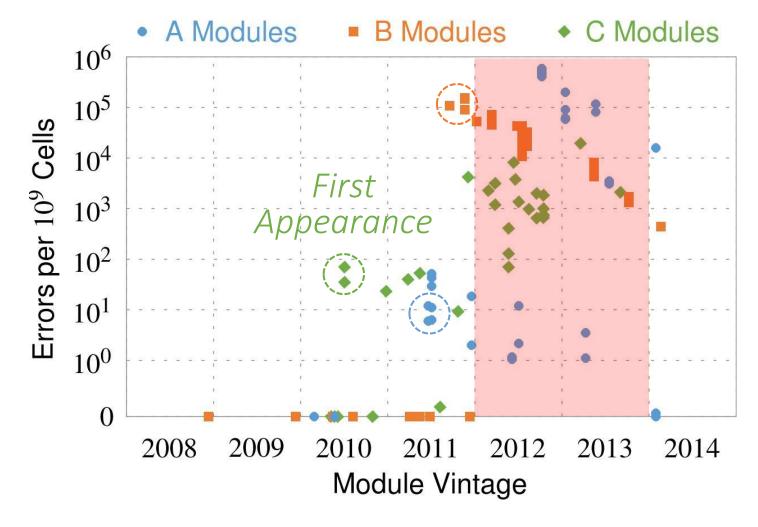
Recent DRAM Is More Vulnerable



Recent DRAM Is More Vulnerable



Recent DRAM Is More Vulnerable



All modules from 2012–2013 are vulnerable

The Reliability & Security Perspectives

Onur Mutlu, "The RowHammer Problem and Other Issues We May Face as Memory Becomes Denser" Invited Paper in Proceedings of the Design, Automation, and Test in Europe Conference (DATE), Lausanne, Switzerland, March 2017. [Slides (pptx) (pdf)]

The RowHammer Problem and Other Issues We May Face as Memory Becomes Denser

Onur Mutlu ETH Zürich onur.mutlu@inf.ethz.ch https://people.inf.ethz.ch/omutlu

SAFARI https://people.inf.ethz.ch/omutlu/pub/rowhammer-and-other-memory-issues_date17.pdf 65

First RowHammer Analysis

Yoongu Kim, Ross Daly, Jeremie Kim, Chris Fallin, Ji Hye Lee, Donghyuk Lee, Chris Wilkerson, Konrad Lai, and Onur Mutlu,
 "Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors"
 Proceedings of the <u>41st International Symposium on Computer Architecture</u> (ISCA), Minneapolis, MN, June 2014.
 [Slides (pptx) (pdf)] [Lightning Session Slides (pptx) (pdf)] [Source Code and Data] [Lecture Video (1 hr 49 mins), 25 September 2020]
 One of the 7 papers of 2012-2017 selected as Top Picks in Hardware and Embedded Security for IEEE TCAD (link).

Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors

Yoongu Kim¹ Ross Daly^{*} Jeremie Kim¹ Chris Fallin^{*} Ji Hye Lee¹ Donghyuk Lee¹ Chris Wilkerson² Konrad Lai Onur Mutlu¹ ¹Carnegie Mellon University ²Intel Labs

RowHammer: 2019 and Beyond...

Onur Mutlu and Jeremie Kim,
 "RowHammer: A Retrospective"
 IEEE Transactions on Computer-Aided Design of Integrated Circuits and Systems (TCAD) Special Issue on Top Picks in Hardware and Embedded Security, 2019.
 [Preliminary arXiv version]
 [Slides from COSADE 2019 (pptx)]
 [Slides from VLSI-SOC 2020 (pptx) (pdf)]
 [Talk Video (1 hr 15 minutes, with Q&A)]

RowHammer: A Retrospective

Onur Mutlu§‡Jeremie S. Kim‡§§ETH Zürich‡Carnegie Mellon University

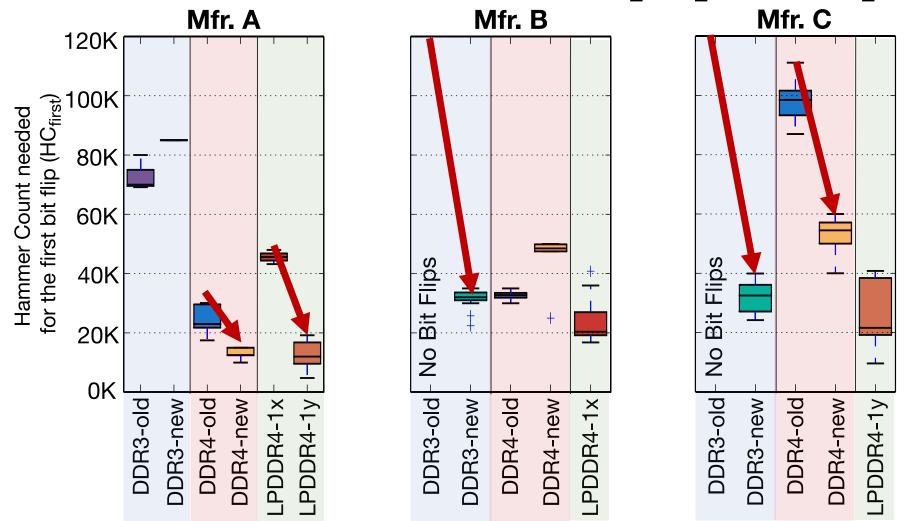
RowHammer in 2020 (I)

 Jeremie S. Kim, Minesh Patel, A. Giray Yaglikci, Hasan Hassan, Roknoddin Azizi, Lois Orosa, and Onur Mutlu,
 "Revisiting RowHammer: An Experimental Analysis of Modern Devices and Mitigation Techniques"
 Proceedings of the <u>47th International Symposium on Computer</u> <u>Architecture</u> (ISCA), Valencia, Spain, June 2020.
 [Slides (pptx) (pdf)]
 [Lightning Talk Slides (pptx) (pdf)]
 [Talk Video (20 minutes)]
 [Lightning Talk Video (3 minutes)]

Revisiting RowHammer: An Experimental Analysis of Modern DRAM Devices and Mitigation Techniques

Jeremie S. Kim^{§†} Minesh Patel[§] A. Giray Yağlıkçı[§] Hasan Hassan[§] Roknoddin Azizi[§] Lois Orosa[§] Onur Mutlu^{§†} [§]ETH Zürich [†]Carnegie Mellon University

5. First RowHammer Bit Flips per Chip



Newer chips from a given DRAM manufacturer **more** vulnerable to RowHammer

5. First RowHammer Bit Flips per Chip



In a DRAM type, HC_{first} reduces significantly from old to new chips, i.e., DDR3: 69.2k to 22.4k, DDR4: 17.5k to 10k, LPDDR4: 16.8k to 4.8k



There are chips whose weakest cells fail after only 4800 hammers

DDF

Newer chips from a given DRAM manufacturer **more** vulnerable to RowHammer

DF

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RowHammer in 2020 (II)

 Lucian Cojocar, Jeremie Kim, Minesh Patel, Lillian Tsai, Stefan Saroiu, Alec Wolman, and Onur Mutlu,
 "Are We Susceptible to Rowhammer? An End-to-End Methodology for Cloud Providers"
 Proceedings of the <u>41st IEEE Symposium on Security and</u> Privacy (S&P), San Francisco, CA, USA, May 2020.
 [Slides (pptx) (pdf)]
 [Talk Video (17 minutes)]

Are We Susceptible to Rowhammer? An End-to-End Methodology for Cloud Providers

Lucian Cojocar, Jeremie Kim^{§†}, Minesh Patel[§], Lillian Tsai[‡], Stefan Saroiu, Alec Wolman, and Onur Mutlu^{§†} Microsoft Research, [§]ETH Zürich, [†]CMU, [‡]MIT

RowHammer in 2020 (III)

 Pietro Frigo, Emanuele Vannacci, Hasan Hassan, Victor van der Veen, Onur Mutlu, Cristiano Giuffrida, Herbert Bos, and Kaveh Razavi,
 "TRRespass: Exploiting the Many Sides of Target Row Refresh" Proceedings of the <u>41st IEEE Symposium on Security and Privacy</u> (S&P), San Francisco, CA, USA, May 2020.
 [Slides (pptx) (pdf)]
 [Lecture Slides (pptx) (pdf)]
 [Lecture Slides (pptx) (pdf)]
 [Lecture Video (17 minutes)]
 [Source Code]
 [Web Article]
 Best paper award.

Pwnie Award 2020 for Most Innovative Research. Pwnie Awards 2020

TRRespass: Exploiting the Many Sides of Target Row Refresh

Pietro Frigo^{*†} Emanuele Vannacci^{*†} Hasan Hassan[§] Victor van der Veen[¶] Onur Mutlu[§] Cristiano Giuffrida^{*} Herbert Bos^{*} Kaveh Razavi^{*}

*Vrije Universiteit Amsterdam

[§]ETH Zürich

[¶]Qualcomm Technologies Inc.

Two RowHammer Papers at MICRO 2021

 Lois Orosa, Abdullah Giray Yaglikci, Haocong Luo, Ataberk Olgun, Jisung Park, Hasan Hassan, Minesh Patel, Jeremie S. Kim, and <u>Onur Mutlu</u>, "A Deeper Look into RowHammer's Sensitivities: Experimental Analysis of Real DRAM Chips and Implications on Future Attacks and Defenses" *Proceedings of the 54th International Symposium on Microarchitecture (MICRO)*, Virtual, October 2021. [Slides (pptx) (pdf)]
 [Short Talk Slides (pptx) (pdf)]
 [Lightning Talk Slides (pptx) (pdf)]
 [Lightning Talk Video (1.5 minutes)]
 [arXiv version]

A Deeper Look into RowHammer's Sensitivities: Experimental Analysis of Real DRAM Chips and Implications on Future Attacks and Defenses

Lois Orosa* ETH Zürich A. Giray Yağlıkçı*

Haocong Luo ETH Zürich Ataberk Olgun ETH Zürich, TOBB ETÜ Jisung Park ETH Zürich

Hasan Hassan ETH Zürich Minesh Patel ETH Zürich Jeremie S. Kim ETH Zürich

Onur Mutlu ETH Zürich

Two RowHammer Papers at MICRO 2021

 Hasan Hassan, Yahya Can Tugrul, Jeremie S. Kim, Victor van der Veen, Kaveh Razavi, and <u>Onur Mutlu</u>, <u>"Uncovering In-DRAM RowHammer Protection Mechanisms: A</u> <u>New Methodology, Custom RowHammer Patterns, and</u> <u>Implications"</u> *Proceedings of the <u>54th International Symposium on Microarchitecture</u> (<i>MICRO*), Virtual, October 2021.

[Slides (pptx) (pdf)] [Short Talk Slides (pptx) (pdf)] [Lightning Talk Slides (pptx) (pdf)] [Talk Video (25 minutes)] [Lightning Talk Video (100 seconds)] [arXiv version]

Uncovering In-DRAM RowHammer Protection Mechanisms: A New Methodology, Custom RowHammer Patterns, and Implications

Hasan Hassan †	Yahya Can Tuğrul ^{†‡}	Jeremie S. Kir	\mathbf{n}^{\dagger} Victor van der Veen ^{σ}
	Kaveh Razavi †	Onur Mutlu	†
†ETH Zürich	[‡] TOBB University of Economics & Technology		$^{\sigma}$ Qualcomm Technologies Inc.

A New RowHammer Paper at DSN 2022

 A. Giray Yağlıkçı, Haocong Luo, Geraldo F. de Oliviera, Ataberk Olgun, Minesh Patel, Jisung Park, Hasan Hassan, Jeremie S. Kim, Lois Orosa, and <u>Onur Mutlu</u>, <u>"Understanding RowHammer Under Reduced Wordline Voltage: An</u> <u>Experimental Study Using Real DRAM Devices"</u> *Proceedings of the <u>52nd Annual IEEE/IFIP International Conference on</u> <u>Dependable Systems and Networks</u> (DSN), Baltimore, MD, USA, June 2022. [Slides (pptx) (pdf)]
 [Lightning Talk Slides (pptx) (pdf)]
 [Talk Video (34 minutes, including Q&A)]
 [Lightning Talk Video (2 minutes)]*

Understanding RowHammer Under Reduced Wordline Voltage: An Experimental Study Using Real DRAM Devices

A. Giray Yağlıkçı¹ Haocong Luo¹ Geraldo F. de Oliviera¹ Ataberk Olgun¹ Minesh Patel¹ Jisung Park¹ Hasan Hassan¹ Jeremie S. Kim¹ Lois Orosa^{1,2} Onur Mutlu¹ ¹ETH Zürich ²Galicia Supercomputing Center (CESGA)



RowHammer is still an open problem

Security by obscurity is likely not a good solution

Major Trends Affecting Main Memory (V)

- DRAM scaling has already become increasingly difficult
 - Increasing cell leakage current, reduced cell reliability, increasing manufacturing difficulties [Kim+ ISCA 2014], [Liu+ ISCA 2013], [Mutlu IMW 2013], [Mutlu DATE 2017]
 - Difficult to significantly improve capacity, energy

Emerging memory technologies are promising

-	

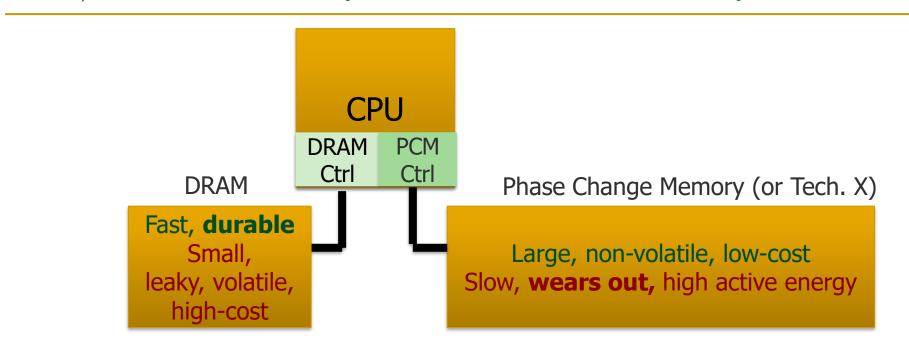
Major Trends Affecting Main Memory (V)

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 - Difficult to significantly improve capacity, energy

Emerging memory technologies are promising

3D-Stacked DRAM	higher bandwidth	smaller capacity
Reduced-Latency DRAM (e.g., RL/TL-DRAM, FLY-RAM)	lower latency	higher cost
Low-Power DRAM (e.g., LPDDR3, LPDDR4, Voltron)	lower power	higher latency higher cost
Non-Volatile Memory (NVM) (e.g., PCM, STTRAM, ReRAM, 3D Xpoint)	larger capacity	higher latency higher dynamic power lower endurance

Major Trend: Hybrid Main Memory



Hardware/software manage data allocation and movement to achieve the best of multiple technologies

Meza+, "Enabling Efficient and Scalable Hybrid Memories," IEEE Comp. Arch. Letters, 2012. Yoon+, "Row Buffer Locality Aware Caching Policies for Hybrid Memories," ICCD 2012 Best Paper Award.



Main Memory Needs Intelligent Controllers

Industry Is Writing Papers About It, Too

DRAM Process Scaling Challenges

* Refresh

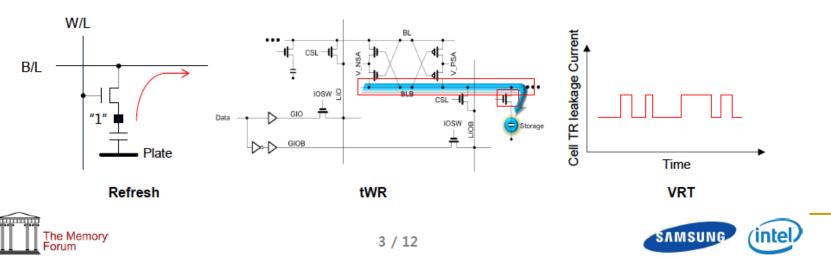
- · Difficult to build high-aspect ratio cell capacitors decreasing cell capacitance
- · Leakage current of cell access transistors increasing

✤ tWR

- · Contact resistance between the cell capacitor and access transistor increasing
- · On-current of the cell access transistor decreasing
- · Bit-line resistance increasing

VRT

Occurring more frequently with cell capacitance decreasing



Call for Intelligent Memory Controllers

DRAM Process Scaling Challenges

* Refresh

Difficult to build high-aspect ratio cell capacitors decreasing cell capacitance
THE MEMORY FORUM 2014

Co-Architecting Controllers and DRAM to Enhance DRAM Process Scaling

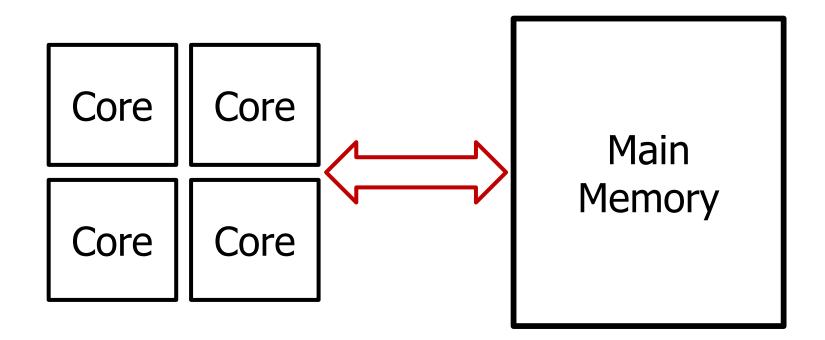
Uksong Kang, Hak-soo Yu, Churoo Park, *Hongzhong Zheng, **John Halbert, **Kuljit Bains, SeongJin Jang, and Joo Sun Choi



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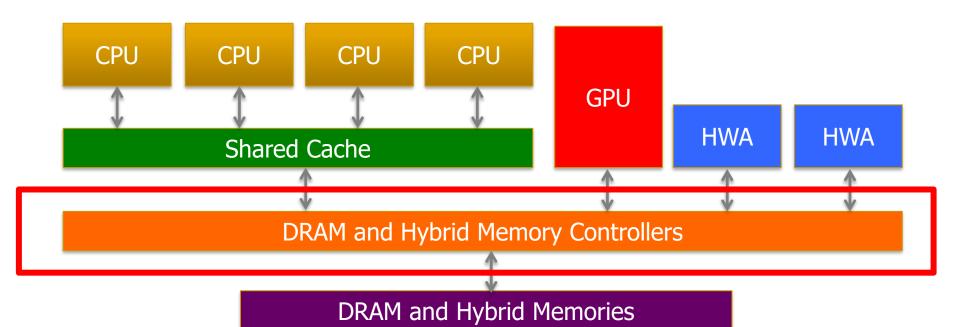
Samsung Electronics, Hwasung, Korea / *Samsung Electronics, San Jose / **Intel

An Orthogonal Issue: Memory Interference



Cores' interfere with each other when accessing shared main memory Uncontrolled interference leads to many problems (QoS, performance)

Goal: Predictable Performance in Complex Systems



- Heterogeneous agents: CPUs, GPUs, and HWAs
- Main memory interference between CPUs, GPUs, HWAs

How to allocate resources to heterogeneous agents to mitigate interference and provide predictable performance?



Main Memory Needs Intelligent Controllers

Solving the Memory Problem

How Do We Solve The Memory Problem?

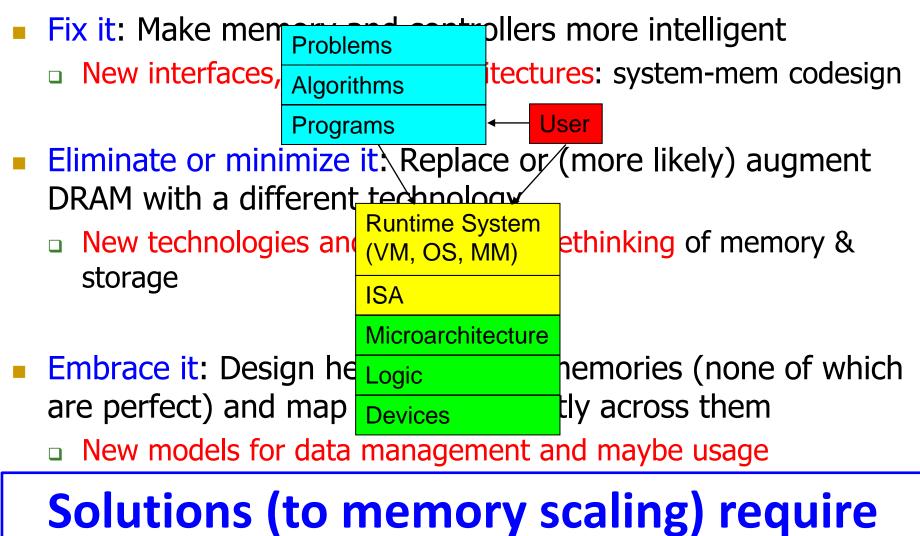
- Fix it: Make memory and controllers more intelligent
 New interfaces, functions, architectures: system-mem codesign
- Eliminate or minimize it: Replace or (more likely) augment DRAM with a different technology
 - New technologies and system-wide rethinking of memory & storage
- Embrace it: Design heterogeneous memories (none of which are perfect) and map data intelligently across them
 New models for data management and maybe usage

How Do We Solve The Memory Problem?

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Solutions (to memory scaling) require software/hardware/device cooperation

How Do We Solve The Memory Problem?



software/hardware/device cooperation

Solution 1: New Memory Architectures

- Overcome memory shortcomings with
 - Memory-centric system design
 - Novel memory architectures, interfaces, functions
 - Better waste management (efficient utilization)
- Key issues to tackle
 - Enable reliability at low cost \rightarrow high capacity
 - Reduce energy
 - Reduce latency
 - Improve bandwidth
 - Reduce waste (capacity, bandwidth, latency)
 - Enable computation close to data

Solution 1: New Memory Architectures

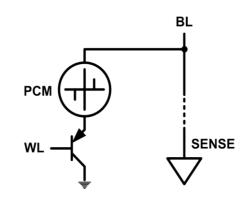
Kim+, "A Case for Exploiting Subarray-Level Parallelism in DRAM," ISCA 2012 Lee+, "Tiered-Latency DRAM: A Low Latency and Low Cost DRAM Architecture," HPCA 2013. Liu+, "An Experimental Study of Data Retention Behavior in Modern DRAM Devices," ISCA 2013 Seshadri+, "RowClone: Fast and Efficient In-DRAM Copy and Initialization of Bulk Data," MICRO 2013 Pethimenko+, "Linearly Compressed Pages: A Main Memory Compression Framework," MICRO 2013. Chang+, "Improving DRAM Performance by Parallelizing Refreshes with Accesses," HPCA 2014. Khan+, "The Efficacy of Error Mitigation Techniques for DRAM Retention Failures: A Comparative Experimental Study," SIGMETRICS 2014 Luo+, "Characterizing Application Memory Error Vulnerability to Optimize Data Center Cost." DSN 2014. Kim+, "Flipping Bits in Memory Without Accessing Them: An Experimental Study of DRAM Disturbance Errors," ISCA 2014 Lee+, "Adaptive-Latency DRAM: Optimizing DRAM Timing for the Common-Case," HPCA 2015. Qureshi+, "AVATAR: A Variable-Retention-Time (VRT) Aware Refresh for DRAM Systems," DSN 2015. Meza+, "Revisiting Memory Errors in Large-Scale Production Data Centers: Analysis and Modeling of New Trends from the Field," DSN 2015. Kim+, "Ramulator: A Fast and Extensible DRAM Simulator," IEEE CAL 2015. Seshadri+, "Fast Bulk Bitwise AND and OR in DRAM," IEEE CAL 2015. Ahn+, "A Scalable Processing-in-Memory Accelerator for Parallel Graph Processing," ISCA 2015. Ahn+, "PIM-Enabled Instructions: A Low-Overhead, Locality-Aware Processing-in-Memory Architecture," ISCA 2015 Lee+, "Decoupled Direct Memory Access: Isolating CPU and IO Traffic by Leveraging a Dual-Data-Port DRAM," PACT 2015. Seshartri+, "Gather-Scatter DRAM: In-DRAM Address Translation to Improve the Spatial Locality of Non-unit Strided Accesses," MICRO 2015. Lee+, "Simultaneous Multi-Layer Access: Improving 3D-Stacked Memory Bandwidth at Low Cost," TACO 2016. 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Hashemi+, "Continuous Runahead: Transparent Hardware Acceleration for Memory Intensive Workloads," MICRO 2016 Khan+, "A Case for Memory Content-Based Detection and Mitigation of Data-Dependent Failures in DRAM"," IEEE CAL 2016. Hassan+, "SoftMC: A Flexible and Practical Open-Source Infrastructure for Enabling Experimental DRAM Studies," HPCA 2017. Mutlu, "The RowHammer Problem and Other Issues We May Face as Memory Becomes Denser," DATE 2017. Lee+, "Design-Induced Latency Variation in Modern DRAM Chips: Characterization, Analysis, and Latency Reduction Mechanisms," SIGMETRICS 2017. Chang+, "Understanding Reduced-Voltage Operation in Modern DRAM Devices: Experimental Characterization, Analysis, and Mechanisms," SIGMETRICS 2017. Patel+, "The Reach Profiler (REAPER): Enabling the Mitigation of DRAM Retention Failures via Profiling at Aggressive Conditions," ISCA 2017. Seshadri and Mutlu, "Simple Operations in Memory to Reduce Data Movement," ADCOM 2017. 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Liu+, "RAIDR: Retention-Aware Intelligent DRAM Refresh," ISCA 2012

Vijaykumar+, *A Case for Core-Assisted Bottleneck Acceleration in GPUs: Enabling Flexible Data Compression with Assi Pekhimenko+, *Toggle-Aware Bandwidth Compression for GPUs,* HPCA 2016.

Solution 2: Emerging Memory Technologies

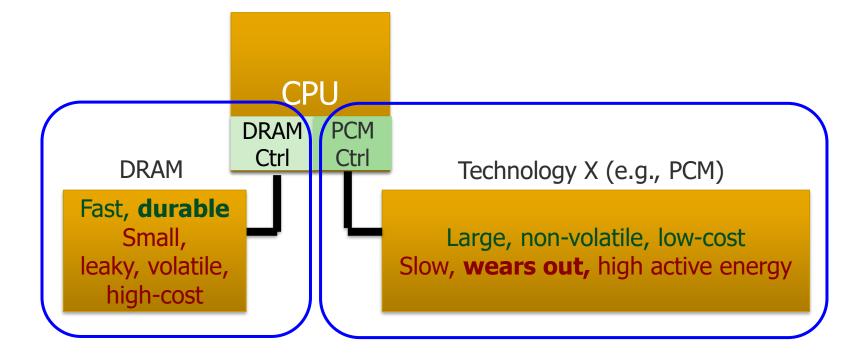
- Some emerging resistive memory technologies seem more scalable than DRAM (and they are non-volatile)
- Example: Phase Change Memory
 - Data stored by changing phase of material
 - Data read by detecting material's resistance
 - Expected to scale to 9nm (2022 [ITRS 2009])
 - Prototyped at 20nm (Raoux+, IBM JRD 2008)
 - Expected to be denser than DRAM: can store multiple bits/cell
- But, emerging technologies have (many) shortcomings
 Can they be enabled to replace/augment/surpass DRAM?



Solution 2: Emerging Memory Technologies

- Lee+, "Architecting Phase Change Memory as a Scalable DRAM Alternative," ISCA'09, CACM'10, IEEE Micro'10.
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- Chauhan+, "NVMove: Helping Programmers Move to Byte-Based Persistence," INFLOW 2016.
- Li+, "Utility-Based Hybrid Memory Management," CLUSTER 2017.
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- Wang+, "Panthera: Holistic Memory Management for Big Data Processing over Hybrid Memories," PLDI 2019.
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- Song+, "Improving Phase Change Memory Performance with Data Content Aware Access," ISMM 2020.
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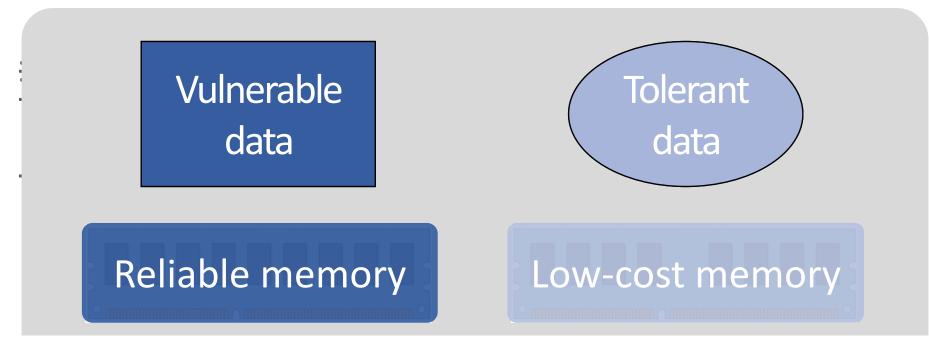
Combination: Hybrid Memory Systems



Hardware/software manage data allocation and movement to achieve the best of multiple technologies

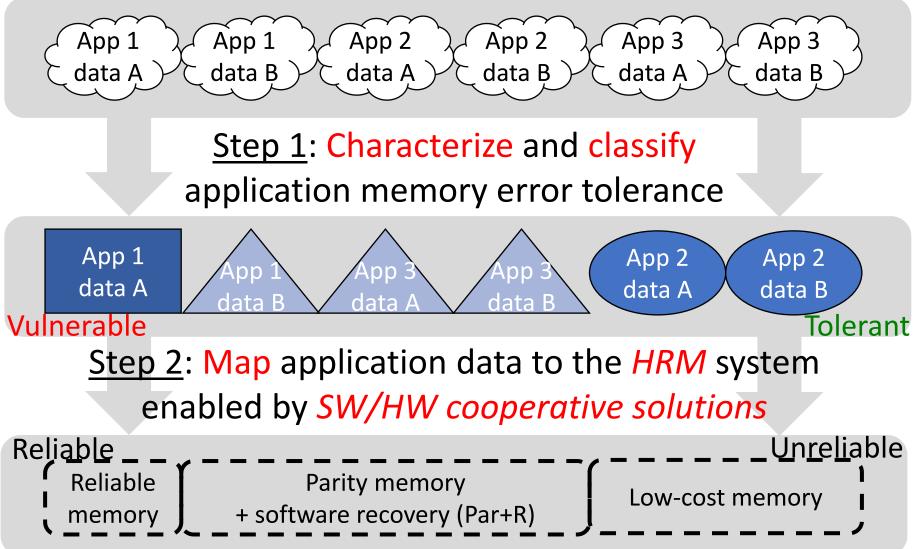
Meza+, "Enabling Efficient and Scalable Hybrid Memories," IEEE Comp. Arch. Letters, 2012. Yoon, Meza et al., "Row Buffer Locality Aware Caching Policies for Hybrid Memories," ICCD 2012 Best Paper Award.

Exploiting Memory Error Tolerance with Hybrid Memory Systems

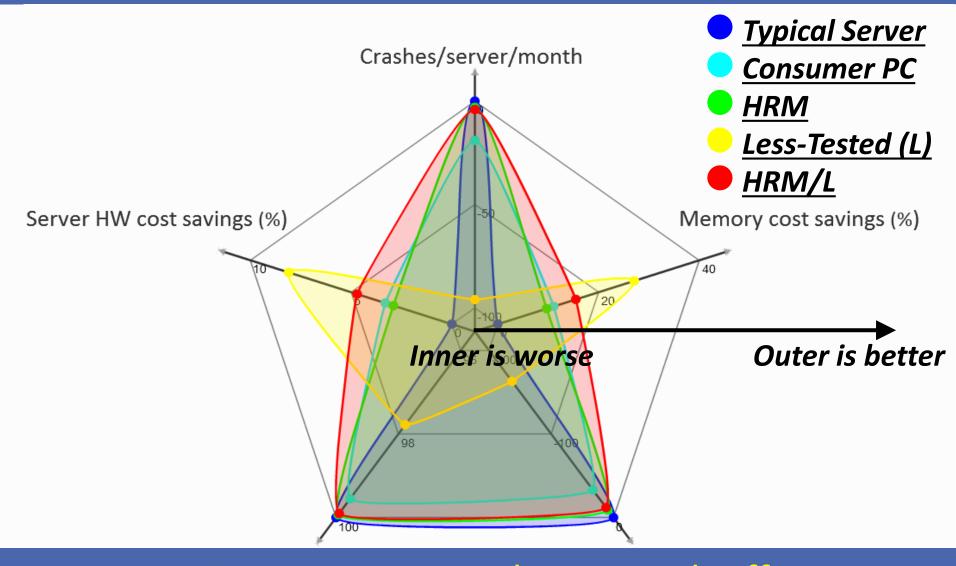


On Microsoft's Web Search workload Reduces server hardware cost by 4.7 % Achieves single server availability target of 99.90 % Heterogeneous-Reliability Memory [DSN 2014]

Heterogeneous-Reliability Memory



Evaluation Results



Bigger area means better tradeoff

More on Heterogeneous Reliability Memory

Yixin Luo, Sriram Govindan, Bikash Sharma, Mark Santaniello, Justin Meza, Aman Kansal, Jie Liu, Badriddine Khessib, Kushagra Vaid, and Onur Mutlu,
 "Characterizing Application Memory Error Vulnerability to Optimize
 Data Center Cost via Heterogeneous-Reliability Memory"
 Proceedings of the <u>44th Annual IEEE/IFIP International Conference on</u>
 Dependable Systems and Networks (DSN), Atlanta, GA, June 2014. [Summary]
 [Slides (pptx) (pdf)] [Coverage on ZDNet]

Characterizing Application Memory Error Vulnerability to Optimize Datacenter Cost via Heterogeneous-Reliability Memory

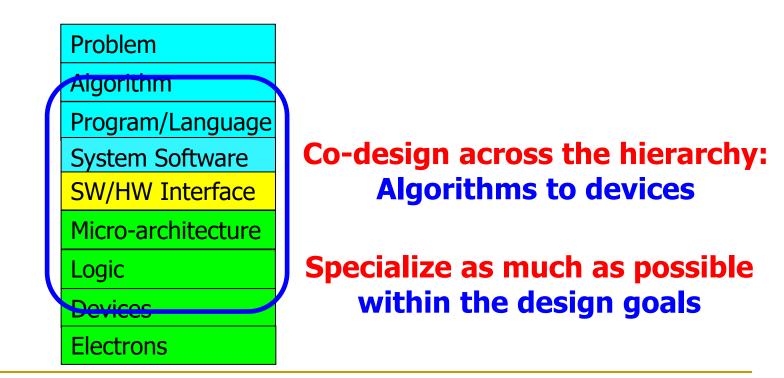
Yixin Luo Sriram Govindan^{*} Bikash Sharma^{*} Mark Santaniello^{*} Justin Meza Aman Kansal^{*} Jie Liu^{*} Badriddine Khessib^{*} Kushagra Vaid^{*} Onur Mutlu Carnegie Mellon University, yixinluo@cs.cmu.edu, {meza, onur}@cmu.edu *Microsoft Corporation, {srgovin, bsharma, marksan, kansal, jie.liu, bkhessib, kvaid}@microsoft.com

HRM is an Example of Our Axiom

To achieve the highest energy efficiency and performance:

we must take the expanded view

of computer architecture



An Orthogonal Issue: Memory Interference

- Problem: Memory interference between cores is uncontrolled
 - \rightarrow unfairness, starvation, low performance
 - \rightarrow uncontrollable, unpredictable, vulnerable system
- Solution: QoS-Aware Memory Systems
 - Hardware designed to provide a configurable fairness substrate
 - Application-aware memory scheduling, partitioning, throttling
 - Software designed to configure the resources to satisfy different QoS goals
- QoS-aware memory systems can provide predictable performance and higher efficiency

Strong Memory Service Guarantees

- Goal: Satisfy performance/SLA requirements in the presence of shared main memory, heterogeneous agents, and hybrid memory/storage
- Approach:
 - Develop techniques/models to accurately estimate the performance loss of an application/agent in the presence of resource sharing
 - Develop mechanisms (hardware and software) to enable the resource partitioning/prioritization needed to achieve the required performance levels for all applications
 - All them while providing high system performance
- Subramanian et al., "MISE: Providing Performance Predictability and Improving Fairness in Shared Main Memory Systems," HPCA 2013.
- Subramanian et al., "The Application Slowdown Model," MICRO 2015.

MISE: Predictable Performance [HPCA'13]

 Lavanya Subramanian, Vivek Seshadri, Yoongu Kim, Ben Jaiyen, and Onur Mutlu,
 "MISE: Providing Performance Predictability and Improving Fairness in Shared Main Memory Systems"
 Proceedings of the 19th International Symposium on High-Performance Computer Architecture (HPCA), Shenzhen, China, February 2013. Slides (pptx)

MISE: Providing Performance Predictability and Improving Fairness in Shared Main Memory Systems

Lavanya Subramanian Vivek Seshadri

ivek Seshadri Yoongu Kim Ben Jaiyen Onur Mutlu Carnegie Mellon University

ASM: Predictable Performance [MICRO'15]

- Lavanya Subramanian, Vivek Seshadri, Arnab Ghosh, Samira Khan, and Onur Mutlu,
 - "The Application Slowdown Model: Quantifying and Controlling the Impact of Inter-Application Interference at Shared Caches and Main Memory"

Proceedings of the <u>48th International Symposium on Microarchitecture</u> (**MICRO**), Waikiki, Hawaii, USA, December 2015. [<u>Slides (pptx) (pdf)</u>] [<u>Lightning Session Slides (pptx) (pdf)</u>] [<u>Poster</u> (<u>pptx) (pdf)</u>] [<u>Source Code</u>]

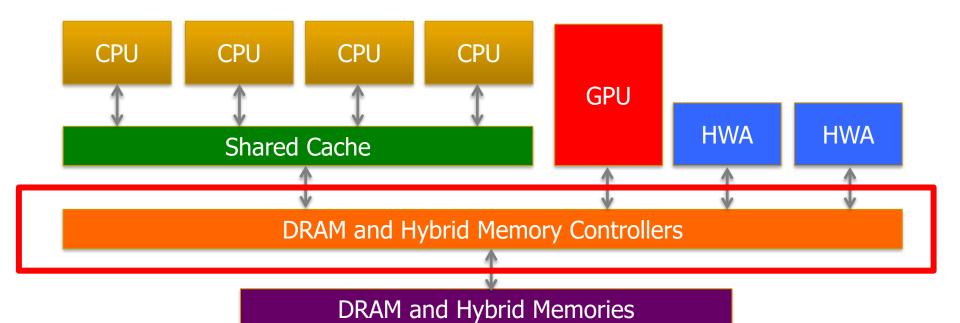
The Application Slowdown Model: Quantifying and Controlling the Impact of Inter-Application Interference at Shared Caches and Main Memory

Lavanya Subramanian^{*}§ Vivek Seshadri^{*} Arnab Ghosh^{*†} Samira Khan^{*‡} Onur Mutlu^{*}

*Carnegie Mellon University §Intel Labs [†]IIT Kanpur [‡]University of Virginia

Memory Controllers

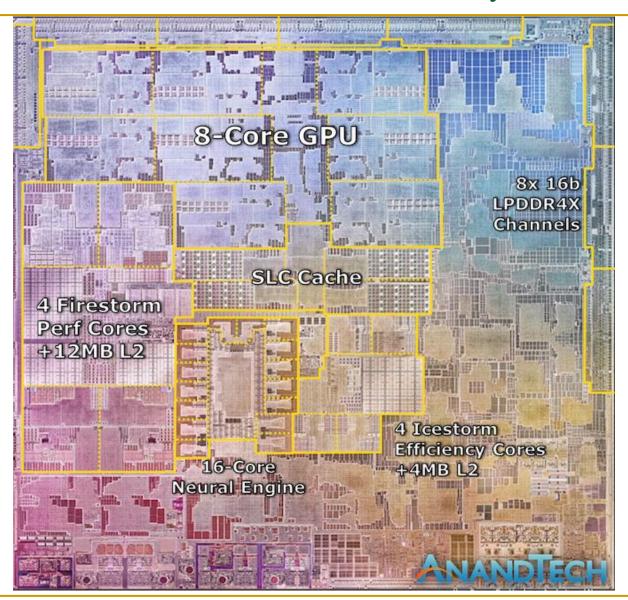
Memory Control is Getting More Complex



- Heterogeneous agents: CPUs, GPUs, and HWAs
- Main memory interference between CPUs, GPUs, HWAs

Many goals, many constraints, many metrics ...

A Similar Picture from Real Systems



Apple M1, 2021

Source: https://www.anandtech.com/show/16252/mac-mini-apple-m1-tested

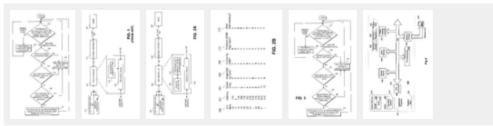
It All Started with FSB Controllers (2001)

Method and apparatus to control memory accesses

Abstract

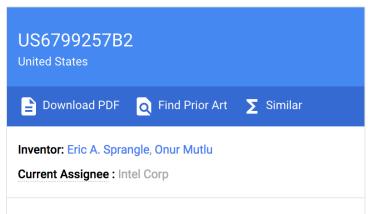
A method and apparatus for accessing memory comprising monitoring memory accesses from a hardware prefetcher and determining whether the memory accesses from the hardware prefetcher are used by an out-of-order core. A front side bus controller switches memory access modes from a minimize memory access latency mode to a maximize memory bus bandwidth mode if a percentage of the memory accesses generated by the hardware prefetcher are used by the out-of-order core.

Images (6)



Classifications

G06F12/0215 Addressing or allocation; Relocation with look ahead addressing means



Worldwide applications

2002 • <u>US</u> 2003 • AU JP DE KR CN WO GB TW 2004 • US 2005 • HK

Application US10/079,967 events ⑦

2002-02-21 • Application filed by Intel Corp

2002-02-21 • Priority to US10/079,967

2002-04-25 • Assigned to INTEL CORPORATION [®]

Memory Performance Attacks [USENIX SEC'07]

 Thomas Moscibroda and Onur Mutlu, <u>"Memory Performance Attacks: Denial of Memory Service</u> <u>in Multi-Core Systems"</u> *Proceedings of the <u>16th USENIX Security Symposium</u> (USENIX SECURITY), pages 257-274, Boston, MA, August 2007. <u>Slides</u> (ppt)*

Memory Performance Attacks: Denial of Memory Service in Multi-Core Systems

Thomas Moscibroda Onur Mutlu Microsoft Research {moscitho,onur}@microsoft.com

STFM [MICRO'07]

 Onur Mutlu and Thomas Moscibroda, "Stall-Time Fair Memory Access Scheduling for Chip Multiprocessors" Proceedings of the <u>40th International Symposium on</u> <u>Microarchitecture</u> (MICRO), pages 146-158, Chicago, IL, December 2007. [Summary] [Slides (ppt)]

Stall-Time Fair Memory Access Scheduling for Chip Multiprocessors

Onur Mutlu Thomas Moscibroda

Microsoft Research {onur,moscitho}@microsoft.com

PAR-BS [ISCA'08]

 Onur Mutlu and Thomas Moscibroda, <u>"Parallelism-Aware Batch Scheduling: Enhancing both</u> <u>Performance and Fairness of Shared DRAM Systems"</u> *Proceedings of the <u>35th International Symposium on Computer</u> <u>Architecture</u> (ISCA), pages 63-74, Beijing, China, June 2008. [Summary] [Slides (ppt)]*

Parallelism-Aware Batch Scheduling: Enhancing both Performance and Fairness of Shared DRAM Systems

Onur Mutlu Thomas Moscibroda Microsoft Research {onur,moscitho}@microsoft.com Variants implemented in Samsung SoC memory controllers

Effective platform level approach and DRAM accesses are crucial to system performance. This paper touches this topics and suggest a superior approach to current known techniques. **Review from ISCA 2008**

ATLAS Memory Scheduler [HPCA'10]

 Yoongu Kim, Dongsu Han, Onur Mutlu, and Mor Harchol-Balter, <u>"ATLAS: A Scalable and High-Performance Scheduling</u> <u>Algorithm for Multiple Memory Controllers"</u> *Proceedings of the <u>16th International Symposium on High-</u> <u>Performance Computer Architecture</u> (HPCA), Bangalore, India, January 2010. <u>Slides (pptx)</u>*

ATLAS: A Scalable and High-Performance Scheduling Algorithm for Multiple Memory Controllers

Yoongu Kim Dongsu Han Onur Mutlu Mor Harchol-Balter Carnegie Mellon University

Thread Cluster Memory Scheduling [MICRO'10]

 Yoongu Kim, Michael Papamichael, Onur Mutlu, and Mor Harchol-Balter,
 "Thread Cluster Memory Scheduling: Exploiting Differences in Memory Access Behavior"
 Proceedings of the <u>43rd International Symposium on</u> Microarchitecture (MICRO), pages 65-76, Atlanta, GA, December 2010. Slides (pptx) (pdf)

Thread Cluster Memory Scheduling: Exploiting Differences in Memory Access Behavior

Yoongu Kim yoonguk@ece.cmu.edu

Michael Papamichael papamix@cs.cmu.edu

Onur Mutlu onur@cmu.edu

Mor Harchol-Balter harchol@cs.cmu.edu

Carnegie Mellon University

BLISS [ICCD'14, TPDS'16]

 Lavanya Subramanian, Donghyuk Lee, Vivek Seshadri, Harsha Rastogi, and Onur Mutlu,
 "The Blacklisting Memory Scheduler: Achieving High Performance and Fairness at Low Cost"
 Proceedings of the 32nd IEEE International Conference on Computer Design (ICCD), Seoul, South Korea, October 2014.
 [Slides (pptx) (pdf)]

The Blacklisting Memory Scheduler: Achieving High Performance and Fairness at Low Cost

Lavanya Subramanian, Donghyuk Lee, Vivek Seshadri, Harsha Rastogi, Onur Mutlu Carnegie Mellon University {lsubrama,donghyu1,visesh,harshar,onur}@cmu.edu

Staged Memory Scheduling: CPU-GPU [ISCA'12]

 Rachata Ausavarungnirun, Kevin Chang, Lavanya Subramanian, Gabriel Loh, and Onur Mutlu,
 "Staged Memory Scheduling: Achieving High Performance and Scalability in Heterogeneous Systems" Proceedings of the <u>39th International Symposium on Computer</u> <u>Architecture</u> (ISCA), Portland, OR, June 2012. <u>Slides (pptx)</u>

Staged Memory Scheduling: Achieving High Performance and Scalability in Heterogeneous Systems

Rachata Ausavarungnirun[†] Kevin Kai-Wei Chang[†] Lavanya Subramanian[†] Gabriel H. Loh[‡] Onur Mutlu[†]

[†]Carnegie Mellon University {rachata,kevincha,lsubrama,onur}@cmu.edu

SAFARI

[‡]Advanced Micro Devices, Inc. gabe.loh@amd.com

DASH: Heterogeneous Systems [TACO'16]

 Hiroyuki Usui, Lavanya Subramanian, Kevin Kai-Wei Chang, and Onur Mutlu,

"DASH: Deadline-Aware High-Performance Memory Scheduler for Heterogeneous Systems with Hardware Accelerators"

ACM Transactions on Architecture and Code Optimization (TACO),

Vol. 12, January 2016. Presented at the <u>11th HiPEAC Conference</u>, Prague, Czech Republic, January 2016. [<u>Slides (pptx) (pdf)</u>] [<u>Source Code</u>]

DASH: Deadline-Aware High-Performance Memory Scheduler for Heterogeneous Systems with Hardware Accelerators

HIROYUKI USUI, LAVANYA SUBRAMANIAN, KEVIN KAI-WEI CHANG, and ONUR MUTLU, Carnegie Mellon University

MISE: Predictable Performance [HPCA'13]

 Lavanya Subramanian, Vivek Seshadri, Yoongu Kim, Ben Jaiyen, and Onur Mutlu,
 "MISE: Providing Performance Predictability and Improving Fairness in Shared Main Memory Systems"
 Proceedings of the 19th International Symposium on High-Performance Computer Architecture (HPCA), Shenzhen, China, February 2013. Slides (pptx)

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The Application Slowdown Model: Quantifying and Controlling the Impact of Inter-Application Interference at Shared Caches and Main Memory

Lavanya Subramanian^{*}§ Vivek Seshadri^{*} Arnab Ghosh^{*†} Samira Khan^{*‡} Onur Mutlu^{*}

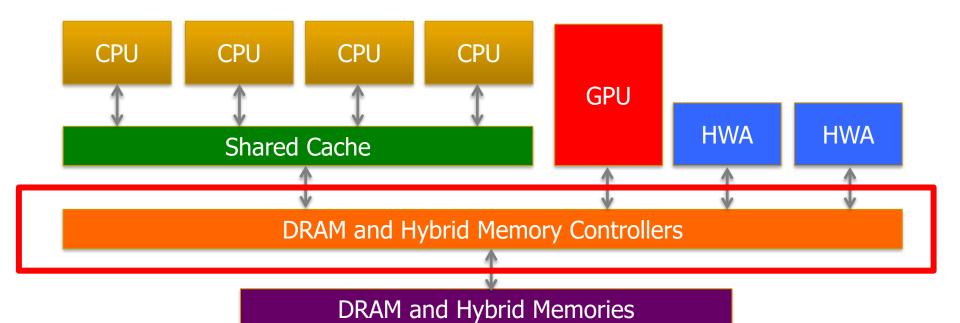
*Carnegie Mellon University §Intel Labs [†]IIT Kanpur [‡]University of Virginia



Memory Controllers are critical to research

They will become even more important

Memory Control is Getting More Complex



- Heterogeneous agents: CPUs, GPUs, and HWAs
- Main memory interference between CPUs, GPUs, HWAs

Many goals, many constraints, many metrics ...

Memory Control w/ Machine Learning [ISCA'08]

 Engin Ipek, Onur Mutlu, José F. Martínez, and Rich Caruana, "Self Optimizing Memory Controllers: A Reinforcement Learning <u>Approach</u>" *Proceedings of the <u>35th International Symposium on Computer Architecture</u> (ISCA), pages 39-50, Beijing, China, June 2008. <u>Slides (pptx)</u>*

Self-Optimizing Memory Controllers: A Reinforcement Learning Approach

Engin İpek^{1,2} Onur Mutlu² José F. Martínez¹ Rich Caruana¹

¹Cornell University, Ithaca, NY 14850 USA

² Microsoft Research, Redmond, WA 98052 USA



Memory Controllers: Many New Problems





Main Memory Needs Intelligent Controllers

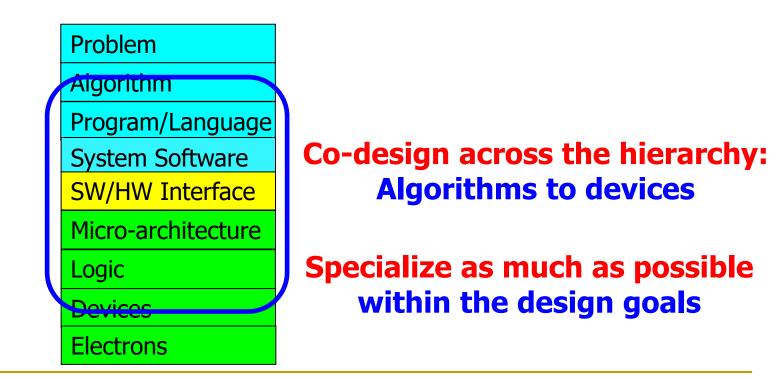


We Will See More Examples of This

To achieve the highest energy efficiency and performance:

we must take the expanded view

of computer architecture



Recommended Interview





SAFARI

https://www.youtube.com/watch?v=8ffSEKZhmvo

Recommended Interview

Computing Research and Education (@ ISCA 2019)

https://www.youtube.com/watch?v=8ffSEKZhmvo&list=PL5Q2 soXY2Zi 4oP9LdL3cc8G6NIjD2Ydz

Maurice Wilkes Award Speech (10 minutes)

https://www.youtube.com/watch?v=tcQ3zZ3JpuA&list=PL5Q2 soXY2Zi8D_5MGV6EnXEJHnV2YFBJl&index=15

Onur Mutlu, "Some Reflections (on DRAM)" *Award Speech for <u>ACM SIGARCH Maurice Wilkes Award</u>, at the ISCA Awards Ceremony, Phoenix, AZ, USA, 25 June 2019. [Slides (pptx) (pdf)] [Video of Award Acceptance Speech (Youtube; 10 minutes) (Youku; 13 minutes)] [Video of Interview after Award Acceptance (Youtube; 1 hour 6 minutes) (Youku; 1 hour 6 minutes)] [News Article on "ACM SIGARCH Maurice Wilkes Award goes to Prof. Onur Mutlu"]*

What We Will Cover in The Next Several Lectures

Agenda for The Next Several Lectures

- Computation in Memory (Processing in/near Memory)
- Some Key Issues: Data Retention & Memory Interference
- RowHammer: Memory Reliability and Security
- Low-Latency Memory
- Data-Driven and Data-Aware Architectures
- Memory Controllers and Memory QoS
- Guiding Principles & Research Topics

PIM Review and Open Problems

A Modern Primer on Processing in Memory

Onur Mutlu^{a,b}, Saugata Ghose^{b,c}, Juan Gómez-Luna^a, Rachata Ausavarungnirun^d

SAFARI Research Group

^aETH Zürich ^bCarnegie Mellon University ^cUniversity of Illinois at Urbana-Champaign ^dKing Mongkut's University of Technology North Bangkok

Onur Mutlu, Saugata Ghose, Juan Gomez-Luna, and Rachata Ausavarungnirun, <u>"A Modern Primer on Processing in Memory"</u> *Invited Book Chapter in <u>Emerging Computing: From Devices to Systems -</u> <u>Looking Beyond Moore and Von Neumann</u>, Springer, to be published in 2023*

A Modern Primer on Processing in Memory

Onur Mutlu^{a,b}, Saugata Ghose^{b,c}, Juan Gómez-Luna^a, Rachata Ausavarungnirun^d

SAFARI Research Group

^aETH Zürich ^bCarnegie Mellon University ^cUniversity of Illinois at Urbana-Champaign ^dKing Mongkut's University of Technology North Bangkok

Abstract

Modern computing systems are overwhelmingly designed to move data to computation. This design choice goes directly against at least three key trends in computing that cause performance, scalability and energy bottlenecks: (1) data access is a key bottleneck as many important applications are increasingly data-intensive, and memory bandwidth and energy do not scale well, (2) energy consumption is a key limiter in almost all computing platforms, especially server and mobile systems, (3) data movement, especially off-chip to on-chip, is very expensive in terms of bandwidth, energy and latency, much more so than computation. These trends are especially severely-felt in the data-intensive server and energy-constrained mobile systems of today.

At the same time, conventional memory technology is facing many technology scaling challenges in terms of reliability, energy, and performance. As a result, memory system architects are open to organizing memory in different ways and making it more intelligent, at the expense of higher cost. The emergence of 3D-stacked memory plus logic, the adoption of error correcting codes inside the latest DRAM chips, proliferation of different main memory standards and chips, specialized for different purposes (e.g., graphics, low-power, high bandwidth, low latency), and the necessity of designing new solutions to serious reliability and security issues, such as the RowHammer phenomenon, are an evidence of this trend.

This chapter discusses recent research that aims to practically enable computation close to data, an approach we call *processing-in-memory* (PIM). PIM places computation mechanisms in or near where the data is stored (i.e., inside the memory chips, in the logic layer of 3D-stacked memory, or in the memory controllers), so that data movement between the computation units and memory is reduced or eliminated. While the general idea of PIM is not new, we discuss motivating trends in applications as well as memory circuits/technology that greatly exacerbate the need for enabling it in modern computing systems. We examine at least two promising new approaches to designing PIM systems to accelerate important data-intensive applications: (1) *processing using memory* by exploiting analog operational properties of DRAM chips to perform massively-parallel operations in memory, with low-cost changes, (2) *processing near memory* by exploiting 3D-stacked memory technology design to provide high memory bandwidth and low memory latency to in-memory logic. In both approaches, we describe and tackle relevant cross-layer research, design, and adoption challenges in devices, architecture, systems, and programming models. Our focus is on the development of in-memory processing designs that can be adopted in real computing platforms at low cost. We conclude by discussing work on solving key challenges to the practical adoption of PIM.

Keywords: memory systems, data movement, main memory, processing-in-memory, near-data processing, computation-in-memory, processing using memory, processing near memory, 3D-stacked memory, non-volatile memory, energy efficiency, high-performance computing, computer architecture, computing paradigm, emerging technologies, memory scaling, technology scaling, dependable systems, robust systems, hardware security, system security, latency, low-latency computing

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3	The Need for Intelligent Memory Controllers to Enhance Memory Scaling			
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1. Introduction

Main memory, built using the Dynamic Random Access Memory (DRAM) technology, is a major component in nearly all computing systems, including servers, cloud platforms, mobile/embedded devices, and sensor systems. Across all of these systems, the data working set sizes of modern applications are rapidly growing, while the need for fast analysis of such data is increasing. Thus, main memory is becoming an increasingly significant bottleneck across a wide variety of computing systems and applications [1-26]. Alleviating the main memory bottleneck requires the memory capacity, energy, cost, and performance to all scale in an efficient manner across technology generations. Unfortunately, it has become increasingly difficult in recent years, especially the past decade, to scale all of these dimensions [1, 2, 27-59], and thus the main memory bottleneck has been worsening.

A major reason for the main memory bottleneck is the high energy and latency cost associated with data movement. In modern computers, to perform any operation on data that resides in main memory, the processor must retrieve the data from main memory. This requires the memory controller to issue commands to a DRAM module across a relatively slow and power-hungry off-chip bus (known as the memory channel). The DRAM module sends the requested data across the memory channel, after which the data is placed in the caches and registers. The CPU can perform computation on the data once the data is in its registers. Data movement from the DRAM to the CPU incurs long latency and consumes a significant amount of energy [7-9, 60-64]. These costs are often exacerbated by the fact that much of the data brought into the caches is not reused by the CPU [62, 63, 65, 66], providing little benefit in return for the high latency and energy cost.

The cost of data movement is a fundamental issue with the *processor-centric* nature of contemporary computer systems. The CPU is considered to be the master in the system, and computation is performed only in the processor (and accelerators). In contrast, data storage and communication units, including the main memory, are treated as unintelligent workers that are incapable of computation. As a result of this processor-centric design paradigm, data moves a lot in the system between the computation units and communication/storage units so that computation can be done on it. With the increasingly *data-centric* nature of contemporary and emerging applications, the processor-centric design paradigm leads to great inefficiency in performance, energy and cost. For example, most of the real estate within a single compute

СЛ	ЕЛ	DI
SA	ΓΑ	RI

PIM Review and Open Problems (II)

A Workload and Programming Ease Driven Perspective of Processing-in-Memory

Saugata Ghose†Amirali Boroumand†Jeremie S. Kim†§Juan Gómez-Luna§Onur Mutlu§††Carnegie Mellon University§ETH Zürich

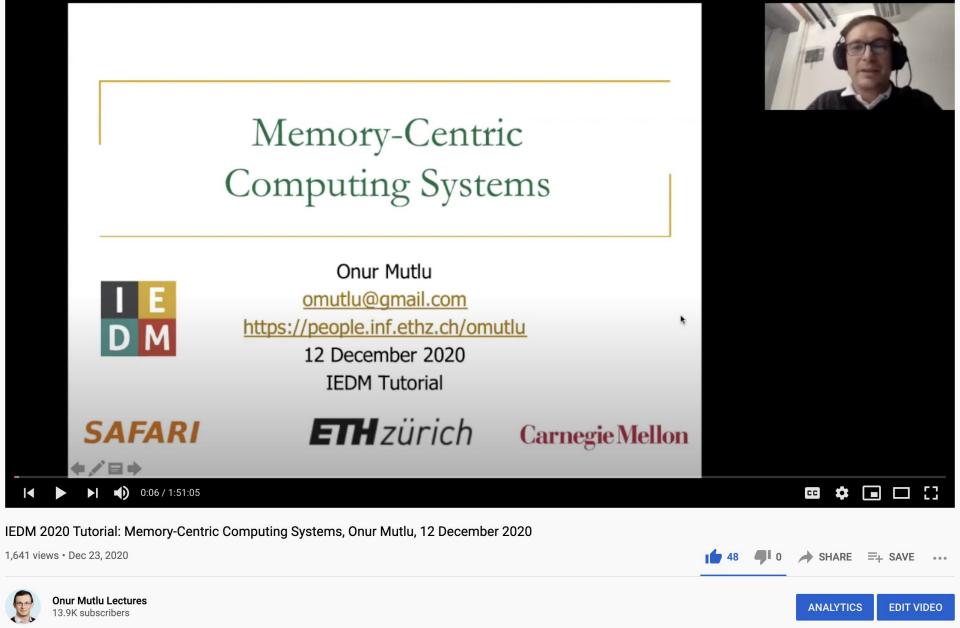
Saugata Ghose, Amirali Boroumand, Jeremie S. Kim, Juan Gomez-Luna, and Onur Mutlu, "Processing-in-Memory: A Workload-Driven Perspective" *Invited Article in IBM Journal of Research & Development, Special Issue on Hardware for Artificial Intelligence*, to appear in November 2019. [Preliminary arXiv version]

A Tutorial on PIM

Onur Mutlu, "Memory-Centric Computing Systems" Invited Tutorial at 66th International Electron Devices Meeting (IEDM), Virtual, 12 December 2020. Slides (pptx) (pdf) [Executive Summary Slides (pptx) (pdf)] [Tutorial Video (1 hour 51 minutes)] [Executive Summary Video (2 minutes)] [Abstract and Bio] [Related Keynote Paper from VLSI-DAT 2020] [Related Review Paper on Processing in Memory]

https://www.youtube.com/watch?v=H3sEaINPBOE

SAFARI https://www.youtube.com/onurmutlulectures



https://www.youtube.com/onurmutlulectures

An "Early" Position Paper [IMW'13]

 Onur Mutlu,
 <u>"Memory Scaling: A Systems Architecture Perspective"</u> *Proceedings of the <u>5th International Memory</u> <u>Workshop</u> (IMW), Monterey, CA, May 2013. <u>Slides</u> (pptx) (pdf) EETimes Reprint*

Memory Scaling: A Systems Architecture Perspective

Onur Mutlu Carnegie Mellon University onur@cmu.edu http://users.ece.cmu.edu/~omutlu/

https://people.inf.ethz.ch/omutlu/pub/memory-scaling_memcon13.pdf

An Extended Version: Memory Scaling

Onur Mutlu, <u>"Main Memory Scaling: Challenges and Solution Directions"</u> Invited Book Chapter in <u>More than Moore Technologies for Next</u> <u>Generation Computer Design</u>, pp. 127-153, Springer, 2015.

Chapter 6 Main Memory Scaling: Challenges and Solution Directions

Onur Mutlu, Carnegie Mellon University

Part of your Homework 1 assignment

A Recent Retrospective Paper [TCAD'19]

Onur Mutlu and Jeremie Kim,
 "RowHammer: A Retrospective"
 IEEE Transactions on Computer-Aided Design of Integrated Circuits and Systems (TCAD) Special Issue on Top Picks in Hardware and Embedded Security, 2019.
 [Preliminary arXiv version]
 [Slides from COSADE 2019 (pptx)]
 [Slides from VLSI-SOC 2020 (pptx) (pdf)]
 [Talk Video (1 hr 15 minutes, with Q&A)]

RowHammer: A Retrospective

Onur Mutlu§‡Jeremie S. Kim‡§§ETH Zürich‡Carnegie Mellon University

Challenges in Memory Scaling

- Data retention (need for refresh)
- Reliability and vulnerabilities (e.g., RowHammer)
- Latency and parallelism (e.g., bank conflicts)
- Energy & power
- Memory's inability to do anything more than just store data

Computer Architecture Lecture 2a: Memory Systems: Challenges and Opportunities

Prof. Onur Mutlu ETH Zürich Fall 2022 30 September 2022