

Drammer: Deterministic Rowhammer Attacks on Mobile Platforms

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Outline

- Summary
- Problem
- Background
- Goal
- Mechanisms
- Key Results: Methodology and Evaluation
- Novelty, Key Approach and Ideas
- Strengths
- Weaknesses
- Thoughts and Ideas
- Takeaways
- Questions/Open Discussion

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Summary

- ARM based devices can be vulnerable to Drammer as well.
- Drammer:
 - Memory templating
 - Scan memory for vulnerable bits
 - Land sensitive data
 - Reproduce the bit flip
- Root access exploitation possible with high reliability.
 - By modifying entries in Page Table Pages (PTP)
- Severe consequences for numerous devices that are currently in use.

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Problem

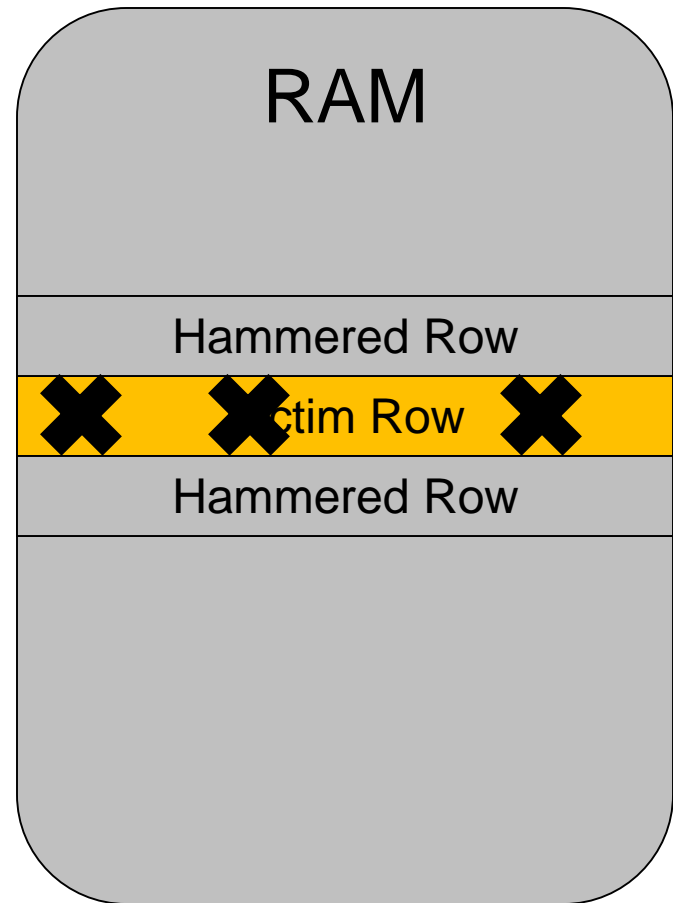
- Rowhammer failure mechanism only exploitable in a probabilistic way
 - Can it be done deterministically?
- Not clear if Rowhammer attacks are possible on ARM
 - Some researcher thought that it might be impossible on ARM

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Rowhammer Failure Mechanism

- Access adjacent rows at a high frequency (hammer)
- This causes voltage leakage in the victim row
- Bit flips are induced



Primitives to exploit Rowhammer Bug – x86

- P1: Fast uncached memory access
 - Explicit cache flush (c1flush instruction)
 - Cache eviction sets
 - Non-temporal access instructions
- P2: Physical memory massaging
 - Page-table spraying (probabilistic)
- P3: Physical memory addressing
 - Pagemap interface
 - Huge pages

Rowhammer Bug on x86

- Probabilistic
 - Page-table spraying used because we don't know exactly where bitflips will occur.
- Countermeasures
 - Disable `clflush`
 - Error Correcting Codes (ECC)
 - Probabilistic Adjacent Row Activation (PARA)
 - Many more

Primitives to exploit Rowhammer Bug – ARM

- P1: Fast uncached memory access
 - Explicit cache flush (c1flush instruction) Privileged instruction
 - Cache eviction sets Too slow
 - Non-temporal access instructions Only suggests to not cache it
- P2: Physical memory massaging
 - Page-table spraying (probabilistic) Can crash the system
- P3: Physical memory addressing
 - Pagemap interface No unprivileged access anymore
 - Huge pages Disabled on stock Android

Rowhammer Failure Mechanism

- x86 Architectures are known to be vulnerable if the DRAM is modern enough.
 - Are ARM architectures vulnerable as well?
- Before this paper:
 - Probabilistic Rowhammer attacks on x86 based devices
 - Low reliability
 - Limited impact in practice
- After this paper:
 - Deterministic Rowhammer attacks
 - High reliability
 - Allows to completely subvert any vulnerable system
 - Requires to trick the OS to put a page table in a known and vulnerable memory location.

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Goal

- *"Deterministic Rowhammer Attack on Mobile Platforms"*
 - Deterministic:
 - Memory templating
 - Land sensitive Data → Phys Feng Shui
 - Reproduce the bit flip
 - *Note: This approach is not completely deterministic but much more reliable than the probabilistic approach that we know from Rowhammer attacks on x86.*
 - Rowhammer Attack:
 - Vulnerable system required
 - Mobile Platforms:
 - ARMv7, ARMv8 running Android

Primitives to exploit Rowhammer Bug - ARM

- P1: Fast uncached memory access
 - Provided by DMA buffer management APIs (ION)
- P2: Physical memory massaging
 - Memory Templating
 - Physical Memory Allocator
 - Buddy allocator
 - Phys Feng Shui
- P3: Physical memory addressing
 - Provided by DMA buffer management APIs (ION)

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Phys Feng Shui

- Variant of Flip Feng Shui
 - FFS was used to mount attacks against other guest OSes running on the same hypervisor.
- Goal
 - Force the OS to place a page-table in a vulnerable memory location such that we can modify an entry in a deterministic way.
- Some size definitions
 - S chunk = chunk of the size of a page (typically 4KB)
 - M chunk = chunk of the size of a row (has to be determined)
 - L chunk = largest possible contiguous chunk

Determining Row Size

- Access time for page pairs
- If access time increases then the pages are on different rows
 - Therefore we can determine the row size

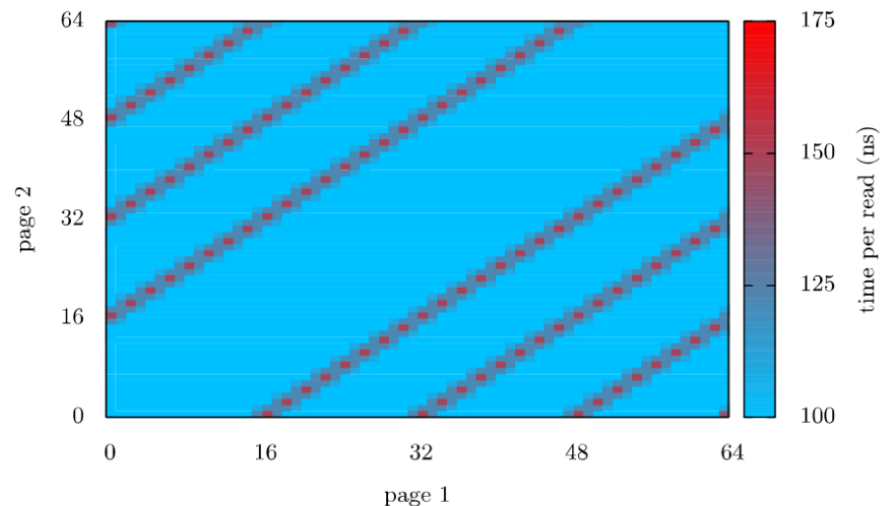
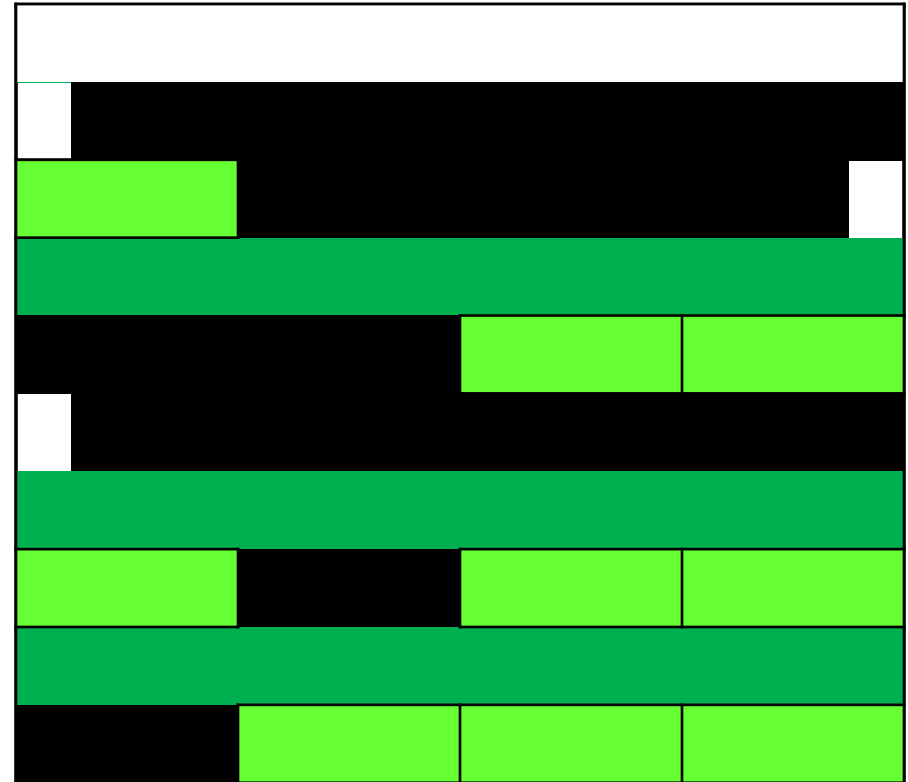


Figure 3: Heatmap representing the time required to access a given pair of pages on a LG Nexus 5. The diagonal pattern clearly indicates that the row size is 16 pages = 64K.

(0,1) (0,2) (0,3) ... (0,16)

Phys Feng Shui

- Step 1: Fill in as many L chunks as possible and create templates
- Step 2: Fill in as many M chunks as possible
- Step 3: Free one L chunk where we want to launch the attack



Free(L*)

Phys Feng Shui

- Step 4: Fill L^* with M chunks
- Step 5: Free M^* and all L chunks (M^* is where we launch the attack)
- Step 6: Fill in S chunks until the first one falls into M^*
- Step 7: Add padding to align victim page table
- Step 8: Launch attack on the victim page table that then points to a page table which we created in L^*



Map(M)

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Empirical Results

25

20

15

10

5

0

Device	Hardware Details			Analysis Results								
	SoC	DRAM	RS	MB	ns	#flips	KB	# 1-to-0	# 0-to-1	# exploitable	1 st	
Nexus 5x	MSM8974A	2 GB	64	441	70	1,058	426	1,011	47	62 (5.86%)	116s	
Nexus 5	MSM8974A	2 GB	64	472	69	284,428	2	261,232	23,196	14,852 (5.22%)	1s	
Nexus 4	MSM8974A	2 GB	64	461	69	547,949	1	534,695	13,254	32,715 (5.97%)	1s	
Nexus 7	MSM8974A	2 GB	64	616	71	0	-	-	-	-	-	
Nexus 7	MSM8974A	2 GB	64	630	69	747,013	1	704,824	42,189	46,609 (6.24%)	1s	
Nexus 7	MSM8974A	2 GB	64	512	69	215,233	3	207,856	7,377	13,365 (6.21%)	3s	
Nexus 7	MSM8974A	2 GB	64	485	70	32,328	15	28,500	3,828	1,894 (5.86%)	4s	
Nexus 7	MSM8974A	2 GB	64	569	69	476,170	2	434,086	42,084	30,190 (6.34%)	0s	
Nexus 7	MSM8974A	2 GB	64	406	69	160,245	3	150,485	9,760	8,701 (5.43%)	1s	
ARMv7 Nexus 7	MSM8974A	2 GB	64	613	70	0	-	-	-	-	-	
Nexus 7	MSM8974A	2 GB	64	600	70	17,384	35	16,767	617	1,241 (7.14%)	16s	
Nexus 7	MSM8974A	2 GB	64	575	69	161,514	4	160,473	1,041	10,378 (6.43%)	355s	
Nexus 7	MSM8974A	2 GB	64	576	69	295,537	2	277,708	17,829	18,900 (6.40%)	1s	
Nexus 7	MSM8974A	2 GB	64	573	69	38,969	15	35,515	3,454	2,775 (7.12%)	11s	
ARMv7 Nexus 7	MSM8974A	2 GB	64	621	70	0	-	-	-	-	-	
Galaxy S3	Exynos 4210	1 GB	32	207	82	0	-	-	-	-	-	
OnePlus	MSM8974A	2 GB	64	292	71	3,981	75	2,924	1,057	242 (6.08%)	942s	
Motorola	MSM8974A	2 GB	64	1189	69	1,992	611	942	1,050	94 (4.72%)	326s	
Motorola	MSM8974A	2 GB	64	134	127	429	275	-	-	-	11s	
Nexus 7	MSM8974A	2 GB	64	151	127	1,577	98	-	-	-	92s	
Nexus 7	MSM8974A	2 GB	64	82	18	1,328	64	-	-	-	7s	
ARMv8 Nexus 7	MSM8974A	2 GB	64	271	63	0	-	-	-	-	-	
Galaxy S3	Exynos 4210	1 GB	32	234	82	0	-	-	-	-	-	
K3	Exynos 4210	1 GB	32	423	218	0	-	-	-	-	-	
Mi 4i	MSM8939	2 GB	64	327	159	0	-	-	-	-	-	
Desire 510	MSM8916	1 GB	32	186	122	0	-	-	-	-	-	
G4	ARMv7 MSM8974A	3 GB	64	833	64	117,496	8	117,260	117,260 (5.58%)	5s	ARMv8	

[†]MSM8974AA [‡]MSM8974AC ^{*}LPDDR2 [°]LPDDR4

■ Vulnerable ■ Resilient

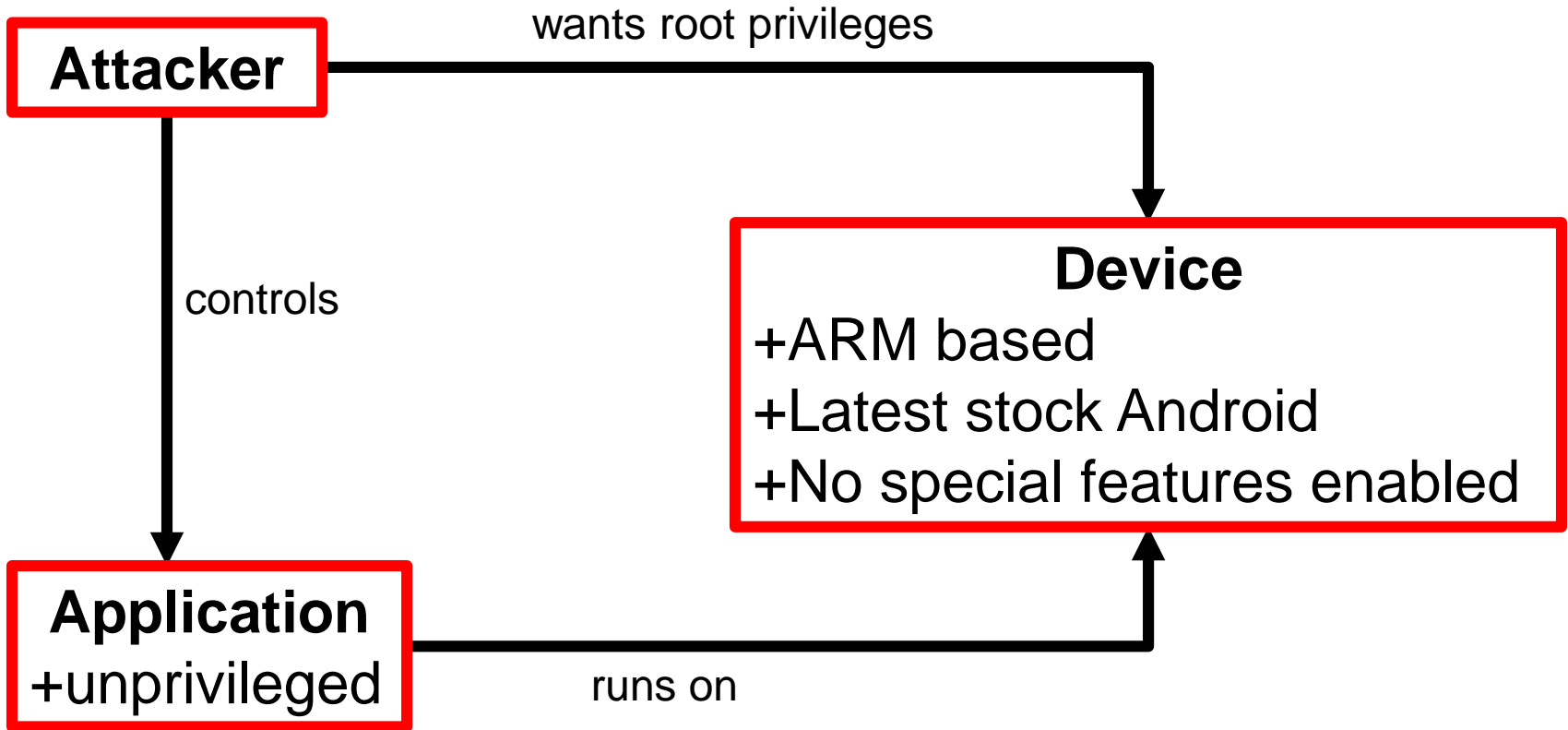
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Threat Model

- Set up:
 - ❑ Attacker has control over an unprivileged app running on a ARM based Android device.
 - ❑ No permissions at all.
 - ❑ Latest stock version of Android with all the latest security updates installed.
 - ❑ No special features enabled.
- Goal of the attacker:
 - ❑ To mount an privilege escalation attack to acquire root privileges.

Threat Model



Novelty, Key Approach and Ideas

- First deterministic Rowhammer attack on ARM architecture
- Generalization of deterministic Rowhammer attacks on x86 architecture
- Mounting a Drammer attack using Direct Memory Access (DMA) bypasses existing defenses (i.e. disabling `clflush`) on x86 architectures

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Strengths

- Shows a severe problem current mobile devices have.
 - First paper that shows a Rowhammer attack on ARM architecture based devices.
- The attack is deterministic.
 - Much more serious consequences in practice.
- Very detailed description of the work that was done

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Weaknesses

- All countermeasures have an overhead
- ARMv8 sample size too small
 - No representative conclusion possible
 - No explanation why they seem more resilient
- Tested only Android smartphones (i.e. no Smartwatches, no iPhones, etc.)
- Considering that the occurrence of bitflips also depends on environmental aspects, the sample size is clearly too small.
- Not so easy to understand (Phys Feng Shui)

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Thoughts and Ideas

- Future research needs to test the influence factors like temperature, age of the components, etc. have on the number of vulnerable bits.
- A broader range of devices needs to be tested (also running other OSes than Android).
- More effective and efficient countermeasures need to be found.

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Takeaways

- Rowhammer on ARM
- Deterministic exploitation
- Practical impact

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QUESTIONS?

Discussion Question

- What do you consider being more valuable regarding security?
 - Opensource Android
 - Known that Drammer is possible
 - Countermeasures proposed by third parties
 - Proprietary iOS
 - Unsure if Drammer is possible
 - If possible then there are not yet any countermeasures

Discussion Question

Opensource Android

- Positive:
 - Huge research community doing research on Android
 - More vulnerabilities get found and can be patched
- Negative:
 - Attacker can learn more about the OS
 - Practical attacks happen faster

Proprietary iOS

- Positive:
 - Attackers have it harder to mount a practical attack since they know less about the OS
- Negative:
 - Once an attacker is successful, he can be active for a very long time
 - Fewer security researchers

Discussion Question

- Do you think it is okay that the researchers published the paper before Google was able to patch it's devices?

Discussion Question

- Researchers reported the attack to Google 91 days before the release of the paper at CCS 2016
 - Including some mitigation techniques
- Google asked them to delay the release of the paper
- The researchers refused
- Google asked them to obfuscate parts of the paper
- The researchers refused again
- → Did the researchers act responsibly?

Discussion Question

- Can you think of countermeasures that have little to no overhead?

Discussion Question

- Hardware based:
 - Probabilistic Adjacent Row Activation (PARA) looks promising
 - Better isolation between rows in DRAM?
- Software based:
 - Disallow features that can be used to satisfy the primitives P1-P3
 - Detect access patterns that could imply that an attack is happening

Supplementary Material

- [Video] CCS 2016 – Drammer: Deterministic Rowhammer Attacks on Mobile Platforms
 - <https://www.youtube.com/watch?v=ITaMvBN1PoA>
- [Video] Computer Architecture – Lecture 2: RowHammer and Beyond (ETH Zürich, Fall 2018)
 - <https://www.youtube.com/watch?v=560JzQ-oeLE>