Digital Design & Computer Arch.

Lecture 7: Hardware Description Languages and Verilog

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Required Readings (This Week)

- Hardware Description Languages and Verilog
  - H&H Chapter 4 in full

- Timing and Verification
  - H&H Chapters 2.9 and 3.5 + (start Chapter 5)

- By tomorrow, make sure you are done with
  - **P&P Chapters 1-3 + H&H Chapters 1-4**

Required Readings (Next Week)

- Von Neumann Model, LC-3, and MIPS
  - P&P, Chapter 4, 5
  - H&H, Chapter 6
  - P&P, Appendices A and C (ISA and microarchitecture of LC-3)
  - H&H, Appendix B (MIPS instructions)

- Programming
  - P&P, Chapter 6

- **Recommended:** Digital Building Blocks
  - H&H, Chapter 5

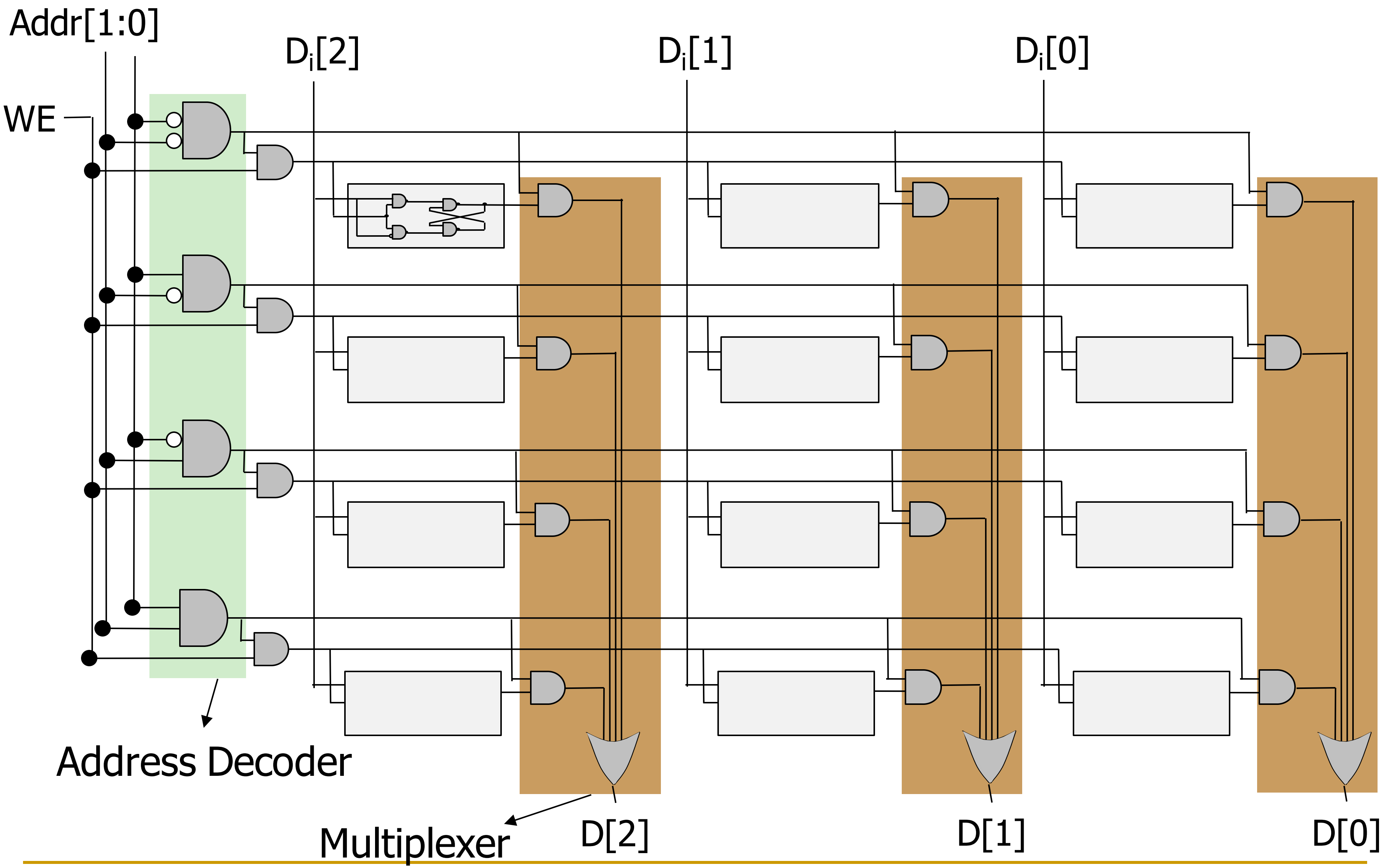
Agenda

- Hardware Description Languages
- Implementing Combinational Logic (in Verilog)
- Implementing Sequential Logic (in Verilog)

- The Verilog slides constitute a tutorial. We will not cover all.
- All slides will be beneficial for your labs.

Aside: Implementing Logic Functions Using Memory

Recall: A Bigger Memory Array (4 locations X 3 bits)



Memory-Based Lookup Table Example

- Memory arrays can also perform Boolean Logic functions
  - 2^N-location M-bit memory can perform any N-input, M-output function
  - Lookup Table (LUT): Memory array used to perform logic functions
  - Each address: row in truth table; each data bit: corresponding output value

Truth Table

A	B	Y
0	0	0
0	1	0
1	0	0
1	1	1

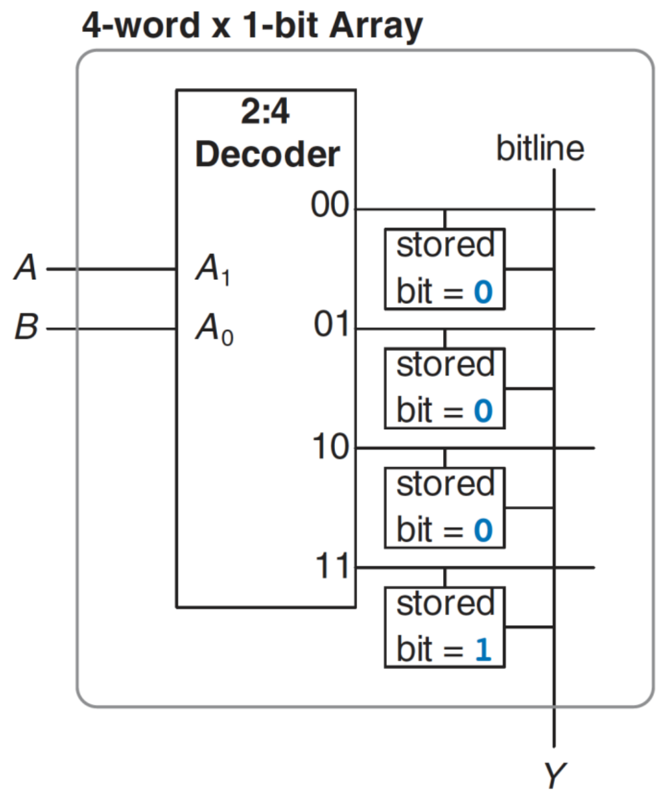
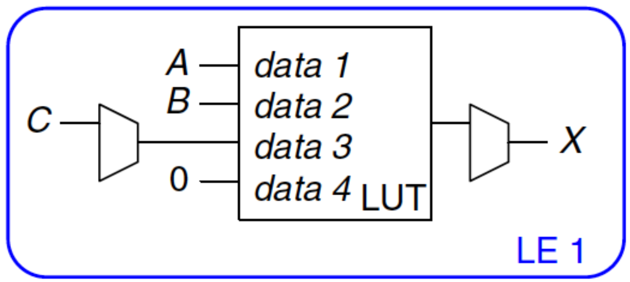


Figure 5.52 4-word × 1-bit memory array used as a lookup table

Lookup Tables (LUTs)

- LUTs are commonly used in FPGAs
  - To enable programmable/reconfigurable logic functions
  - To enable easy integration of combinational and sequential logic

(A) data 1	(B) data 2	(C) data 3	data 4	(X) LUT output
0	0	0	X	0
0	0	1	X	1
0	1	0	X	0
0	1	1	X	0
1	0	0	X	0
1	0	1	X	0
1	1	0	X	1
1	1	1	X	0



(A) data 1	(B) data 2	data 3	data 4	(Y) LUT output
0	0	X	X	0
0	1	X	X	0
1	0	X	X	1
1	1	X	X	0

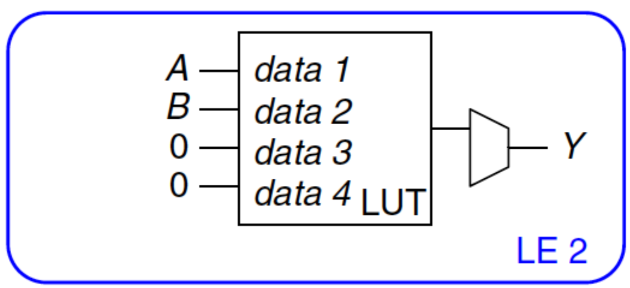
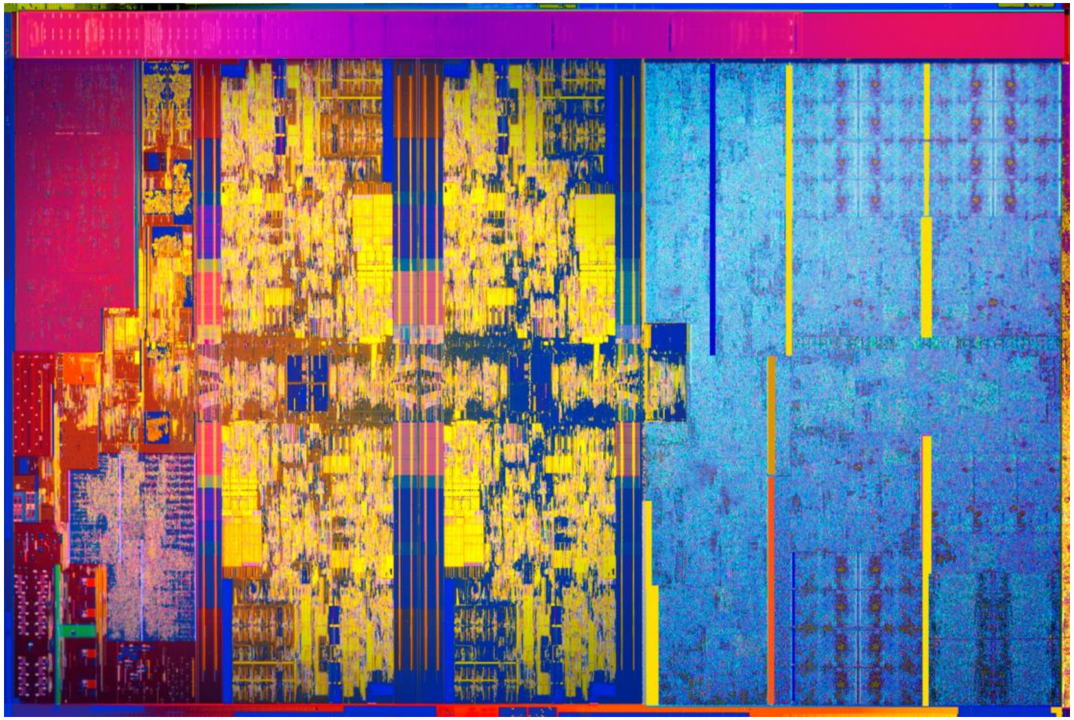


Figure 5.59 **LE configuration for two functions of up to four inputs each**

Read H&H Chapter 5.6.2

Hardware Description Languages & Verilog

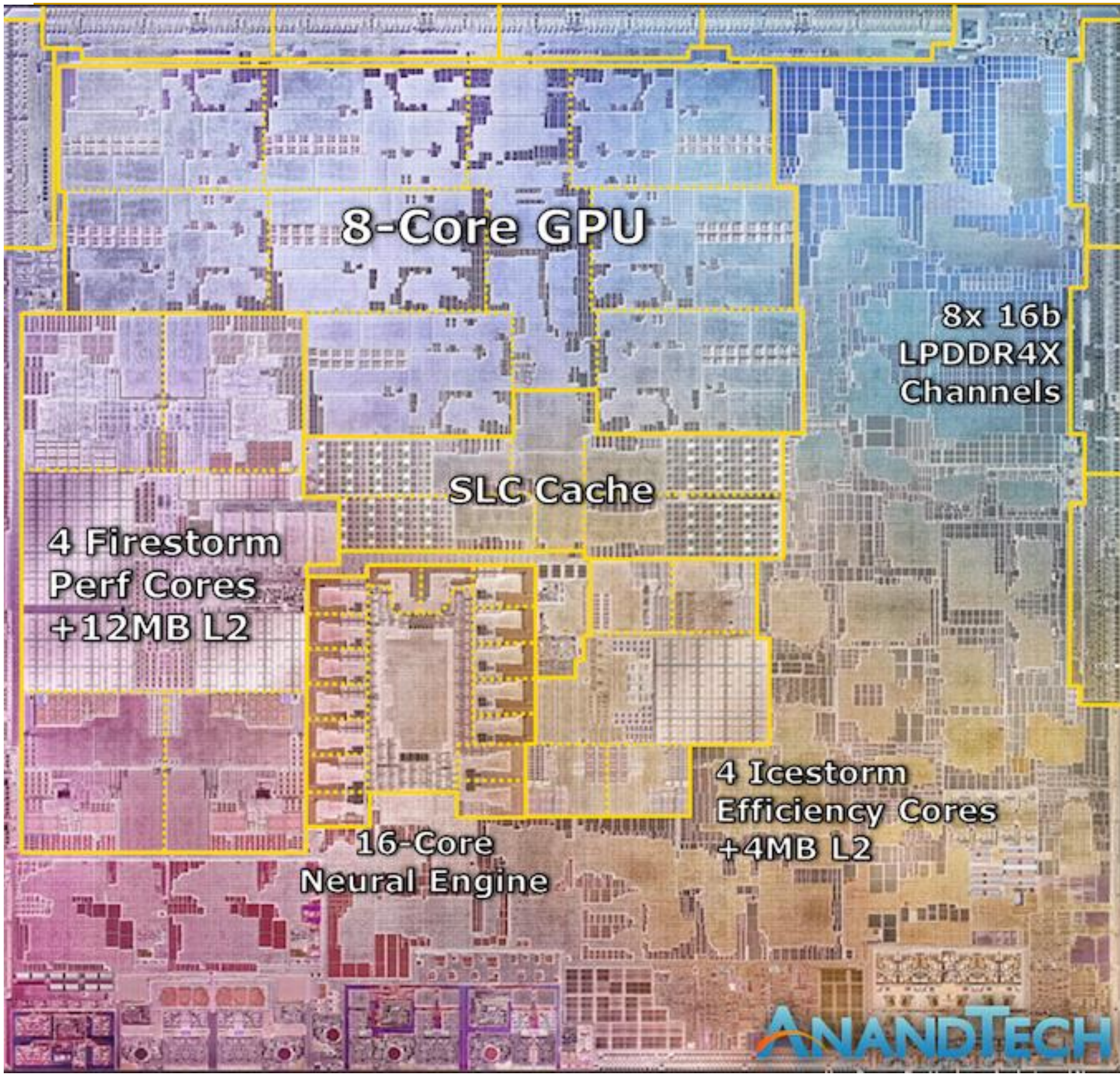
2017: Intel Kaby Lake



https://en.wikichip.org/wiki/intel/microarchitectures/kaby_lake

- 64-bit processor
- 4 cores, 8 threads
- 14-19 stage pipeline
- 3.9 GHz clock freq.
- 1.75B transistors
- In ~47 years, about 1,000,000-fold growth in transistor count and performance!

2021: Apple M1

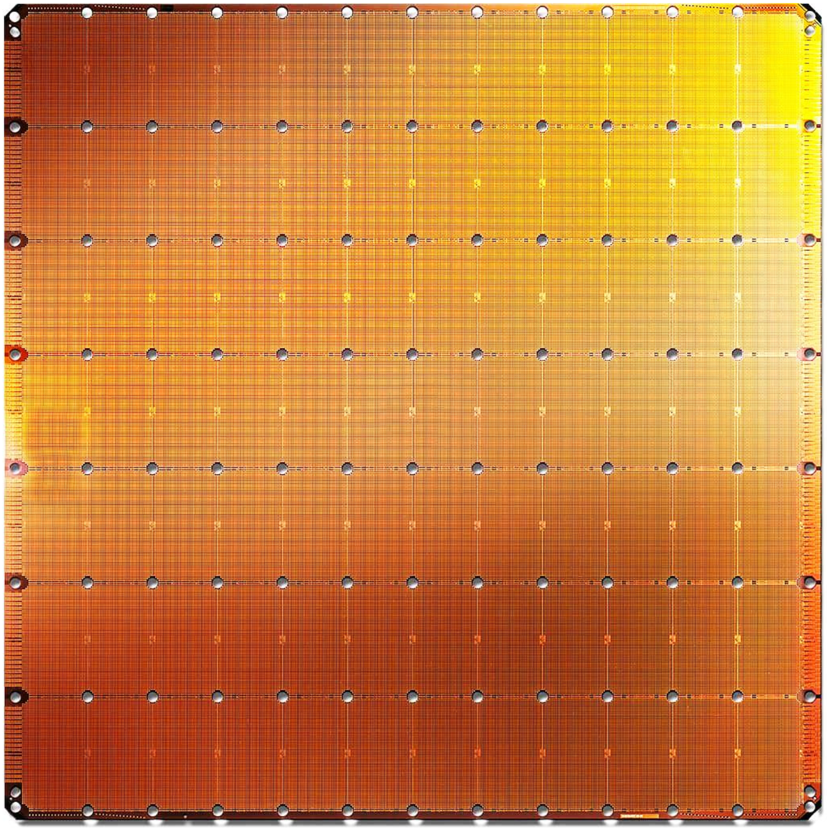


- 4 High-Perf GP Cores
- 4 Efficient GP Cores
- 8-Core GPU
- 16-Core Neural Engine
- Lots of Cache
- Many Caches
- 8x Memory Channels

- 16B transistors

Source: https://www.anandtech.com/show/16252/mac-mini-apple-m1-tested

2019: Cerebras Wafer Scale Engine



- The largest ML accelerator chip
- 400,000 cores

Cerebras WSE
1.2 Trillion transistors
46,225 mm^2

Largest GPU
21.1 Billion transistors
815 mm^2

NVIDIA TITAN V

https://www.anandtech.com/show/14758/hot-chips-31-live-blogs-cerebras-wafer-scale-deep-learning

https://www.cerebras.net/cerebras-wafer-scale-engine-why-we-need-big-chips-for-deep-learning/

How to Deal with This Complexity?

- Hardware Description Languages!

- Needs and wants:
  - Ability to specify complex designs
  - … and to simulate their behavior (functional & timing)
  - … and to synthesize (automatically design) portions of it
    - have an error-free path to implementation

- Hardware Description Languages enable all of the above
  - Languages designed to describe and specify hardware
  - There are similarly-featured HDLs (e.g., **Verilog**, VHDL, ...)
    - if you learn one, it is not hard to learn another
    - mapping between languages is typically mechanical, especially for the commonly used subset

Hardware Description Languages

- **Two well-known hardware description languages**

- **Verilog**
  - Developed in 1984 by Gateway Design Automation
  - Became an IEEE standard (1364) in 1995
  - More popular in US

- **VHDL (VHSIC Hardware Description Language)**
  - Developed in 1981 by the US Department of Defense
  - Became an IEEE standard (1076) in 1987
  - More popular in Europe

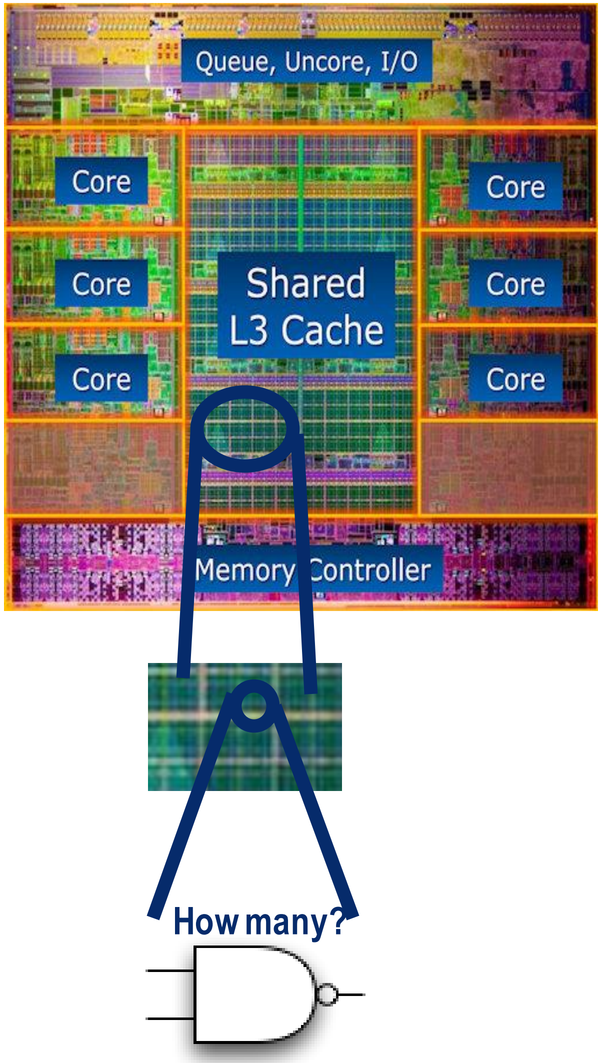
- We will use Verilog in this course

Hardware Design Using HDL

Principle: Hierarchical Design

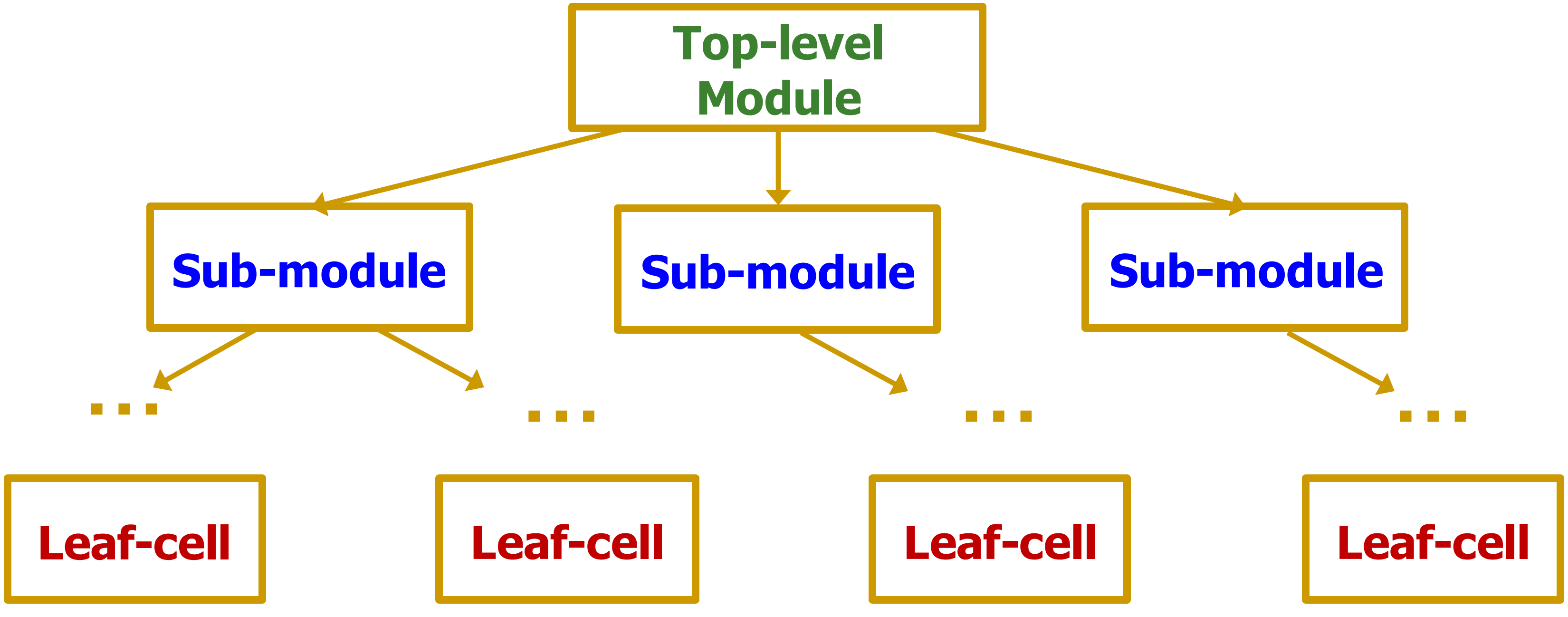
- **Design a hierarchy of modules**
  - Predefined "primitive" gates (AND, OR, …)
  - Simple modules are built by instantiating these gates (components like MUXes)
  - Complex modules are built by instantiating simple modules, …
- **Hierarchy controls complexity**
  - Analogous to the use of function/method abstraction in programming
- **Complexity is a BIG deal**
  - In real world, how big is the size of a module (that is described in HDL and then synthesized to gates)?

https://techreport.com/review/21987/intel-core-i7-3960x-processor



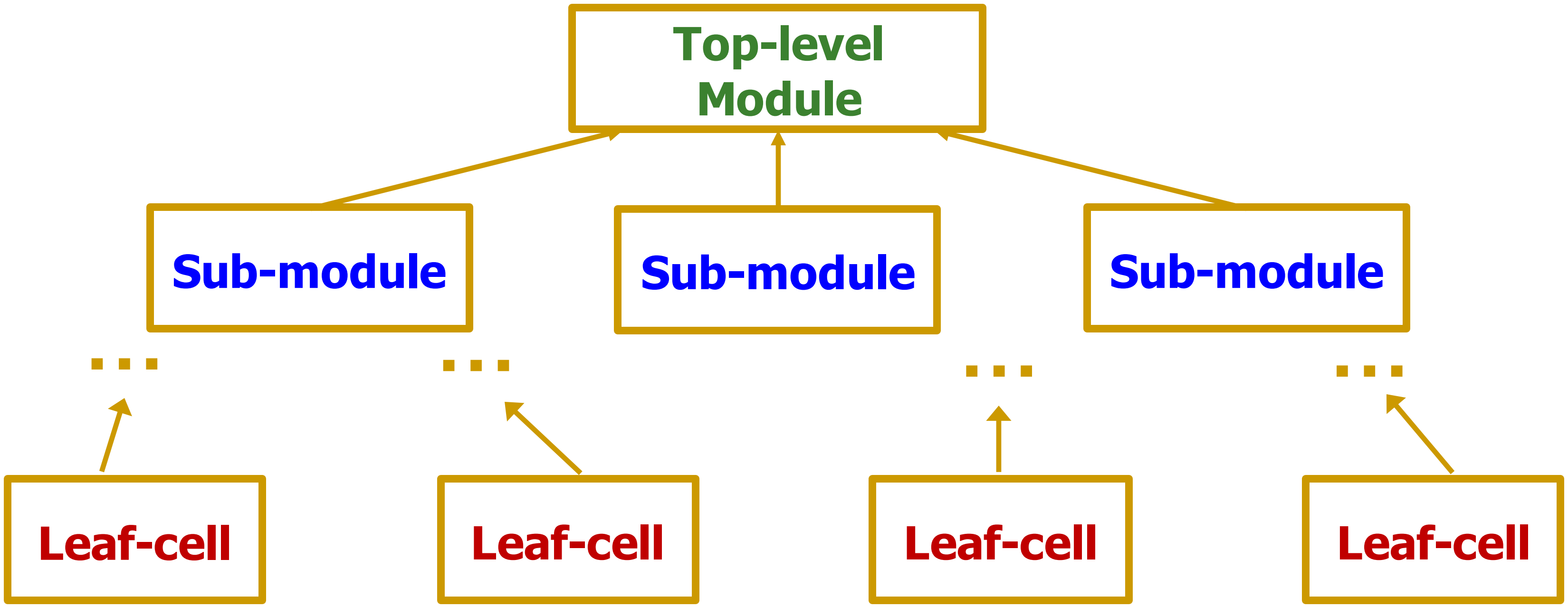
Top-Down Design Methodology

- We define the top-level module and identify the sub-modules necessary to build the top-level module
- Subdivide the sub-modules until we come to leaf cells
  - Leaf cell: circuit components that cannot further be divided (e.g., *logic gates, cell libraries*)



Bottom-Up Design Methodology

- We first identify the building blocks that are available to us
- Build bigger modules, using these building blocks
- These modules are then used for higher-level modules until we build the top-level module in the design

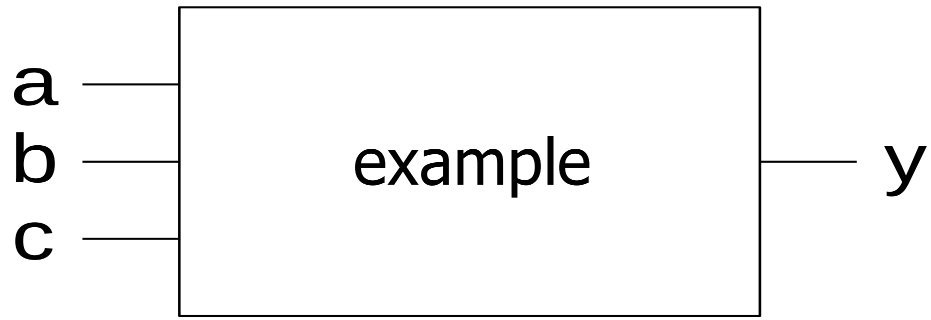


Defining a Module in Verilog

- A module is the main building block in Verilog

- We first need to define:
  - Name of the module
  - Directions of its ports (e.g., input, output)
  - Names of its ports
- Then:
  - Describe the functionality of the module

Implementing a Module in Verilog



name of module

Port list (inputs and outputs)

ports have a declared type

a module definition

```
module example (a, b, c, y);
        input a;
        input b;
        input c;
        output y;

// here comes the circuit description

endmodule
```

A Question of Style

- **The following two codes are functionally identical**

```
module test ( a, b, y );
        input a;
        input b;
        output y;

endmodule
```

```
module test ( input a,
              input b,
              output y );

endmodule
```

port name and direction declaration can be combined

What If We Have Multi-bit Input/Output?

- **You can also define multi-bit Input/Output (Bus)**
  - [range_end : range_start]
  - **Number of bits:** range_end – range_start + 1
- **Example**:

```
input  [31:0] a;   // a[31], a[30] .. a[0]
output [15:8] b1;  // b1[15], b1[14] .. b1[8]
output [7:0]  b2;  // b2[7], b2[6] .. b2[0]
input         c;   // single signal
```

- **a** represents a 32-bit value, so we prefer to define it as: `[31:0] a`
- It is preferred over `[0:31] a` which resembles *array* definition
- It is good practice to be consistent with the representation of multi-bit signals, i.e., always [31:0] or always [0:31]

Manipulating Bits

- Bit Slicing
- Concatenation
- Duplication

Basic Syntax

- Verilog is case sensitive
  - `SomeName` and `somename` are not the same!
- Names cannot start with numbers:
  - `2good` is not a valid name
- Whitespaces are ignored

```
// Single line comments start with a //

/* Multiline comments
   are defined like this */
```

Two Main Styles of HDL Implementation

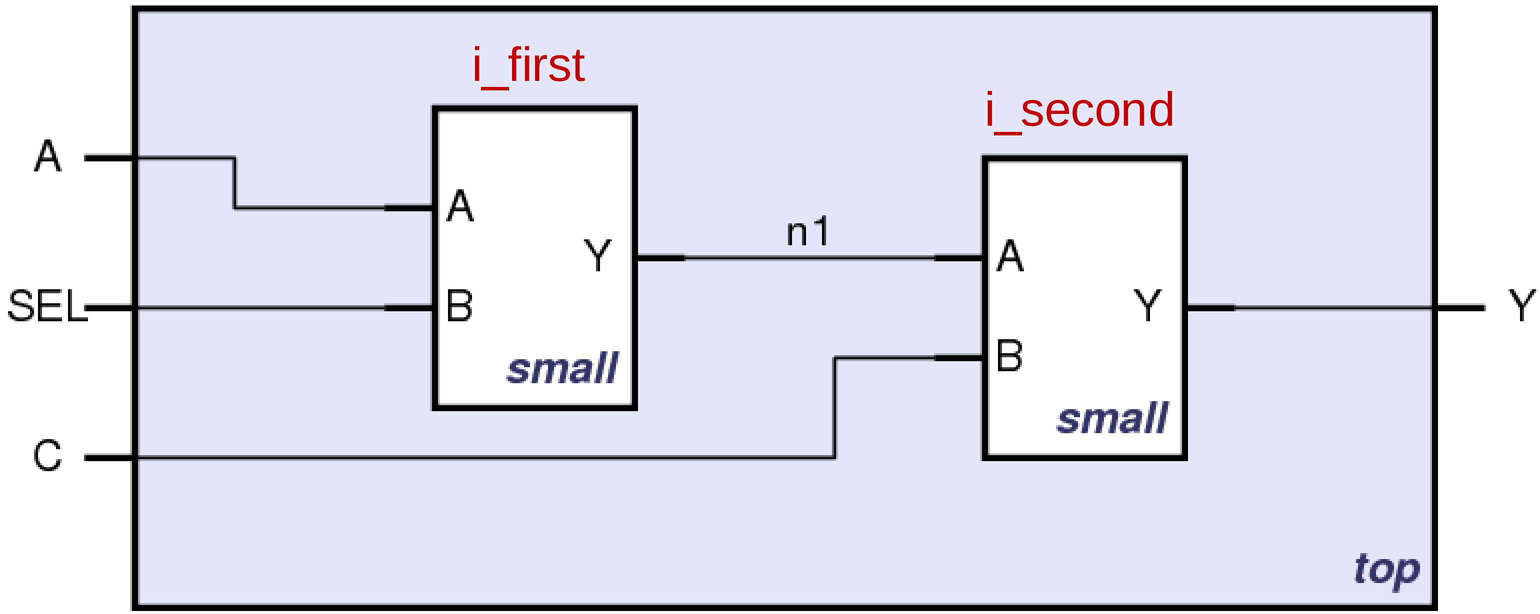
- **Structural (Gate-Level)**
  - The module body contains gate-level description of the circuit
  - Describe how modules are interconnected
  - Each module contains other modules (instances)
  - … and interconnections between those modules
  - Describes a hierarchy of modules defined as gates

- **Behavioral**
  - The module body contains functional description of the circuit
  - Contains logical and mathematical operators
  - **Level of abstraction is higher than gate-level**
    - Many possible gate-level realizations of a behavioral description

- **Many practical designs use a combination of both**

Structural (Gate-Level) HDL

Structural HDL: Instantiating a Module



Schematic of module "top" that is built from two instances of module "small"

Structural HDL Example

- ## Module Definitions in Verilog

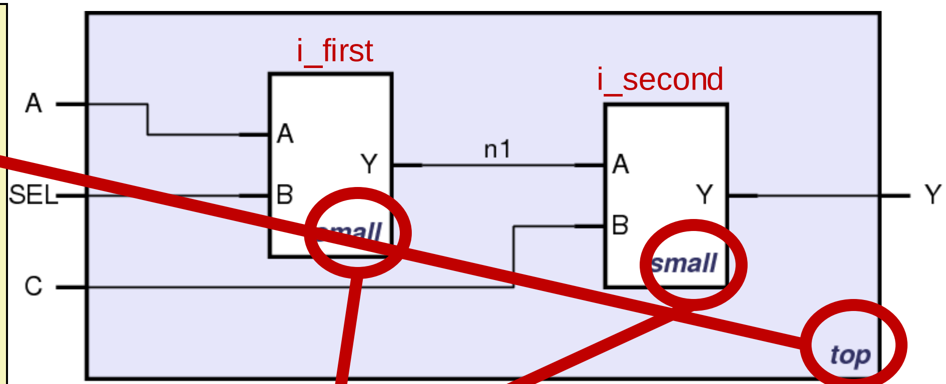
```
module top (A, SEL, C, Y);
  input A, SEL, C;
  output Y;
  wire n1;







endmodule
```




```
module small (A, B, Y);
  input A;
  input B;
  output Y;

// description of small

endmodule
```

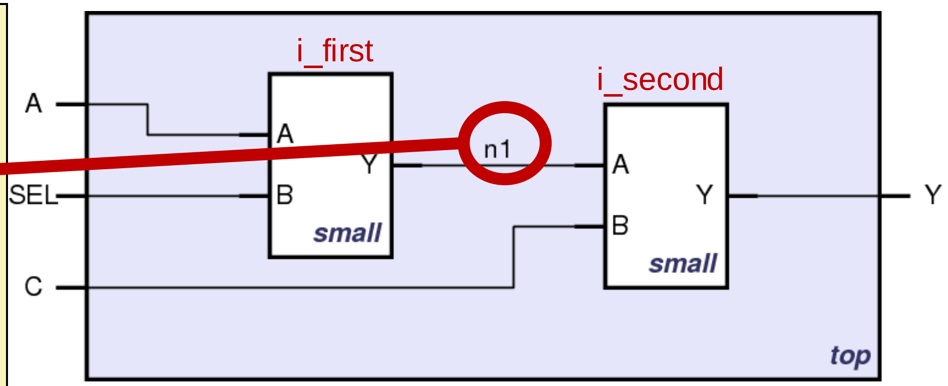
Structural HDL Example

- **Defining wires (module interconnections)**

```
module top (A, SEL, C, Y);
  input A, SEL, C;
  output Y;
  wire n1;



endmodule
```




```
module small (A, B, Y);
  input A;
  input B;
  output Y;

// description of small

endmodule
```

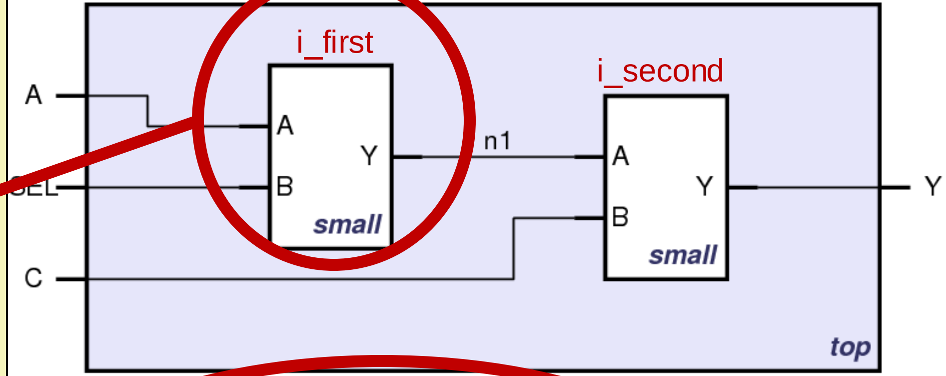
Structural HDL Example

- **The first instantiation of the “small” module**

```
module top (A, SEL, C, Y);
  input A, SEL, C;
  output Y;
  wire n1;

// instantiate small once
small i_first ( .A(A),
                .B(SEL),
                .Y(n1)   );

endmodule
```




```
module small (A, B, Y);
  input A;
  input B;
  output Y;

// description of small

endmodule
```

Structural HDL Example

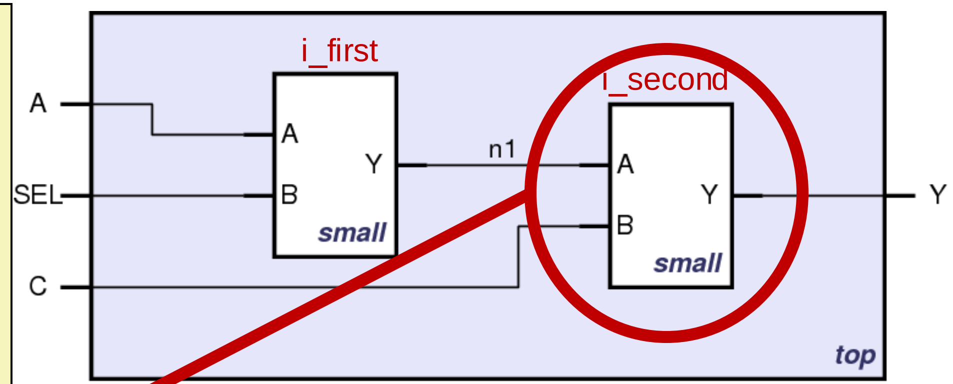
- **The second instantiation of the "small" module**

```
module top (A, SEL, C, Y);
  input A, SEL, C;
  output Y;
  wire n1;

// instantiate small once
small i_first ( .A(A),
                .B(SEL),
                .Y(n1)   );

// instantiate small second time
small i_second ( .A(n1),
                 .B(C),
                 .Y(Y) );

endmodule
```




```
module small (A, B, Y);
  input A;
  input B;
  output Y;

// description of small

endmodule
```

Structural HDL Example

- **Short form of module instantiation**

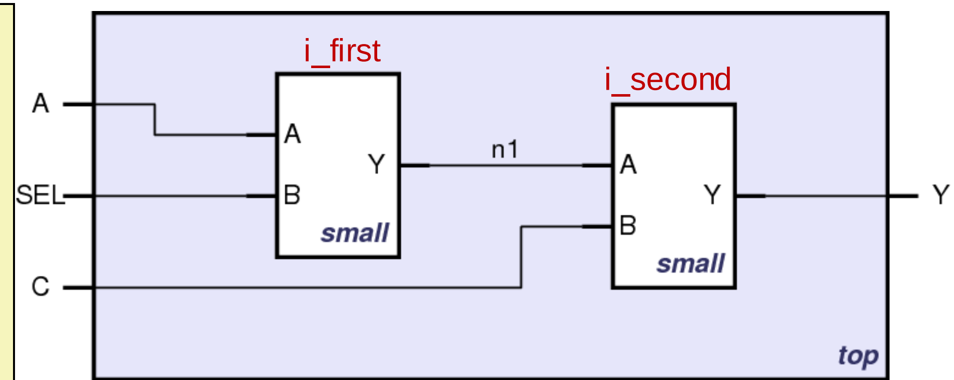
```
module top (A, SEL, C, Y);
  input A, SEL, C;
  output Y;
  wire n1;

// alternative
small i_first ( A, SEL, n1 );

/* Shorter instantiation,
   pin order very important */

// any pin order, safer choice
small i_second ( .B(C),
                 .Y(Y),
                 .A(n1) );

endmodule
```




```
module small (A, B, Y);
  input A;
  input B;
  output Y;

// description of small

endmodule
```

Short form is not good practice
as it reduces code maintainability

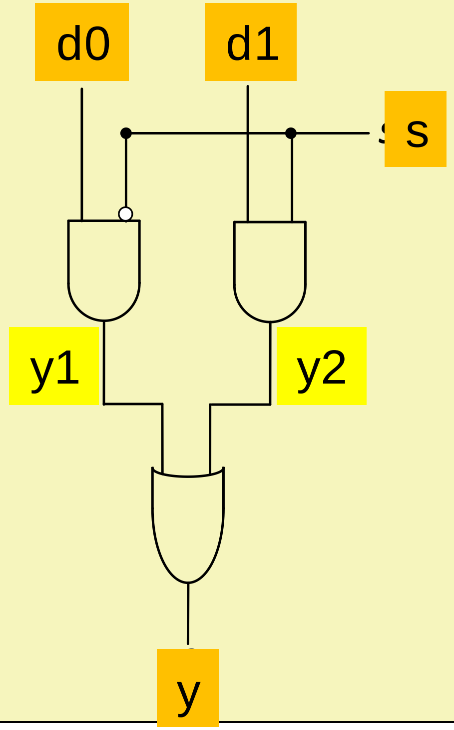
Structural HDL Example (II)

- Verilog supports basic logic gates as predefined *primitives*
  - These primitives are instantiated like modules except that they are predefined in Verilog and *do not need a module definition*

```
module mux2(input d0, d1,
            input s,
            output y);
  wire ns, y1, y2;

  not  g1 (ns, s);
  and  g2 (y1, d0, ns);
  and  g3 (y2, d1, s);
  or   g4 (y, y1, y2);

endmodule
```



Behavioral HDL

Recall: Two Main Styles of HDL Implementation

- **Structural (Gate-Level)**
  - The module body contains gate-level description of the circuit
  - Describe how modules are interconnected
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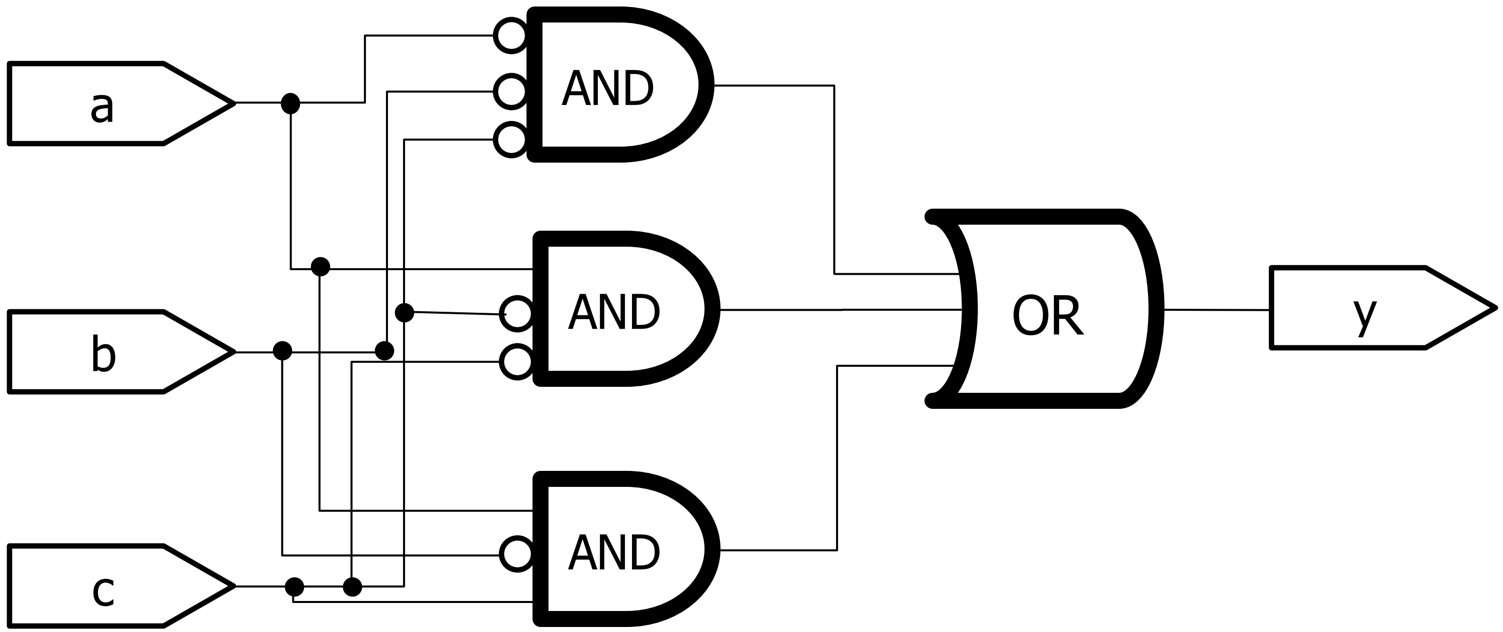
  - … and interconnections between those modules
  - Describes a hierarchy of modules defined as gates
- **Behavioral**
  - The module body contains functional description of the circuit
  - Contains logical and mathematical operators
  - **Level of abstraction is higher than gate-level**
    - Many possible gate-level realizations of a behavioral description
- **Many practical designs use a combination of both**

Behavioral HDL: Defining Functionality

```
module example (a, b, c, y);
       input a;
       input b;
       input c;
       output y;

// here comes the circuit description
assign y = ~a & ~b & ~c |
            a & ~b & ~c |
            a & ~b &  c;

endmodule
```

Behavioral HDL: Schematic View

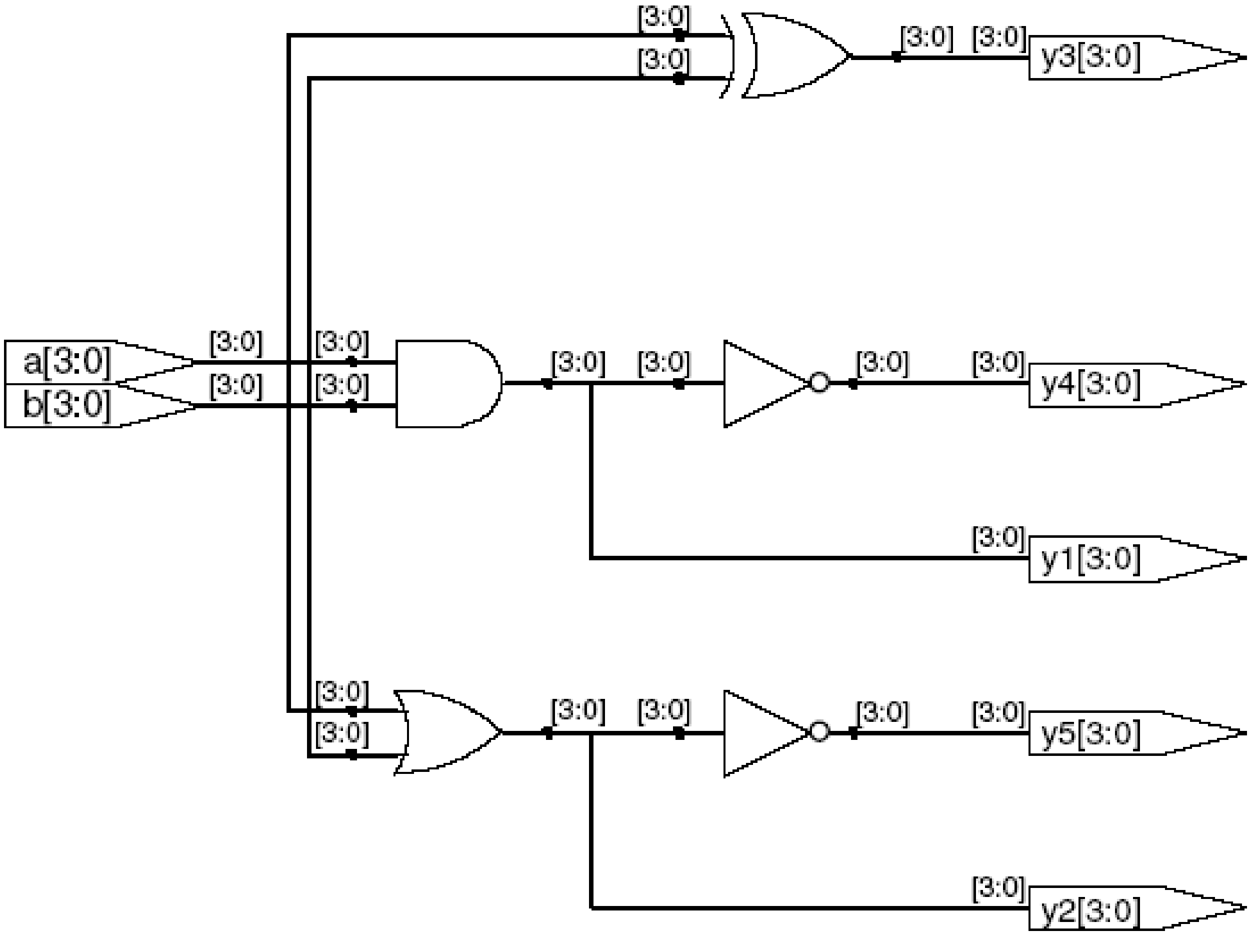
A behavioral implementation still models a hardware circuit!



Bitwise Operators in Behavioral Verilog

```
module gates(input  [3:0]  a, b,
             output [3:0] y1, y2, y3, y4, y5);

   /* Five different two-input logic
      gates acting on 4 bit buses */

   assign y1 = a & b;     // AND
   assign y2 = a | b;     // OR
   assign y3 = a ^ b;     // XOR
   assign y4 = ~(a & b); // NAND
   assign y5 = ~(a | b); // NOR

endmodule
```

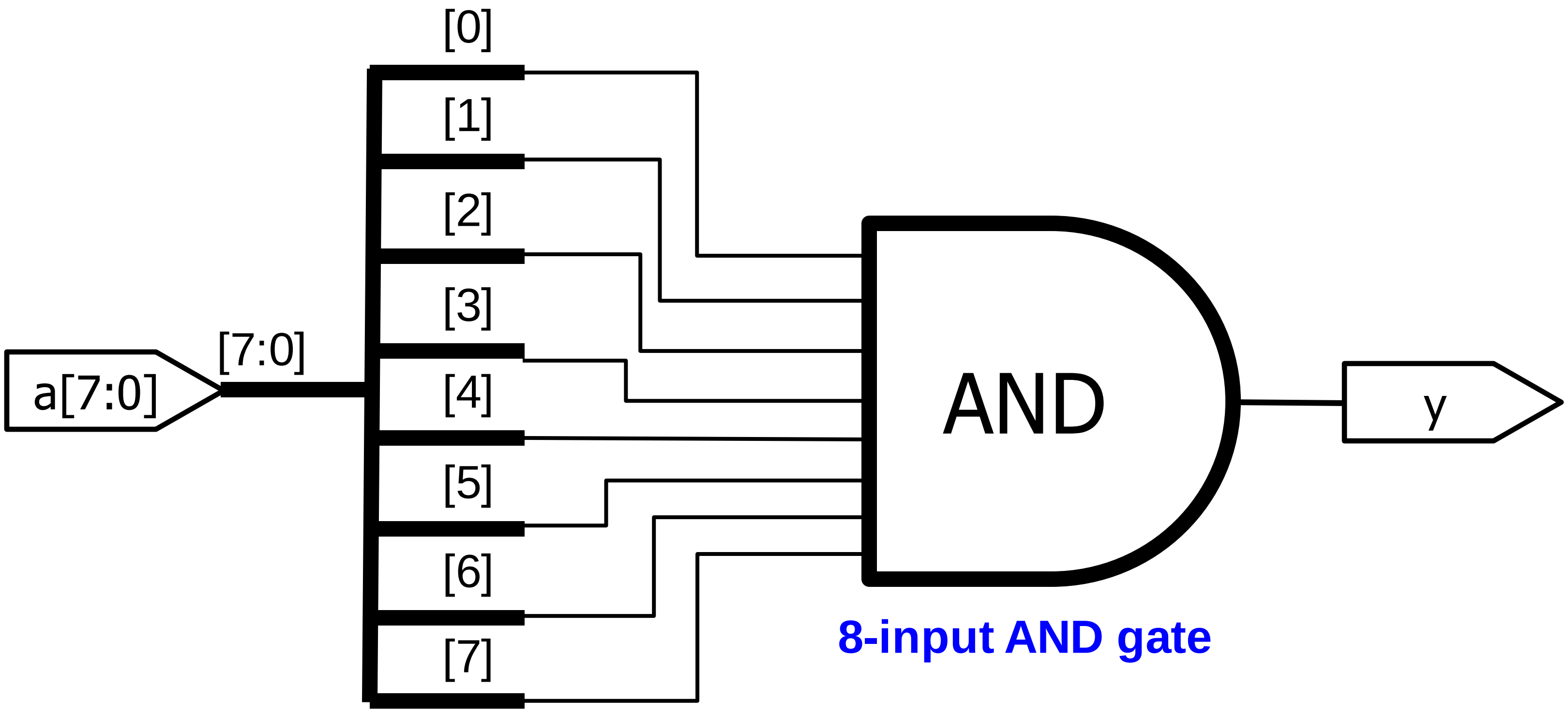
Bitwise Operators: Schematic View



Reduction Operators in Behavioral Verilog

```
module and8(input  [7:0] a,
            output       y);

   assign y = &a;

   // &a is much easier to write than
   // assign y = a[7] & a[6] & a[5] & a[4] &
   //            a[3] & a[2] & a[1] & a[0];

endmodule
```

Reduction Operators: Schematic View



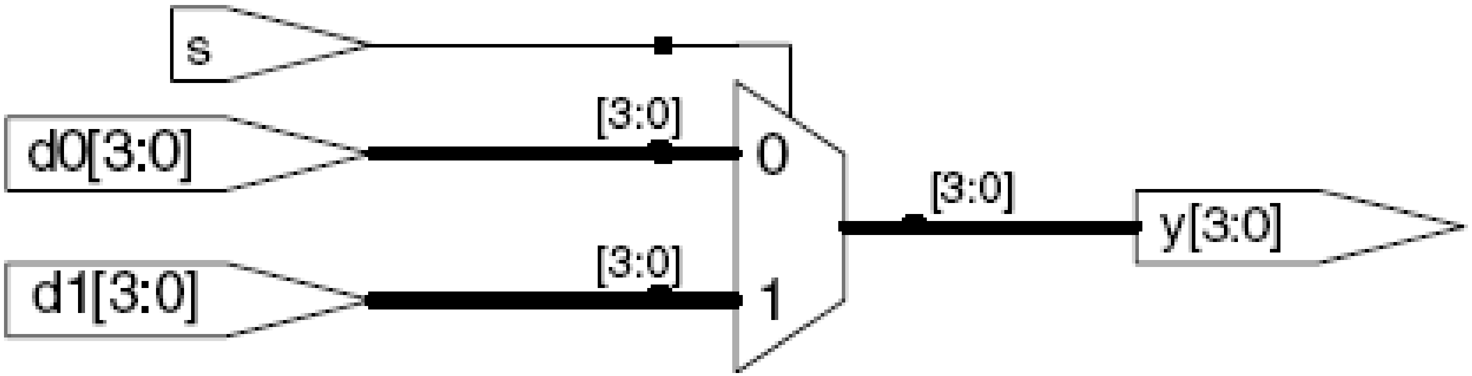
8-input AND gate

Conditional Assignment in Behavioral Verilog

```
module mux2(input  [3:0] d0, d1,
            input        s,
            output [3:0] y);

   assign y = s ? d1 : d0;
   // if (s) then y=d1 else y=d0;

endmodule
```

- ? : is also called a ternary operator as it operates on three inputs:
  - s
  - d1
  - d0

Conditional Assignment: Schematic View



More Complex Conditional Assignments

```
module mux4(input  [3:0] d0, d1, d2, d3
            input  [1:0] s,
            output [3:0] y);

   assign y = s[1] ? ( s[0] ? d3 : d2)
                   : ( s[0] ? d1 : d0);
  // if (s1) then
  //      if (s0) then y=d3 else y=d2
  // else
  //      if (s0) then y=d1 else y=d0

endmodule
```

Even More Complex Conditional Assignments

```
module mux4(input  [3:0] d0, d1, d2, d3
            input  [1:0] s,
            output [3:0] y);

   assign y = (s == 2'b11) ? d3 :
              (s == 2'b10) ? d2 :
              (s == 2'b01) ? d1 :
              d0;
// if      (s = "11" ) then y= d3
// else if (s = "10" ) then y= d2
// else if (s = "01" ) then y= d1
// else                     y= d0

endmodule
```

Precedence of Operations in Verilog

Highest

~	NOT
*, /, %	mult, div, mod
+, -	add,sub
<<, >>	shift
<<<, >>>	arithmetic shift
<, <=, >, >=	comparison
==, !=	equal, not equal
&, ~&	AND, NAND
^, ~^	XOR, XNOR
|, ~|	OR, NOR
?:	ternary operator

Lowest

How to Express Numbers ?

N'Bxx

8'b0000_0001

- (N) Number of bits
  - Expresses how many bits will be used to store the value

- (B) Base
  - Can be b (binary), h (hexadecimal), d (decimal), o (octal)

- (xx) Number
  - The value expressed in base
  - Can also have X (invalid) and Z (floating), as values
  - Underscore _ can be used to improve readability

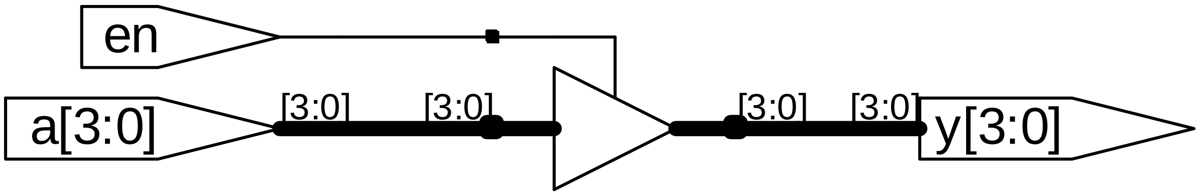
Number Representation in Verilog

Verilog	Stored Number	Verilog	Stored Number
4’b1001	1001	4’d5	0101
8’b1001	0000 1001	12’hFA3	1111 1010 0011
8’b0000_1001	0000 1001	8’o12	00 001 010
8’bxX0X1zZ1	XX0X 1ZZ1	4’h7	0111
‘b01	$\underbrace{\text{0000 .. 0001}}_{\text{32 bits (default)}}$	12’h0	0000 0000 0000

Reminder: Floating Signals (Z)

- Floating signal: Signal that is not driven by any circuit
  - Open circuit, floating wire
- Also known as: high impedance, hi-Z, tri-stated signals

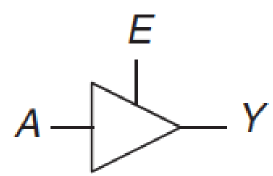
```
module tristate_buffer(input  [3:0] a,
                       input        en,
                       output [3:0] y);

   assign y = en ? a : 4'bz;

endmodule
```



Tri-State Buffer

- A tri-state buffer enables gating of different signals onto a wire

Tristate Buffer



E	A	Y
0	0	Z
0	1	Z
1	0	0
1	1	1

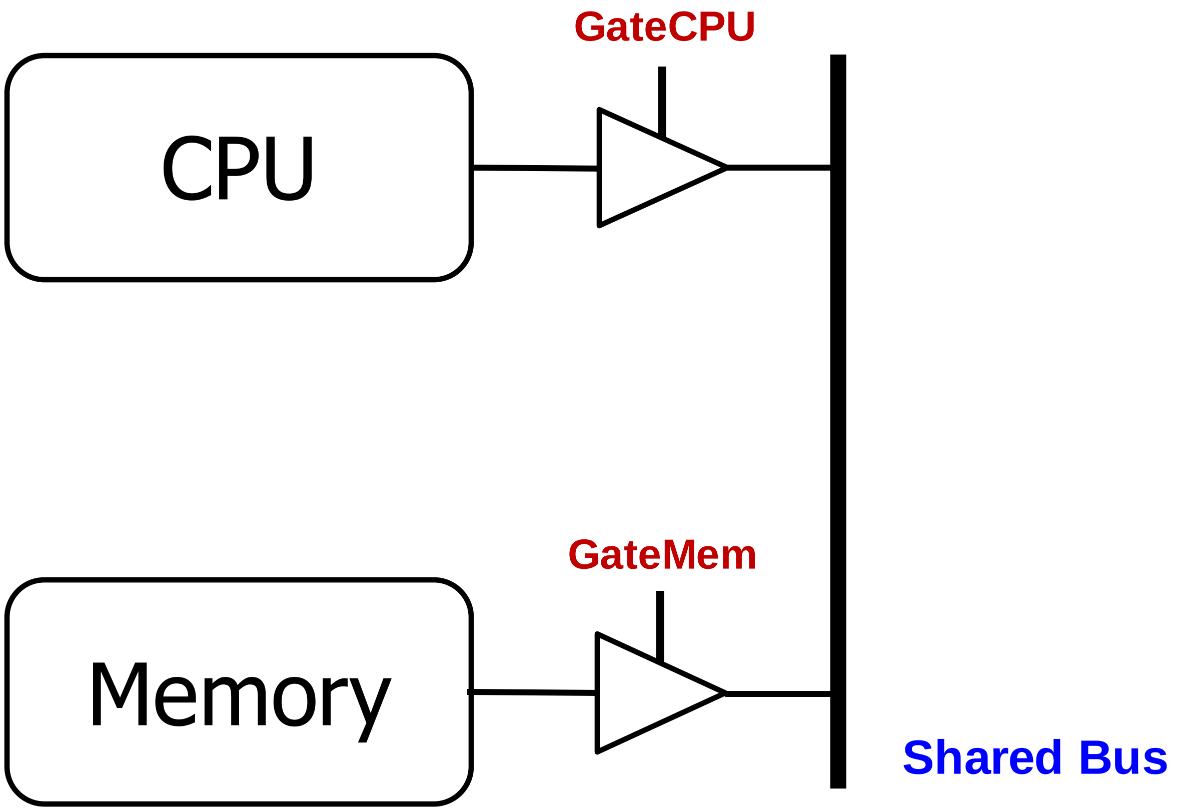
Figure 2.40 **Tristate buffer**

- Floating signal (Z): Signal that is not driven by any circuit
  - Open circuit, floating wire

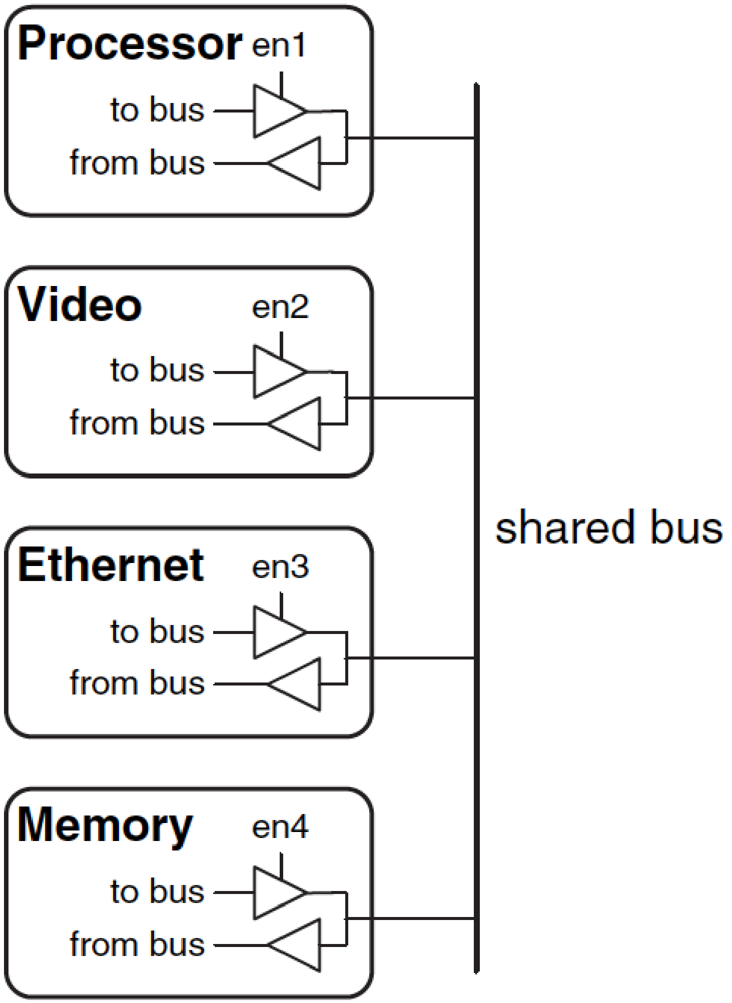
Example: Use of Tri-State Buffers

- Imagine a wire connecting the CPU and memory
  - At any time only the CPU or the memory can place a value on the wire, both not both
  - You can have two tri-state buffers: one driven by CPU, the other memory; and ensure at most one is enabled at any time

Example Design with Tri-State Buffers



Another Example



Truth Table for AND with Z and X

AND		A			
		0	**1**	**Z**	**X**
B	**0**	0	0	0	0
	1	0	1	X	X
	Z	0	X	X	X
	X	0	X	X	X

What Happens with HDL Code?

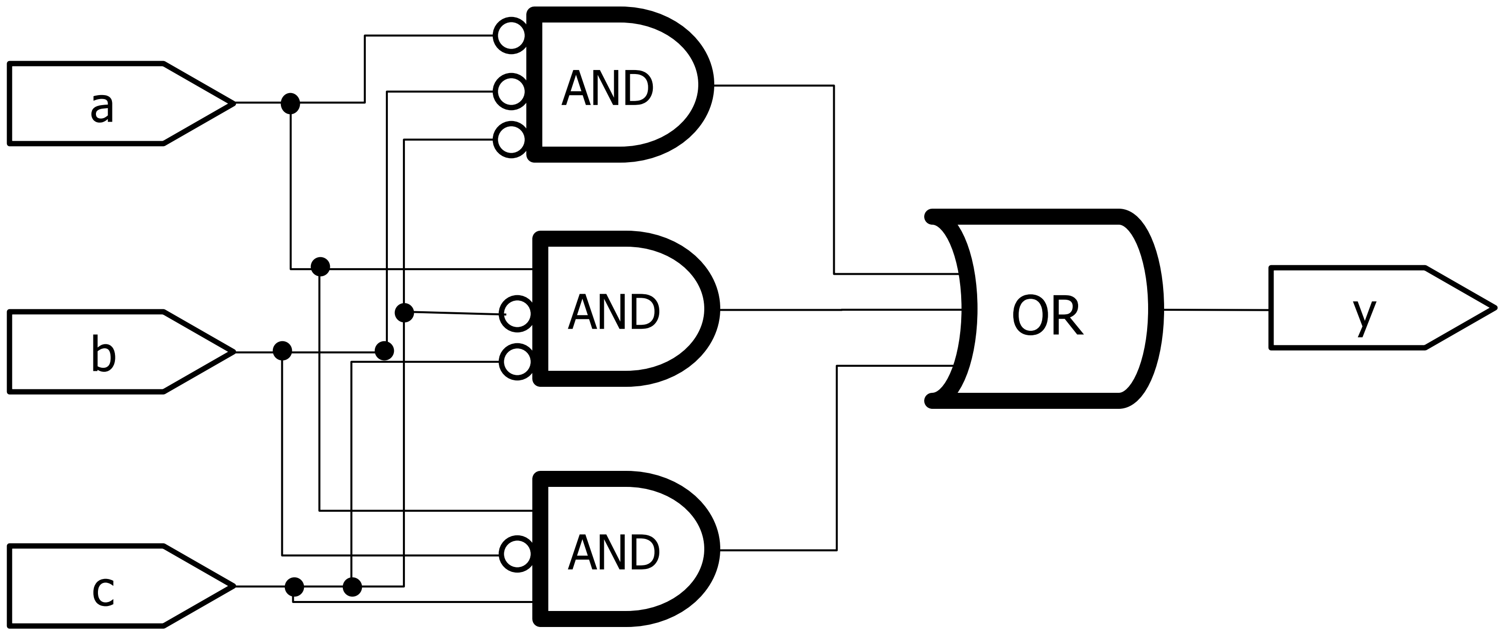
- **Synthesis**
  - Modern tools are able to map ***synthesizable*** *HDL code* into low-level *cell libraries → netlist describing gates and wires*
  - They can perform many optimizations
  - … however they can **not** guarantee that a solution is optimal
    - Mainly due to computationally expensive placement and routing algorithms
  - Most common way of Digital Design these days

- **Simulation**
  - Allows the behavior of the circuit to be verified without actually manufacturing the circuit
  - Simulators can work on *structural* or *behavioral* HDL

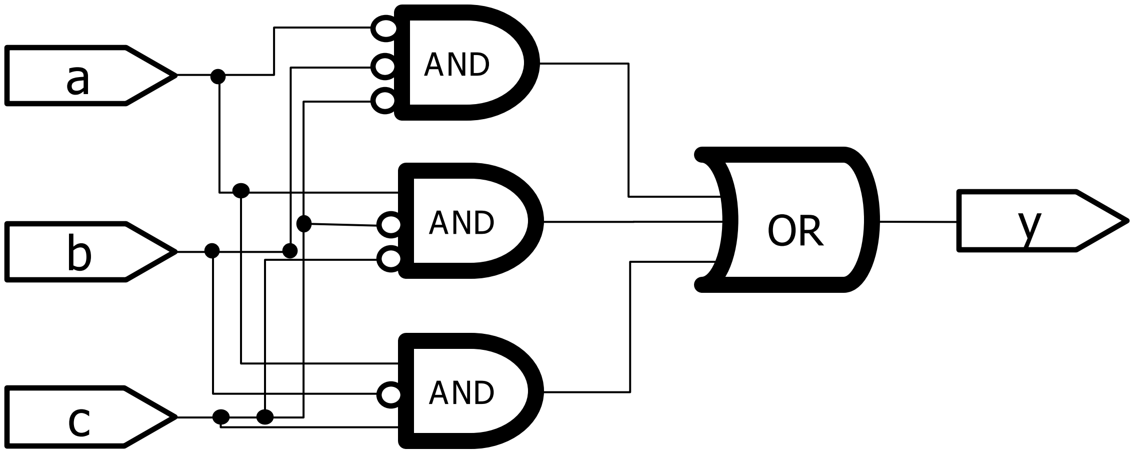
Recall This “example”

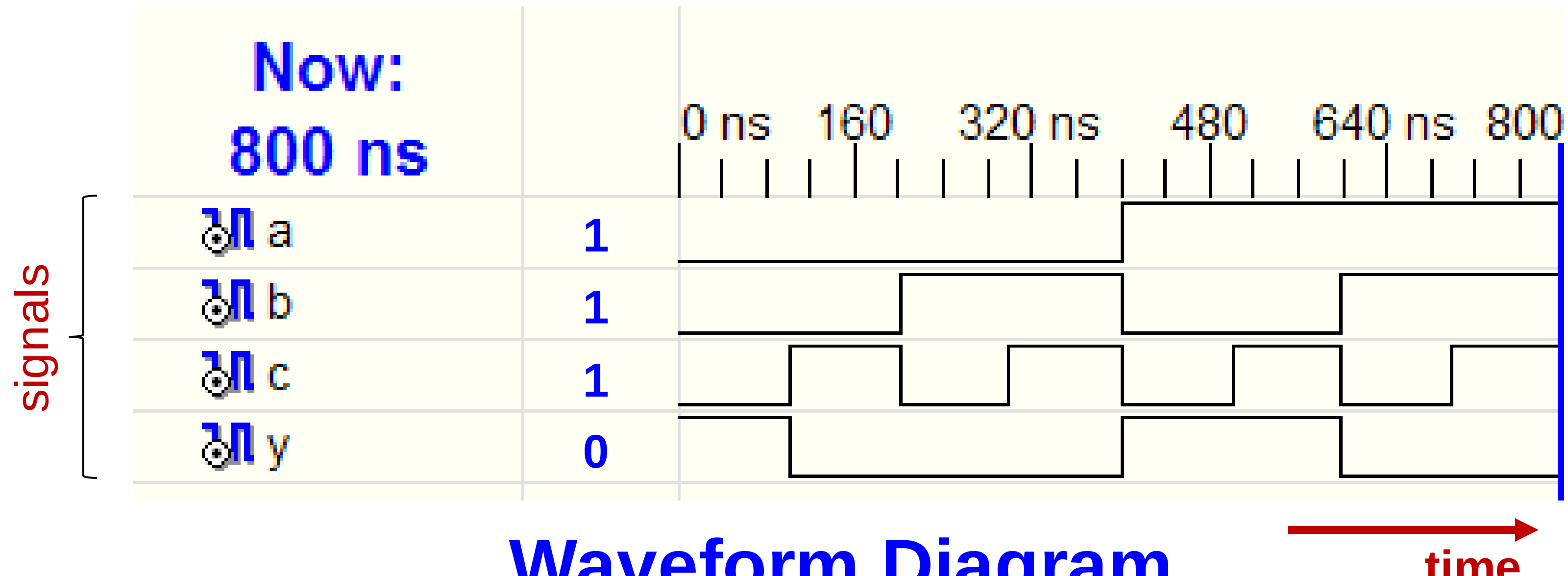
```
module example (a, b, c, y);
       input a;
       input b;
       input c;
       output y;

// here comes the circuit description
assign y = ~a & ~b & ~c |
            a & ~b & ~c |
            a & ~b &  c;

endmodule
```

Synthesizing the "example"



Simulating the "example"





Waveform Diagram

What We Have Seen So Far

- **Describing structural hierarchy with Verilog**
  - Instantiate modules in an other module
- **Describing functionality using behavioral modeling**

- **Writing simple logic equations**
  - We can write AND, OR, XOR, …
- **Multiplexer functionality**
  - If … then … else

- **We can describe constants**

- **But there is more…**

More Verilog Examples

- We can write Verilog code in many different ways

- Let's see how we can express the same functionality by developing Verilog code

  - At a low-level of abstraction
    - Poor readability
    - More optimization opportunities (especially for low-level tools)

  - At a high-level of abstraction
    - Better readability
    - Limited optimization opportunities

Comparing Two Numbers

- **Defining your own gates as new modules**

- We will use our gates to show the different ways of implementing a 4-bit comparator (equality checker)

An XNOR gate

```
module MyXnor (input A, B,
               output Z);

   assign Z = ~(A ^ B); //not XOR

endmodule
```

An AND gate

```
module MyAnd (input A, B,
              output Z);

   assign Z = A & B;    // AND

endmodule
```

Gate-Level Implementation

```
module compare (input a0, a1, a2, a3, b0, b1, b2, b3,
                output eq);
       wire c0, c1, c2, c3, c01, c23;

MyXnor i0 (.A(a0), .B(b0), .Z(c0) ); // XNOR
MyXnor i1 (.A(a1), .B(b1), .Z(c1) ); // XNOR
MyXnor i2 (.A(a2), .B(b2), .Z(c2) ); // XNOR
MyXnor i3 (.A(a3), .B(b3), .Z(c3) ); // XNOR
MyAnd haha (.A(c0), .B(c1), .Z(c01) ); // AND
MyAnd hoho (.A(c2), .B(c3), .Z(c23) ); // AND
MyAnd bubu (.A(c01), .B(c23), .Z(eq) ); // AND

endmodule
```

Using Logical Operators

```
module compare (input a0, a1, a2, a3, b0, b1, b2, b3,
                output eq);
       wire c0, c1, c2, c3, c01, c23;

MyXnor i0 (.A(a0), .B(b0), .Z(c0) ); // XNOR
MyXnor i1 (.A(a1), .B(b1), .Z(c1) ); // XNOR
MyXnor i2 (.A(a2), .B(b2), .Z(c2) ); // XNOR
MyXnor i3 (.A(a3), .B(b3), .Z(c3) ); // XNOR
assign c01 = c0 & c1;
assign c23 = c2 & c3;
assign eq  = c01 & c23;

endmodule
```

Eliminating Intermediate Signals

```
module compare (input a0, a1, a2, a3, b0, b1, b2, b3,
                output eq);
       wire c0, c1, c2, c3;

MyXnor i0 (.A(a0), .B(b0), .Z(c0) ); // XNOR
MyXnor i1 (.A(a1), .B(b1), .Z(c1) ); // XNOR
MyXnor i2 (.A(a2), .B(b2), .Z(c2) ); // XNOR
MyXnor i3 (.A(a3), .B(b3), .Z(c3) ); // XNOR
// assign c01 = c0 & c1;
// assign c23 = c2 & c3;
// assign eq  = c01 & c23;
assign eq  = c0 & c1 & c2 & c3;


endmodule
```

Multi-Bit Signals (Bus)

```
module compare (input [3:0] a, input [3:0] b,
                output eq);
       wire [3:0] c; // bus definition

MyXnor i0 (.A(a[0]), .B(b[0]), .Z(c[0]) ); // XNOR
MyXnor i1 (.A(a[1]), .B(b[1]), .Z(c[1]) ); // XNOR
MyXnor i2 (.A(a[2]), .B(b[2]), .Z(c[2]) ); // XNOR
MyXnor i3 (.A(a[3]), .B(b[3]), .Z(c[3]) ); // XNOR

assign eq  = &c; // short format


endmodule
```

Bitwise Operations

```
module compare (input [3:0] a, input [3:0] b,
                output eq);
       wire [3:0] c; // bus definition

// MyXnor i0 (.A(a[0]), .B(b[0]), .Z(c[0]) );
// MyXnor i1 (.A(a[1]), .B(b[1]), .Z(c[1]) );
// MyXnor i2 (.A(a[2]), .B(b[2]), .Z(c[2]) );
// MyXnor i3 (.A(a[3]), .B(b[3]), .Z(c[3]) );

assign c = ~(a ^ b); // XNOR

assign eq  = &c; // short format


endmodule
```

Highest Abstraction Level: Comparing Two Numbers

```
module compare (input [3:0] a, input [3:0] b,
                output eq);



// assign c = ~(a ^ b); // XNOR

// assign eq  = &c; // short format

assign eq = (a == b) ? 1 : 0; // really short



endmodule
```

Writing More Reusable Verilog Code

- We have a module that can compare two 4-bit numbers

- What if in the overall design we need to compare:
  - **5**-bit numbers?
  - **6**-bit numbers?
  - …
  - **N**-bit numbers?
  - Writing code for each case looks tedious

- What could be a better way?

Parameterized Modules

In Verilog, we can define module parameters

```
module mux2
  #(parameter width = 8)  // name and default value
   (input  [width-1:0] d0, d1,
    input              s,
    output [width-1:0] y);

  assign y = s ? d1 : d0;
endmodule
```

We can set the parameters to different values when instantiating the module

Instantiating Parameterized Modules

```
module mux2
  #(parameter width = 8)  // name and default value
   (input  [width-1:0] d0, d1,
    input              s,
    output [width-1:0] y);

  assign y = s ? d1 : d0;
endmodule
```

What About Timing?

- It is possible to define *timing relations* in Verilog. **BUT:**
  - These are ***ONLY*** for simulation
  - They ***CAN NOT*** be synthesized
  - They are used for *modeling delays* in a circuit

```
‘timescale 1ns/1ps
module simple (input a, output z1, z2);

assign #5 z1 = ~a; // inverted output after 5ns
assign #9 z2 = a;  // output after 9ns

endmodule
```

More to come later today!

Good Practices

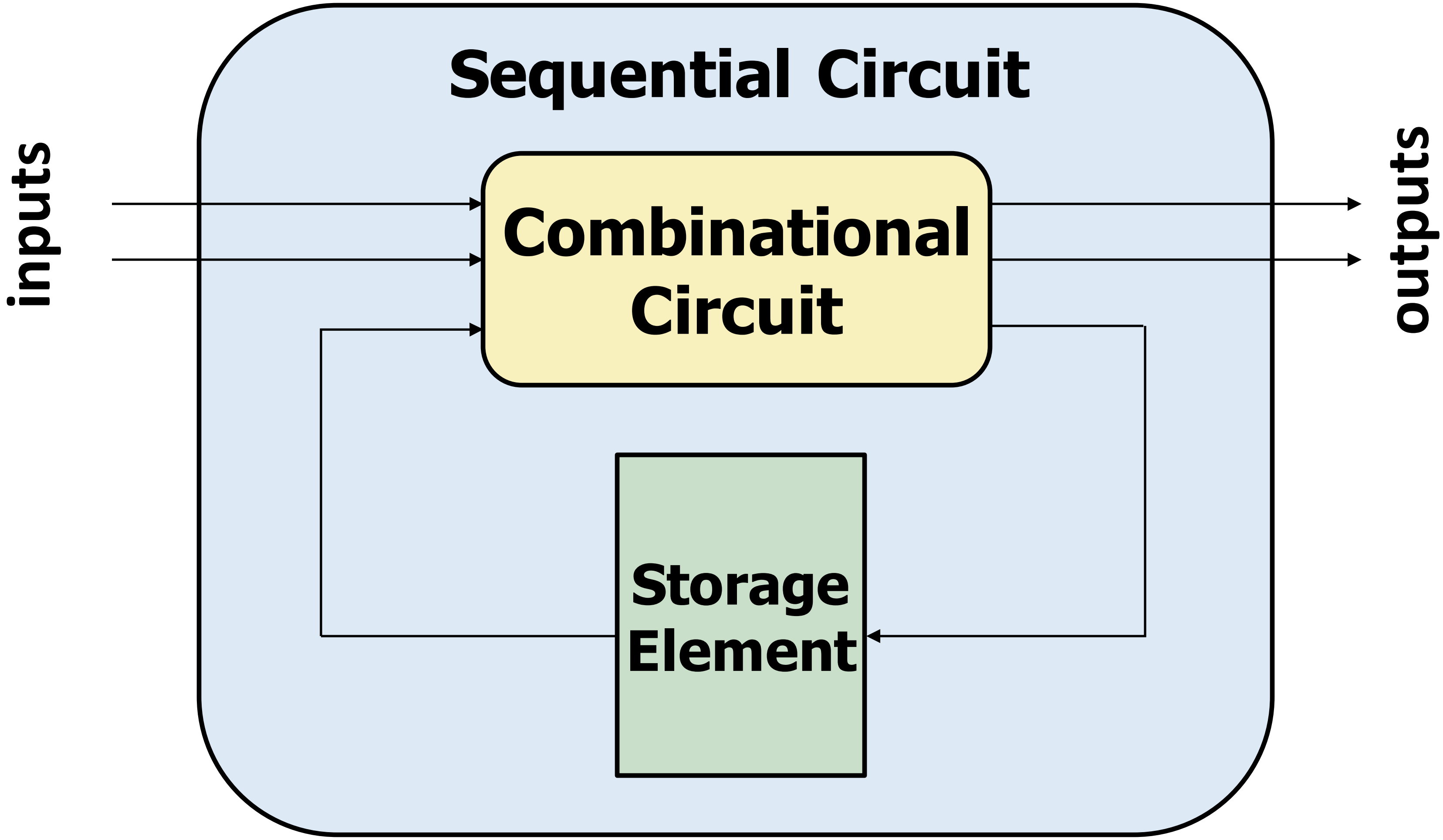
- Develop/use a consistent naming style
- Use MSB to LSB ordering for buses
  - Use "`a[31:0]`", **not** "`a[0:31]`"
- Define one module per file
  - Makes managing your design hierarchy easier
- Use a file name that equals module name
  - e.g., module TryThis is defined in a file called TryThis.v
- Always keep in mind that Verilog describes hardware

Summary (HDL for Combinational Logic)

- We have seen an overview of Verilog
- Discussed structural and behavioral modeling
- Studied combinational logic constructs

Implementing Sequential Logic Using Verilog

Combinational + Memory = Sequential



Sequential Logic in Verilog

- Define blocks that have memory
  - *Flip-Flops, Latches, Finite State Machines*

- Sequential Logic state transition is triggered by a "CLOCK" signal
  - Latches are sensitive to level of the signal
  - Flip-flops are sensitive to the transitioning of signal

- Combinational HDL constructs are **not** sufficient to express sequential logic
  - We need **new constructs**:
    - `always`
    - `posedge/negedge`

The "always" Block

```
always @ (sensitivity list)
      statement;
```

Whenever the event in the sensitivity list occurs,
the statement is executed

Example: D Flip-Flop

```
module flop(input            clk,
            input      [3:0] d,
            output reg [3:0] q);

  always @ (posedge clk)
    q <= d;                // pronounced “q gets d”

endmodule
```

- posedge defines a rising edge (transition from 0 to 1).
- Statement executed when the clk signal rises (posedge of clk)
- Once the clk signal rises: the value of d is copied to q

Example: D Flip-Flop

```
module flop(input            clk,
            input      [3:0] d,
            output reg [3:0] q);

  always @ (posedge clk)
    q <= d;               // pronounced “q gets d”

endmodule
```

- assign statement is **not** used within an always block
- <= describes a non-blocking assignment
  - We will see the difference between blocking assignment and non-blocking assignment soon

Example: D Flip-Flop

```
module flop(input            clk,
            input      [3:0] d,
            output reg [3:0] q);

  always @ (posedge clk)
    q <= d;                 // pronounced “q gets d”

endmodule
```

- Assigned variables need to be declared as reg
- The name reg does not necessarily mean that the value is a register (It could be, but it does not have to be)
- We will see examples later

Asynchronous and Synchronous Reset

- Reset signals are used to initialize the hardware to a known state
  - Usually activated at system start (on power up)

- **Asynchronous Reset**
  - The reset signal is sampled independent of the clock
  - Reset gets the highest priority
  - Sensitive to glitches, may have metastability issues
    - Will be discussed in Lecture 8

- **Synchronous Reset**
  - The reset signal is sampled with respect to the clock
  - The reset should be active long enough to get sampled at the clock edge
  - Results in completely synchronous circuit

D Flip-Flop with Asynchronous Reset

```
module flop_ar (input              clk,
                input              reset,
                input        [3:0] d,
                output reg   [3:0] q);

  always @ (posedge clk, negedge reset)
    begin
       if (reset == 0) q <= 0;   // when reset
       else            q <= d;   // when clk
    end
endmodule
```

- In this example: two events can trigger the process:
  - A ***rising edge*** on clk
  - A ***falling edge*** on reset

D Flip-Flop with Asynchronous Reset

```
module flop_ar (input              clk,
                input              reset,
                input        [3:0] d,
                output reg   [3:0] q);

  always @ (posedge clk, negedge reset)
    begin
      if (reset == 0) q <= 0;   // when reset
      else            q <= d;   // when clk
    end
endmodule
```

- For longer statements, a begin-end pair can be used
  - To improve readability
  - In this example, it was not necessary, but it is a good idea

D Flip-Flop with **Asynchronous** Reset

```
module flop_ar (input                 clk,
                input                 reset,
                input           [3:0] d,
                output reg      [3:0] q);

  always @ (posedge clk, negedge reset)
    begin
       if (reset == 0) q <= 0;  // when reset
       else            q <= d;  // when clk
    end
endmodule
```

- First reset is checked: if reset is 0, q is set to 0.
  - This is an asynchronous reset as the reset can happen independently of the clock (on the negative edge of reset signal)
- If there is no reset, then regular assignment takes effect

D Flip-Flop with **Synchronous** Reset

```
module flop_sr (input              clk,
                input              reset,
                input        [3:0] d,
                output reg   [3:0] q);

  always @ (posedge clk)
    begin
       if (reset == ‘0’) q <= 0;   // when reset
       else              q <= d;   // when clk
    end
endmodule
```

- The process is sensitive to only clock
  - Reset ***happens only*** when the ***clock rises***. This is a synchronous reset

D Flip-Flop with Enable and Reset

```
module flop_en_ar (input              clk,
                   input              reset,
                   input              en,
                   input        [3:0] d,
                   output reg   [3:0] q);

  always @ (posedge clk, negedge reset)
    begin
       if (reset == '0') q <= 0;    // when reset
       else if (en)      q <= d;    // when en AND clk
    end
endmodule
```

- A flip-flop with **enable** and **reset**
  - Note that the en signal is ***not*** in the *sensitivity list*
- q gets d only when clk is rising **and** en is 1

Example: D Latch

```
module latch (input             clk,
              input       [3:0] d,
              output reg  [3:0] q);

  always @ (clk, d)
    if (clk) q <= d;       // latch is transparent when
                           // clock is 1
endmodule
```

Summary: Sequential Statements So Far

- Sequential statements are within an `always` block

- The sequential block is triggered with a change in the sensitivity list

- Signals assigned within an **always** must be declared as `reg`

- We use `<=` for (non-blocking) assignments and do not use `assign` within the always block.

Basics of **always** Blocks

```
module example (input             clk,
                input       [3:0] d,
                output reg  [3:0] q);

  wire [3:0] normal;           // standard wire
  reg  [3:0] special;          // assigned in always

  always @ (posedge clk)
    special <= d;              // first FF array

  assign normal = ~ special; // simple assignment

  always @ (posedge clk)
    q <= normal;               // second FF array
endmodule
```

You can have as many always blocks as needed

Assignment to the same signal in different always blocks is not allowed!

Why Does an **always** Block Remember?

```
module flop (input            clk,
             input      [3:0] d,
             output reg [3:0] q);

  always @ (posedge clk)
    begin
        q <= d;    // when clk rises copy d to q
    end
endmodule
```

- This statement describes what happens to signal q
- … but what happens when the clock is not rising?
- The value of q is preserved (remembered)

An **always** Block Does **NOT** Always Remember

```
module comb (input              inv,
             input        [3:0] data,
             output reg [3:0] result);

  always @ (inv, data)        // trigger with inv, data
    if (inv) result <= ~data;// result is inverted data
    else     result <= data; // result is data

endmodule
```

- This statement describes what happens to signal result
  - When `inv` is 1, result is ~`data`
  - When `inv` is not 1, result is `data`
- The circuit is combinational (no memory)
  - result is assigned a value in all cases of the if .. else block, always

always Blocks for Combinational Circuits

- An always block defines combinational logic if:
  - All outputs are always (**continuously**) updated

1. All right-hand side signals are in the sensitivity list
   - You can use `always @*` for short
2. All left-hand side signals get assigned in every possible condition of if .. else and case blocks

- It is easy to make mistakes and unintentionally describe memorizing elements (latches)
  - Vivado will most likely warn you. Make sure you check the warning messages

- **Always** blocks allow powerful combinational logic statements
  - `if .. else`
  - `case`

Sequential or Combinational?

```
wire enable, data;
reg out_a, out_b;

always @ (*) begin
        out_a = 1'b0;
        if(enable) begin
                out_a = data;
                out_b = data;
        end
end
```

No assignment for ~enable

```
wire enable, data;
reg out_a, out_b;

always @ (data) begin
        out_a = 1'b0;
        out_b = 1'b0;
        if(enable) begin
                out_a = data;
                out_b = data;
        end
end
```

Not in the sensitivity list

Sequential

Sequential

The **always** Block is **NOT** Always Practical/Nice

```
reg  [31:0] result;
wire [31:0] a, b, comb;
wire        sel,

always @ (a, b, sel)    // trigger with a, b, sel
   if (sel) result <= a; // result is a
   else     result <= b; // result is b

assign comb = sel ? a : b;
```

- Both statements describe the **same** multiplexer
- In this case, the always block is more work

always Block for Case Statements (Handy!)

```
module sevensegment (input      [3:0] data,
                     output reg [6:0] segments);

  always @ ( * )                   // * is short for all signals
    case (data)                    // case statement
      4'd0: segments = 7'b111_1110;  // when data is 0
      4'd1: segments = 7'b011_0000;  // when data is 1
      4'd2: segments = 7'b110_1101;
      4'd3: segments = 7'b111_1001;
      4'd4: segments = 7'b011_0011;
      4'd5: segments = 7'b101_1011;
      // etc etc
      default: segments = 7'b000_0000; // required
    endcase

endmodule
```

Summary: **always** Block

- `if .. else` can only be used in always blocks

- The always block is combinational only if all regs within the block are always assigned to a signal
  - Use the `default` case to make sure you do not forget an unimplemented case, which may otherwise result in a latch

- Use `casex` statement to be able to check for don't cares

Non-Blocking and Blocking Assignments

Non-blocking (<=)

```
always @ (a)
begin
   a <= 2'b01;
   b <= a;
// all assignments are made here
// b is not (yet) 2'b01
end
```

- All assignments are made at the end of the block
- All assignments are made in parallel, process flow is not-blocked

Blocking (=)

```
always @ (a)
begin
   a = 2'b01;
// a is 2'b01
   b = a;
// b is now 2'b01 as well
end
```

- Each assignment is made immediately
- Process waits until the first assignment is complete, it blocks progress

Example: Blocking Assignment

- Assume all inputs are initially '0'

```
always @ ( * )
  begin
    p    = a ^ b ;          // p    = 0  1
    g    = a & b ;          // g    = 0  0
    s    = p ^ cin ;        // s    = 0  1
    cout = g | (p & cin) ; // cout = 0  0
  end
```

- If a changes to '1'
  - All values are updated in order

The Same Example: Non-Blocking Assignment

- Assume all inputs are initially '0'

```
always @ ( * )
  begin
    p    <= a ^ b ;          // p    = 0  1
    g    <= a & b ;          // g    = 0  0
    s    <= p ^ cin ;        // s    = 0  0
    cout <= g | (p & cin) ; // cout = 0  0
  end
```

- If a changes to '1'
  - All assignments are concurrent
  - When s is being assigned, p is still 0

The Same Example: Non-Blocking Assignment

- After the first iteration, p has changed to '1' as well

```
always @ ( * )
  begin
    p    <= a ^ b ;          // p    = 1  1
    g    <= a & b ;          // g    = 0  0
    s    <= p ^ cin ;        // s    = 0  1
    cout <= g | (p & cin) ; // cout = 0  0
  end
```

- Since there is a change in p, the process triggers again
- This time s is calculated with p=1

Rules for Signal Assignment

- Use `always @(posedge clk)` and non-blocking assignments (`<=`) to model synchronous sequential logic

```
always @ (posedge clk)
    q <= d; // non-blocking
```

- Use continuous assignments (`assign`) to model simple combinational logic

```
assign y = a & b;
```

Rules for Signal Assignment (Cont.)

- Use `always @ (*)` and blocking assignments (`=`) to model more complicated combinational logic.

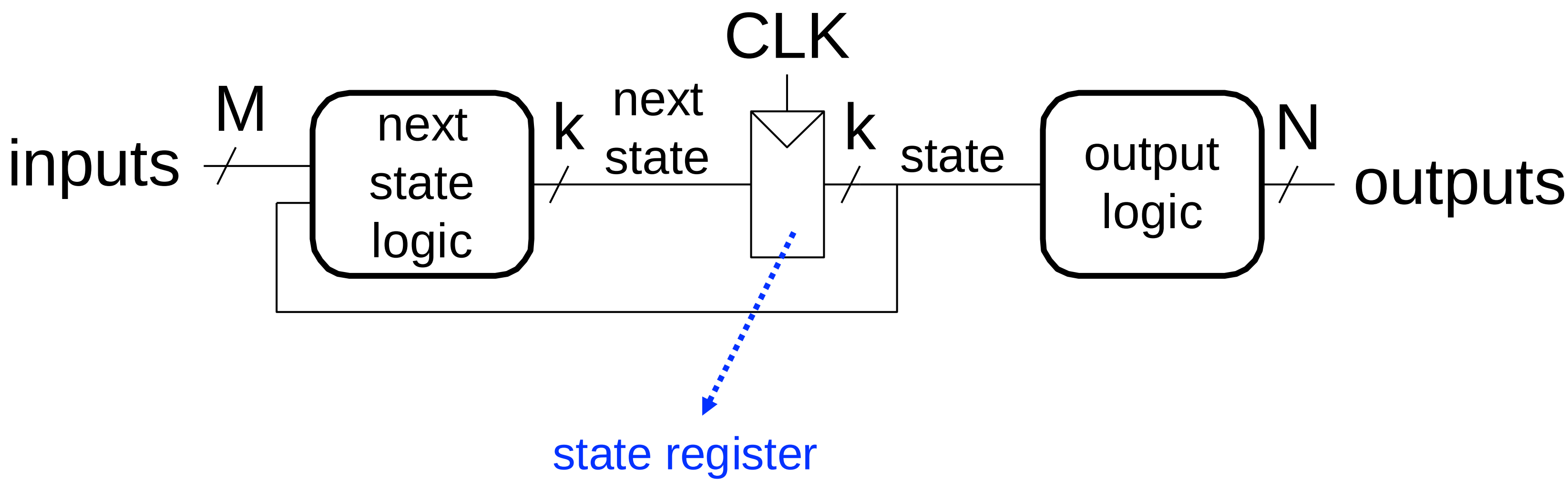
- You cannot make assignments to the same signal in more than one always block or in a *continuous assignment*

```
always @ (*)
        a = b;

always @ (*)
        a = c;
```

```
always @ (*)
        a = b;

assign a = c;
```

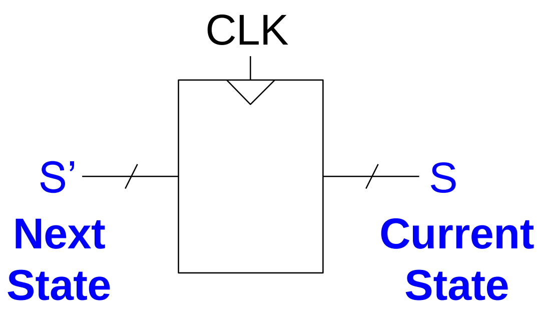
Recall: Finite State Machines (FSMs)

- Each FSM consists of three separate parts:
  - next state logic
  - state register
  - output logic

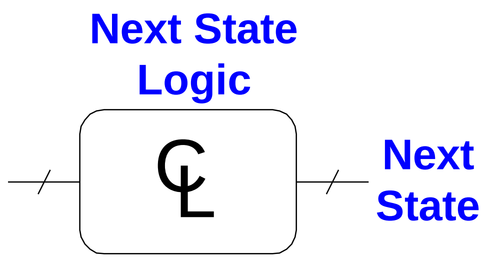


Recall: Finite State Machines (FSMs) Comprise

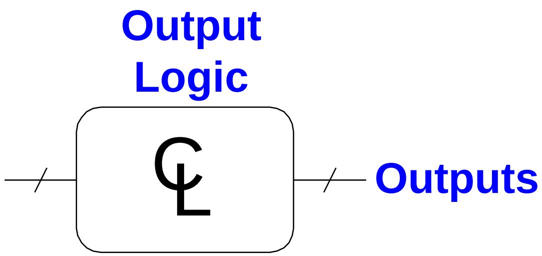
- **Sequential circuits**
  - State register(s)
    - Store the current state and
    - Load the next state at the clock edge



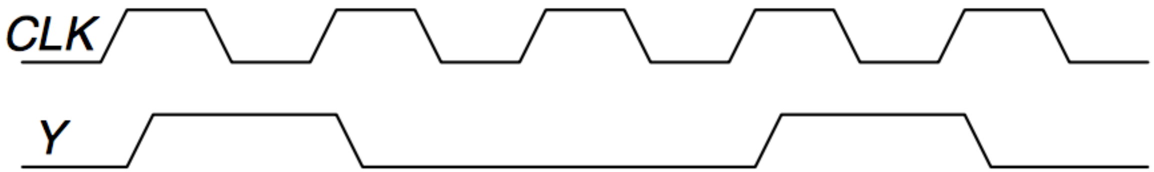
- **Combinational Circuits**
  - Next state logic
    - Determines what the next state will be



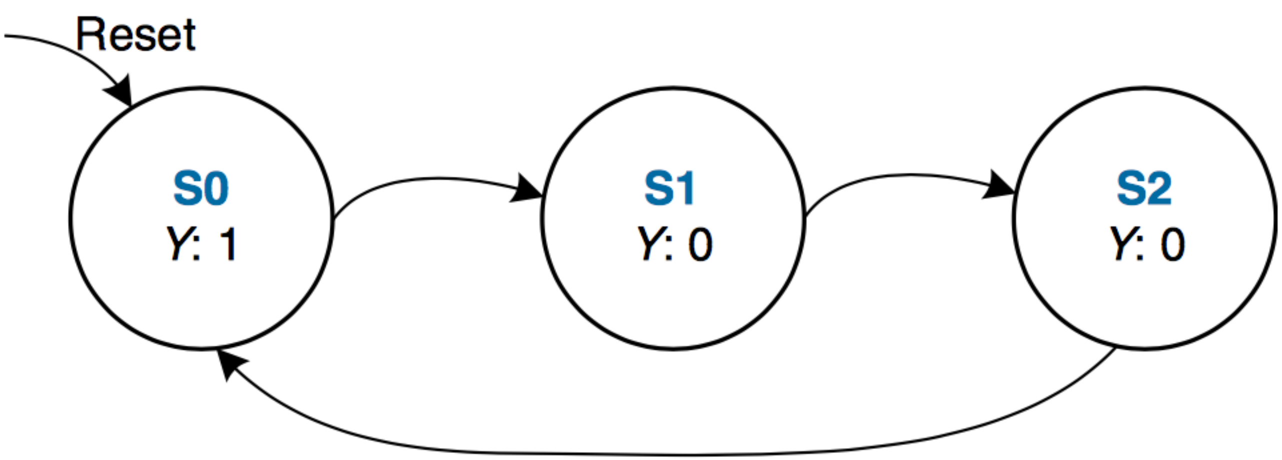
  - Output logic
    - Generates the outputs



FSM Example 1: Divide the Clock Frequency by 3



The output *Y* is HIGH for **one clock cycle out of every *3*.** In other words, the output **divides the frequency of the clock by *3***.

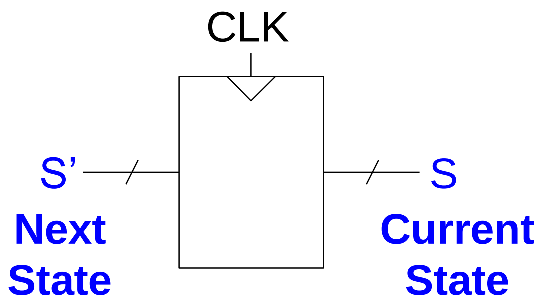


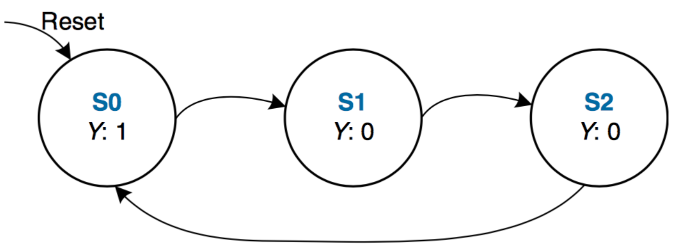
Implementing FSM Example 1: Definitions

```
module divideby3FSM (input clk,
                     input reset,
                     output q);

   reg  [1:0] state, nextstate;

   parameter S0 = 2'b00;
   parameter S1 = 2'b01;
   parameter S2 = 2'b10;
```

- We define `state` and `nextstate` as 2-bit **reg**
- The parameter descriptions are optional, it makes reading easier

Implementing FSM Example 1: State Register

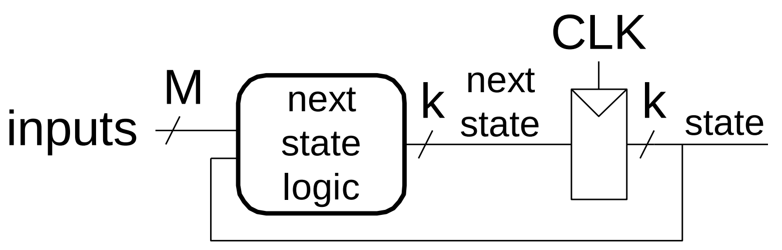


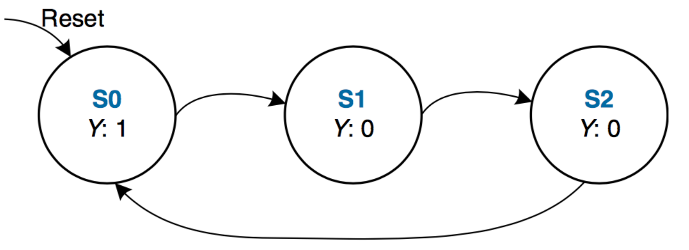



```
// state register
   always @ (posedge clk, posedge reset)
      if (reset) state <= S0;
      else       state <= nextstate;
```

- This part defines the state register (memorizing process)
- Sensitive to only clk, reset
- In this example, reset is active when it is '1' (active-high)

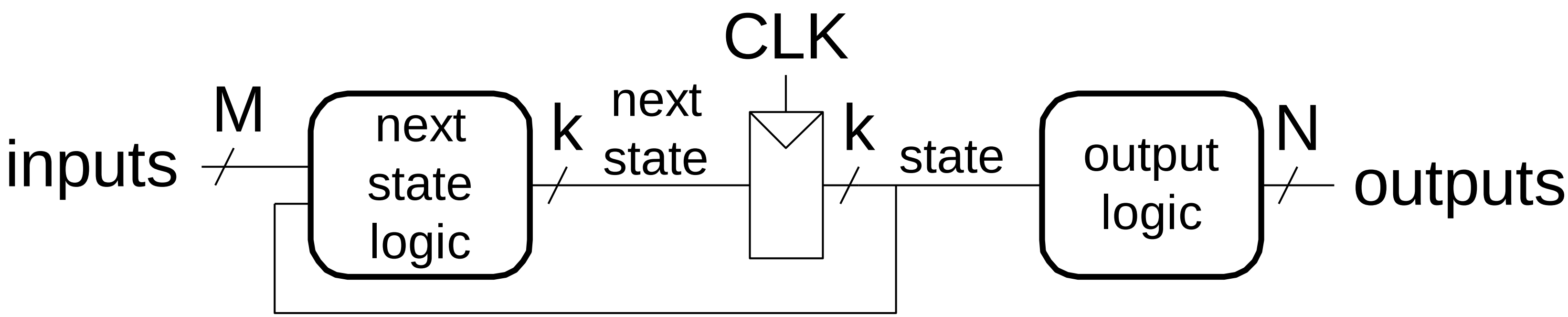
Implementing FSM Example 1: Next State Logic

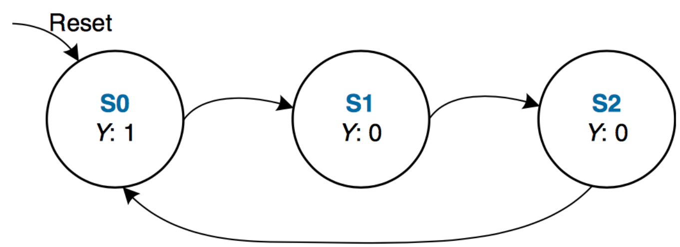





```
// next state logic
  always @ (*)
     case (state)
        S0:      nextstate = S1;
        S1:      nextstate = S2;
        S2:      nextstate = S0;
        default: nextstate = S0;
     endcase
```

Implementing FSM Example 1: Output Logic






```
// output logic
   assign q = (state == S0);
```

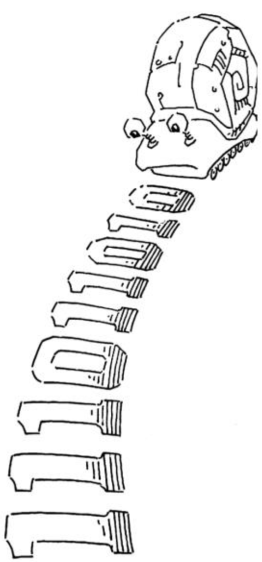
- In this example, output depends only on state
  - **Moore type FSM**

Implementation of FSM Example 1

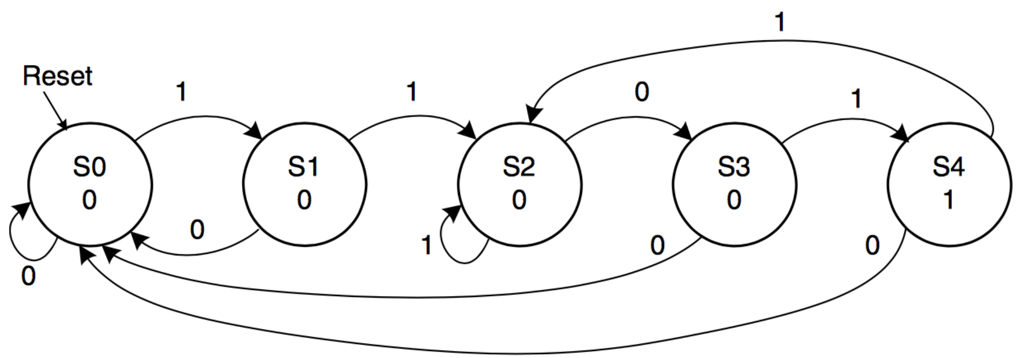
```
module divideby3FSM (input clk, input reset, output q);
   reg  [1:0] state, nextstate;

   parameter S0 = 2'b00; parameter S1 = 2'b01; parameter S2 = 2'b10;

   always @ (posedge clk, posedge reset) // state register
      if (reset) state <= S0;
      else       state <= nextstate;

   always @ (*)                          // next state logic
      case (state)
         S0:      nextstate = S1;
         S1:      nextstate = S2;
         S2:      nextstate = S0;
         default: nextstate = S0;
      endcase
   assign q = (state == S0);            // output logic
endmodule
```

FSM Example 2: Smiling Snail

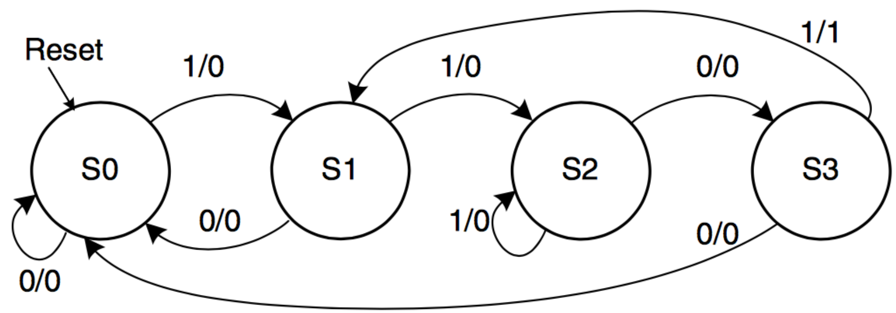
- Alyssa P. Hacker has a snail that crawls down a paper tape with 1's and 0's on it
- The snail smiles whenever the last four digits it has crawled over are 1101
- Design Moore and Mealy FSMs of the snail's brain



Moore

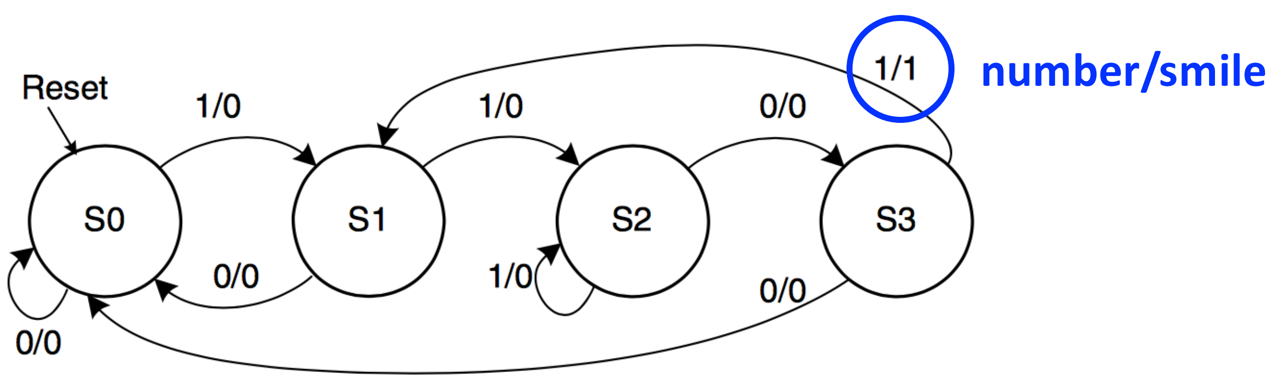


Mealy



Implementing FSM Example 2: Definitions

```
module SmilingSnail (input clk,
                     input reset,
                     input number,
                     output smile);

   reg  [1:0] state, nextstate;

   parameter S0 = 2'b00;
   parameter S1 = 2'b01;
   parameter S2 = 2'b10;
   parameter S3 = 2’b11;
```



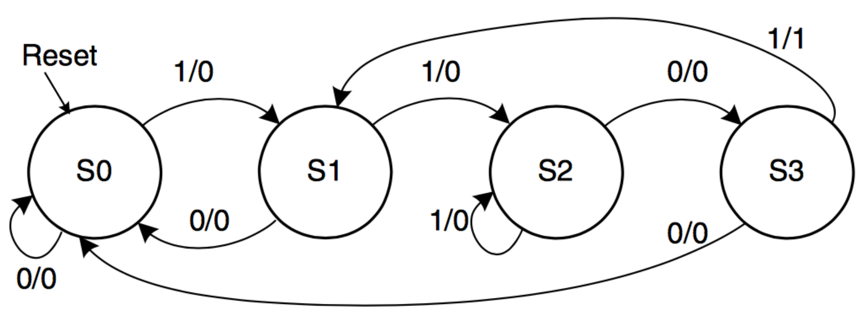
Implementing FSM Example 2: State Register

```
// state register
   always @ (posedge clk, posedge reset)
      if (reset) state <= S0;
      else       state <= nextstate;
```

- This part defines the state register (memorizing process)
- Sensitive to only clk, reset
- In this example reset is active when ‘1’ (active-high)

Implementing FSM Example 2: Next State Logic

```
// next state logic
  always @ (*)
     case (state)
        S0: if (number) nextstate = S1;
            else   nextstate = S0;
        S1: if (number) nextstate = S2;
            else   nextstate = S0;
        S2: if (number) nextstate = S2;
            else   nextstate = S3;
        S3: if (number) nextstate = S1;
            else   nextstate = S0;
        default:   nextstate = S0;
     endcase
```



Implementing FSM Example 2: Output Logic

```
// output logic
   assign smile = (number & state == S3);
```

- In this example, output depends on state and input
  - **Mealy type FSM**
- We used a simple combinational assignment

Implementation of FSM Example 2

```
module SmilingSnail (input clk,
                     input reset,
                     input number,
                     output smile);

   reg  [1:0] state, nextstate;

   parameter S0 = 2'b00;
   parameter S1 = 2'b01;
   parameter S2 = 2'b10;
   parameter S3 = 2’b11;

   // state register
   always @ (posedge clk, posedge
reset)
      if (reset) state <= S0;
      else       state <= nextstate;
```

```
always @ (*) // next state logic
      case (state)
         S0: if (number)
                  nextstate = S1;
             else nextstate = S0;
         S1: if (number)
                  nextstate = S2;
             else nextstate = S0;
         S2: if (number)
                  nextstate = S2;
             else nextstate = S3;
         S3: if (number)
                  nextstate = S1;
             else nextstate = S0;
         default: nextstate = S0;
      endcase
   // output logic
assign smile = (number & state==S3);

endmodule
```

What Did We Learn?

- Basics of describing sequential circuits in Verilog

- The always statement
  - Needed for defining memorizing elements (flip-flops, latches)
  - Can also be used to define combinational circuits

- Blocking vs Non-blocking statements
  - = assigns the value immediately
  - <= assigns the value at the end of the block

- Describing FSMs in Verilog
  - Next state logic
  - State assignment
  - Output logic

Next Lecture:

Timing and Verification

Digital Design & Computer Arch.

Lecture 7: Hardware Description Languages and Verilog

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ETH Zürich

Spring 2021

18 March 2021

Logic Simplification: Karnaugh Maps (K-Maps)

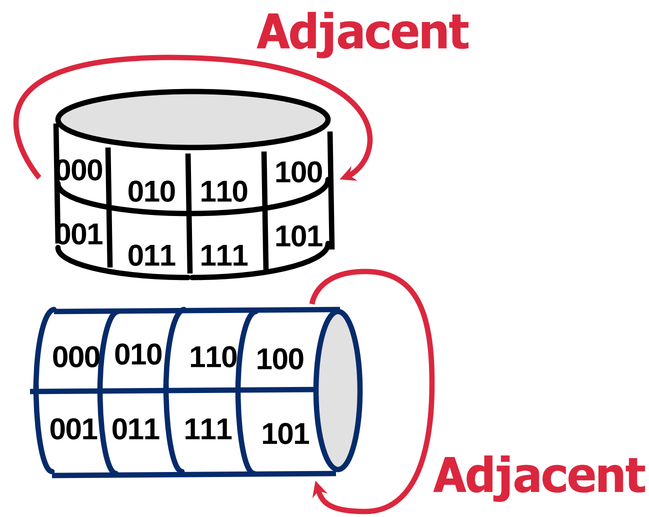
Karnaugh Maps are Fun…

- A pictorial way of minimizing circuits by visualizing opportunities for simplification
- They are for you to **study on your own**…

- See Backup Slides
- Read H&H Section 2.7
- Watch videos of Lectures 5 and 6 from 2019 DDCA course:
  - https://youtu.be/0ks0PeaOUjE?list=PL5Q2soXY2Zi8J58xLKBNFQFHRO3GrXxA9&t=4570
  - https://youtu.be/ozs18ARNG6s?list=PL5Q2soXY2Zi8J58xLKBNFQFHRO3GrXxA9&t=220

Karnaugh Map Methods

A \ BC	00	01	11	10
0	000	001	011	010
1	100	101	111	110



K-map adjacencies go "around the edges"

Wrap around from first to last column

Wrap around from top row to bottom row

Backup Slides on Karnaugh Maps (K-Maps)

Complex Cases

- One example

$$\boldsymbol{Cout = \bar{A}BC + A\bar{B}C + AB\bar{C} + ABC}$$

- Problem
  - Easy to see how to apply Uniting Theorem…
  - Hard to know if you applied it in all the right places…
  - …especially in a function of many more variables

- Question
  - Is there an easier way to find potential simplifications?
  - i.e., potential applications of Uniting Theorem…?

- Answer
  - Need an intrinsically *geometric* representation for Boolean f( )
  - Something we can draw, see…

Karnaugh Map

- Karnaugh Map (K-map) method
  - K-map is an alternative method of representing the truth table that helps visualize adjacencies in up to 6 dimensions
  - Physical adjacency ↔ Logical adjacency

2-variable K-map

$A \backslash B$	0	1
0	00	01
1	10	11

3-variable K-map

$A \backslash BC$	00	01	11	10
0	000	001	011	010
1	100	101	111	110

4-variable K-map

$AB \backslash CD$	00	01	11	10
00	0000	0001	0011	0010
01	0100	0101	0111	0110
11	1100	1101	1111	1110
10	1000	1001	1011	1010

Numbering Scheme: **00, 01, 11, 10 is called a "Gray Code" — only a *single bit (variable) changes* from one code word and the next code word**

K-map Cover - 4 Input Variables

$AB \backslash CD$	00	01	11	10
00	1	0	0	1
01	0	1	0	0
11	1	1	1	1
10	1	1	1	1

$$\mathrm{F(A,B,C,D)} = \sum m(0,2,5,8,9,10,11,12,13,14,15)$$

$$\mathrm{F} = \mathrm{A} + \overline{B}\,\overline{D} + \mathrm{B}\overline{C}D$$

Strategy for "circling" rectangles on Kmap:

Biggest "oops!" that people forget:

Logic Minimization Using K-Maps

- Very simple guideline:
  - Circle all the rectangular blocks of 1's in the map, using the fewest possible number of circles
    - Each circle should be as large as possible
  - Read off the implicants that were circled

- More formally:
  - A Boolean equation is minimized when it is written as a sum of the fewest number of prime implicants
  - Each circle on the K-map represents an implicant
  - The largest possible circles are prime implicants

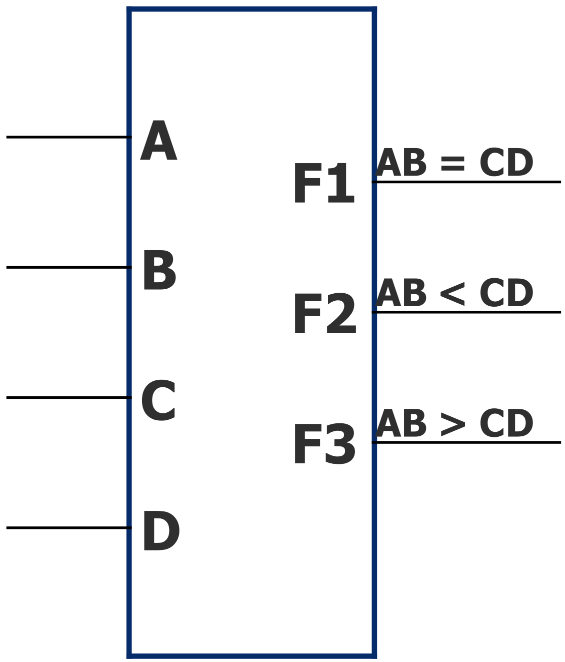
K-map Rules

- **What can be legally combined (circled) in the K-map?**
  - Rectangular groups of size 2^k for any integer k
  - Each cell has the same value (1, for now)
  - All values must be adjacent
    - Wrap-around edge is okay

- **How does a group become a term in an expression?**
  - Determine which literals are constant, and which vary across group
  - Eliminate varying literals, then AND the constant literals
    - constant 1 → use **X**, constant 0 → use $\overline{X}$

- **What is a good solution?**
  - Biggest groupings → eliminate more variables (literals) in each term
  - Fewest groupings → fewer terms (gates) all together
  - OR together all AND terms you create from individual groups

K-map Example: Two-bit Comparator



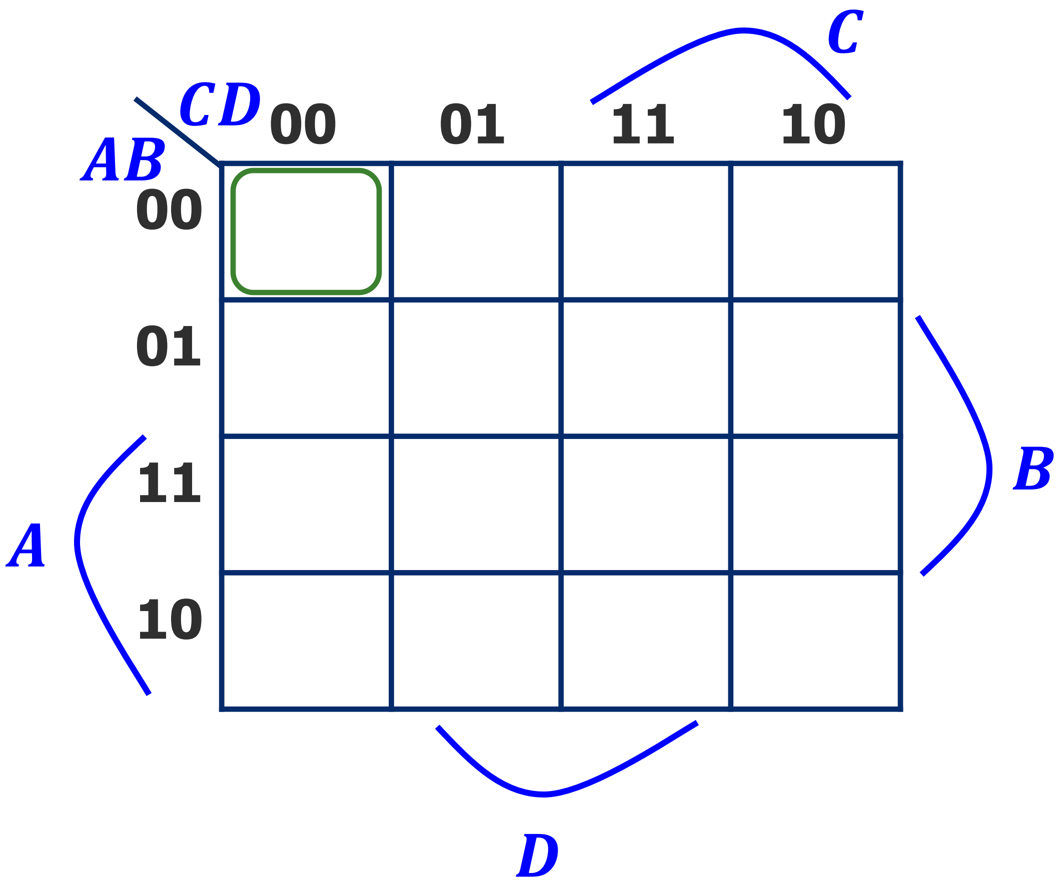
A	B	C	D	F1	F2	F3
0	0	0	0	1	0	0
0	0	0	1	0	1	0
0	0	1	0	0	1	0
0	0	1	1	0	1	0
0	1	0	0	0	0	1
0	1	0	1	1	0	0
0	1	1	0	0	1	0
0	1	1	1	0	1	0
1	0	0	0	0	0	1
1	0	0	1	0	0	1
1	0	1	0	1	0	0
1	0	1	1	0	1	0
1	1	0	0	0	0	1
1	1	0	1	0	0	1
1	1	1	0	0	0	1
1	1	1	1	1	0	0

Design Approach:

Write a 4-Variable K-map for each of the 3 output functions

K-map Example: Two-bit Comparator (2)

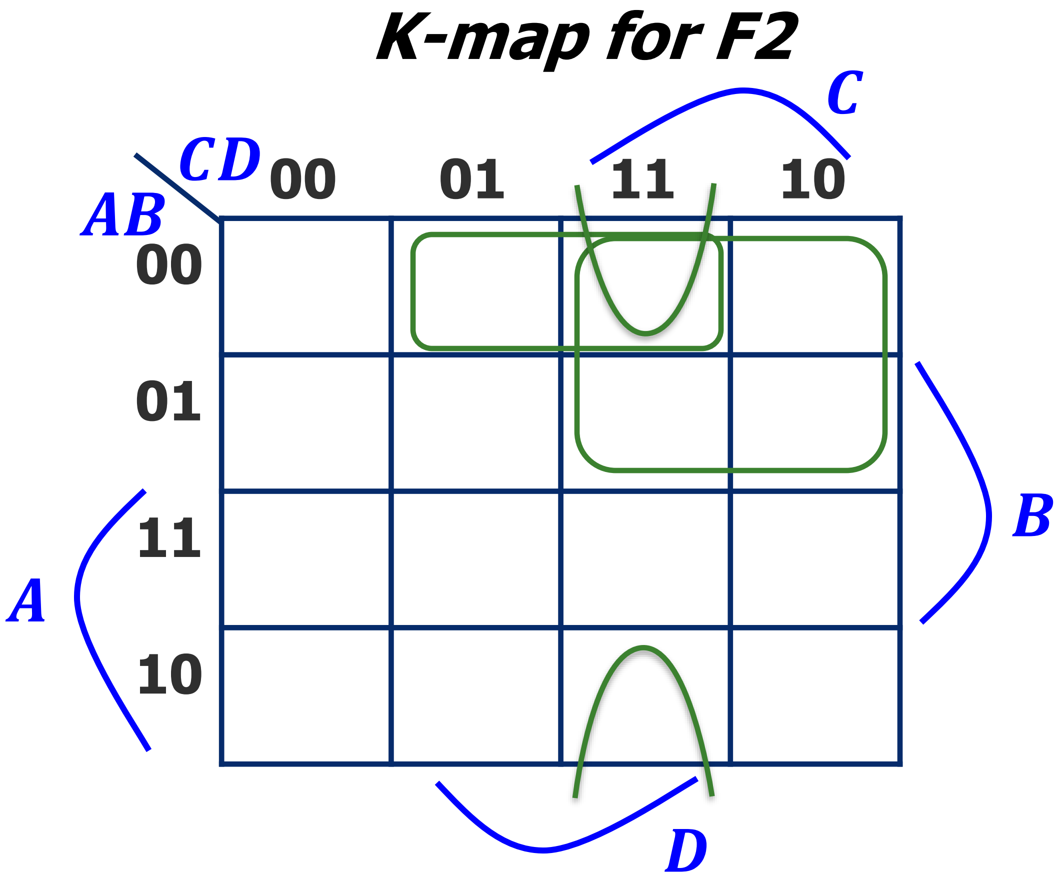
K-map for F1



A	B	C	D	F1	F2	F3
0	0	0	0	1	0	0
0	0	0	1	0	1	0
0	0	1	0	0	1	0
0	0	1	1	0	1	0
0	1	0	0	0	0	1
0	1	0	1	1	0	0
0	1	1	0	0	1	0
0	1	1	1	0	1	0
1	0	0	0	0	0	1
1	0	0	1	0	0	1
1	0	1	0	1	0	0
1	0	1	1	0	1	0
1	1	0	0	0	0	1
1	1	0	1	0	0	1
1	1	1	0	0	0	1
1	1	1	1	1	0	0

F1 =

K-map Example: Two-bit Comparator (3)



F2 =

F3 = ? (Exercise for you)

A	B	C	D	F1	F2	F3
0	0	0	0	1	0	0
0	0	0	1	0	1	0
0	0	1	0	0	1	0
0	0	1	1	0	1	0
0	1	0	0	0	0	1
0	1	0	1	1	0	0
0	1	1	0	0	1	0
0	1	1	1	0	1	0
1	0	0	0	0	0	1
1	0	0	1	0	0	1
1	0	1	0	1	0	0
1	0	1	1	0	1	0
1	1	0	0	0	0	1
1	1	0	1	0	0	1
1	1	1	0	0	0	1
1	1	1	1	1	0	0

K-maps with "Don't Care"

- Don't Care really means *I don't care what my circuit outputs if this appears as input*
  - You have an engineering choice to use DON'T CARE patterns intelligently as 1 or 0 to better simplify the circuit

A	B	C	D	F	G
• • •					
0	1	1	0	X	X
0	1	1	1		
1	0	0	0	X	X
1	0	0	1		
• • •					

I can pick 00, 01, 10, 11 independently of below

I can pick 00, 01, 10, 11 independently of above

Example: BCD Increment Function

- BCD (Binary Coded Decimal) digits
  - Encode decimal digits 0 - 9 with bit patterns 0000_2 — 1001_2
  - When incremented, the decimal sequence is 0, 1, …, 8, 9, 0, 1

A	B	C	D	W	X	Y	Z
0	0	0	0	0	0	0	1
0	0	0	1	0	0	1	0
0	0	1	0	0	0	1	1
0	0	1	1	0	1	0	0
0	1	0	0	0	1	0	1
0	1	0	1	0	1	1	0
0	1	1	0	0	1	1	1
0	1	1	1	1	0	0	0
1	0	0	0	1	0	0	1
1	0	0	1	0	0	0	0
1	0	1	0	X	X	X	X
1	0	1	1	X	X	X	X
1	1	0	0	X	X	X	X
1	1	0	1	X	X	X	X
1	1	1	0	X	X	X	X
1	1	1	1	X	X	X	X

These input patterns should never be encountered in practice (hey -- it's a BCD number!) So, associated output values are "Don't Cares"

K-map for BCD Increment Function

A B
+
W X

Z (without don’t cares) =

Z (with don’t cares) =

10	1		X	X

10			X	X

Y

AB \ CD	00	01	11	10
00		1		1
01		1		1
11	X	X	X	X
10			X	X

Z

AB \ CD	00	01	11	10
00	1			1
01	1			1
11	X	X	X	X
10	1		X	X

K-map Summary

- Karnaugh maps as a formal systematic approach for logic simplification

- 2-, 3-, 4-variable K-maps

- K-maps with "Don't Care" outputs

- H&H Section 2.7